

MIPS® Architecture Extension: nanoMIPS32™ DSP Technical Reference Manual



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About This Book

This chapter describes the terminology and conventions for describing features of the MIPS[®] Architecture such as instructions and control and status registers.

1.1 Typographical Conventions

This section describes the use of *italic*, **bold** and `courier` fonts in this book.

1.1.1 Italic Text

- is used for *emphasis*
- is used for *bits*, *fields*, and *registers* that are important from a software perspective (for instance, address bits used by software, and programmable fields and registers), and various *floating point instruction formats*, such as *S* and *D*
- is used for the memory access types, such as *cached* and *uncached*

1.1.2 Bold Text

- represents a term that is being **defined**
- is used for **bits** and **fields** that are important from a hardware perspective (for instance, **register** bits, which are not programmable but accessible only to hardware)
- is used for ranges of numbers; the range is indicated by an ellipsis. For instance, **5..1** indicates numbers 5 through 1
- is used to emphasize **UNPREDICTABLE** and **UNDEFINED** behavior, as defined below.

1.1.3 Courier Text

`Courier` fixed-width font is used for text that is displayed on the screen, and for examples of code and instruction pseudocode.

1.2 UNPREDICTABLE and UNDEFINED

The terms **UNPREDICTABLE** and **UNDEFINED** are used throughout this book to describe the behavior of the processor in certain cases. **UNDEFINED** behavior or operations can occur only as the result of executing instructions in

a privileged mode (i.e., in Kernel Mode or Debug Mode, or with the CP0 usable bit set in the Status register). Unprivileged software can never cause **UNDEFINED** behavior or operations. Conversely, both privileged and unprivileged software can cause **UNPREDICTABLE** results or operations.

1.2.1 UNPREDICTABLE

UNPREDICTABLE results may vary from processor implementation to implementation, instruction to instruction, or as a function of time on the same implementation or instruction. Software can never depend on results that are **UNPREDICTABLE**. **UNPREDICTABLE** operations may cause a result to be generated or not. If a result is generated, it is **UNPREDICTABLE**. **UNPREDICTABLE** operations may cause arbitrary exceptions.

UNPREDICTABLE results or operations have several implementation restrictions:

- Implementations of operations generating **UNPREDICTABLE** results must not depend on any data source (memory or internal state) which is inaccessible in the current processor mode
- **UNPREDICTABLE** operations must not read, write, or modify the contents of memory or internal state which is inaccessible in the current processor mode. For example, **UNPREDICTABLE** operations executed in user mode must not access memory or internal state that is only accessible in Kernel Mode or Debug Mode or in another process
- **UNPREDICTABLE** operations must not halt or hang the processor

1.2.2 UNDEFINED

UNDEFINED operations or behavior may vary from processor implementation to implementation, instruction to instruction, or as a function of time on the same implementation or instruction. **UNDEFINED** operations or behavior may vary from nothing to creating an environment in which execution can no longer continue. **UNDEFINED** operations or behavior may cause data loss.

UNDEFINED operations or behavior has one implementation restriction:

- **UNDEFINED** operations or behavior must not cause the processor to hang (that is, enter a state from which there is no exit other than powering down the processor). The assertion of any of the reset signals must restore the processor to an operational state

1.2.3 UNSTABLE

UNSTABLE results or values may vary as a function of time on the same implementation or instruction. Unlike **UNPREDICTABLE** values, software may depend on the fact that a sampling of an **UNSTABLE** value results in a legal transient value that was correct at some point in time prior to the sampling.

UNSTABLE values have one implementation restriction:

- Implementations of operations generating **UNSTABLE** results must not depend on any data source (memory or internal state) which is inaccessible in the current processor mode

1.3 Special Symbols in Pseudocode Notation

In this book, algorithmic descriptions of an operation are described using a high-level language pseudocode resembling Pascal. Special symbols used in the pseudocode notation are listed in [Table 1.1](#).

Table 1.1 Symbols Used in Instruction Operation Statements

Symbol	Meaning
♦	Assignment
=, ≠	Tests for equality and inequality
	Bit string concatenation
x^y	A y -bit string formed by y copies of the single-bit value x
$b\#n$	A constant value n in base b . For instance 10#100 represents the decimal value 100, 2#100 represents the binary value 100 (decimal 4), and 16#100 represents the hexadecimal value 100 (decimal 256). If the "b#" prefix is omitted, the default base is 10.
0bn	A constant value n in base 2. For instance 0b100 represents the binary value 100 (decimal 4).
0xn	A constant value n in base 16. For instance 0x100 represents the hexadecimal value 100 (decimal 256).
$x_y z$	Selection of bits y through z of bit string x . Little-endian bit notation (rightmost bit is 0) is used. If y is less than z , this expression is an empty (zero length) bit string.
$x.\text{bit}[y]$	Bit y of bitstring x . Alternative to the traditional MIPS notation x_y .
$x.\text{bits}[y..z]$	Selection of bits y through z of bit string x . Alternative to the traditional MIPS notation $x_y z$.
$x.\text{byte}[y]$	Byte y of bitstring x . Equivalent to the traditional MIPS notation $x_{8*y+7 \ 8*y}$.
$x.\text{bytes}[y..z]$	Selection of bytes y through z of bit string x . Alternative to the traditional MIPS notation $x_{8*y+7 \ 8*z}$.
$x.\text{halfword}[y]$ $x.\text{word}[i]$ $x.\text{doubleword}[i]$	Similar extraction of particular bitfields (used in e.g., MSA packed SIMD vectors).
$x.\text{bit}31, x.\text{byte}0$, etc.	Examples of abbreviated form of $x.\text{bit}[y]$, etc. notation, when y is a constant.
$x.\text{field}_y$	Selection of a named subfield of bitstring x , typically a register or instruction encoding. More formally described as “Field y of register x ”. For example, FIR.D = “the D bit of the Coprocessor 1 Floating-point Implementation Register (FIR)”.
+, −	2’s complement or floating point arithmetic: addition, subtraction
*, ∞	2’s complement or floating point multiplication (both used for either)
div	2’s complement integer division
mod	2’s complement modulo
/	Floating point division
<	2’s complement less-than comparison
>	2’s complement greater-than comparison
≤	2’s complement less-than or equal comparison
≥	2’s complement greater-than or equal comparison
nor	Bitwise logical NOR
xor	Bitwise logical XOR
and	Bitwise logical AND
or	Bitwise logical OR

Table 1.1 Symbols Used in Instruction Operation Statements (Continued)

Symbol	Meaning
not	Bitwise inversion
&&	Logical (non-Bitwise) AND
<<	Logical Shift left (shift in zeros at right-hand-side)
>>	Logical Shift right (shift in zeros at left-hand-side)
GPRLen	The length in bits (32 or 64) of the CPU general-purpose registers
<i>GPR</i> [<i>x</i>]	CPU general-purpose register <i>x</i> . The content of <i>GPR</i> [0] is always zero. In Release 2 of the Architecture, <i>GPR</i> [<i>x</i>] is a short-hand notation for <i>SGPR</i> [<i>SRSCtl</i> _{CSS} , <i>x</i>].
<i>SGPR</i> [<i>s,x</i>]	In Release 2 of the Architecture and subsequent releases, multiple copies of the CPU general-purpose registers may be implemented. <i>SGPR</i> [<i>s,x</i>] refers to GPR set <i>s</i> , register <i>x</i> .
<i>FPR</i> [<i>x</i>]	Floating Point operand register <i>x</i>
<i>FCC</i> [<i>CC</i>]	Floating Point condition code <i>CC</i> . <i>FCC</i> [0] has the same value as <i>COC</i> [1]. Release 6 removes the floating point condition codes.
<i>FPR</i> [<i>x</i>]	Floating Point (Coprocessor unit 1), general register <i>x</i>
<i>CPR</i> [<i>z,x,s</i>]	Coprocessor unit <i>z</i> , general register <i>x</i> , select <i>s</i>
CP2CPR[<i>x</i>]	Coprocessor unit 2, general register <i>x</i>
<i>CCR</i> [<i>z,x</i>]	Coprocessor unit <i>z</i> , control register <i>x</i>
CP2CCR[<i>x</i>]	Coprocessor unit 2, control register <i>x</i>
<i>COC</i> [<i>z</i>]	Coprocessor unit <i>z</i> condition signal
<i>Xlat</i> [<i>x</i>]	Translation of the MIPS16e GPR number <i>x</i> into the corresponding 32-bit GPR number
BigEndianMem	Endian mode as configured at chip reset (0 → Little-Endian, 1 → Big-Endian). Specifies the endianness of the memory interface (see LoadMemory and StoreMemory pseudocode function descriptions) and the endianness of Kernel and Supervisor mode execution.
BigEndianCPU	The endianness for load and store instructions (0 → Little-Endian, 1 → Big-Endian). In User mode, this endianness may be switched by setting the <i>RE</i> bit in the <i>Status</i> register. Thus, BigEndianCPU may be computed as (BigEndianMem XOR ReverseEndian).
ReverseEndian	Signal to reverse the endianness of load and store instructions. This feature is available in User mode only, and is implemented by setting the <i>RE</i> bit of the <i>Status</i> register. Thus, ReverseEndian may be computed as (SR _{RE} and User mode).
<i>LLbit</i>	Bit of virtual state used to specify operation for instructions that provide atomic read-modify-write. <i>LLbit</i> is set when a linked load occurs and is tested by the conditional store. It is cleared, during other CPU operation, when a store to the location would no longer be atomic. In particular, it is cleared by exception return instructions.

Table 1.1 Symbols Used in Instruction Operation Statements (Continued)

Symbol	Meaning						
I , I+n , I-n :	<p>This occurs as a prefix to <i>Operation</i> description lines and functions as a label. It indicates the instruction time during which the pseudocode appears to “execute.” Unless otherwise indicated, all effects of the current instruction appear to occur during the instruction time of the current instruction. No label is equivalent to a time label of I. Sometimes effects of an instruction appear to occur either earlier or later — that is, during the instruction time of another instruction. When this happens, the instruction operation is written in sections labeled with the instruction time, relative to the current instruction I, in which the effect of that pseudocode appears to occur. For example, an instruction may have a result that is not available until after the next instruction. Such an instruction has the portion of the instruction operation description that writes the result register in a section labeled I+1.</p> <p>The effect of pseudocode statements for the current instruction labeled I+1 appears to occur “at the same time” as the effect of pseudocode statements labeled I for the following instruction. Within one pseudocode sequence, the effects of the statements take place in order. However, between sequences of statements for different instructions that occur “at the same time,” there is no defined order. Programs must not depend on a particular order of evaluation between such sections.</p>						
PC	<p>The <i>Program Counter</i> value. During the instruction time of an instruction, this is the address of the instruction word. The address of the instruction that occurs during the next instruction time is determined by assigning a value to <i>PC</i> during an instruction time. If no value is assigned to <i>PC</i> during an instruction time by any pseudocode statement, it is automatically incremented by either 2 (in the case of a 16-bit MIPS16e instruction) or 4 before the next instruction time. A taken branch assigns the target address to the <i>PC</i> during the instruction time of the instruction in the branch delay slot.</p> <p>In the MIPS Architecture, the <i>PC</i> value is only visible indirectly, such as when the processor stores the restart address into a GPR on a jump-and-link or branch-and-link instruction, or into a Coprocessor 0 register on an exception. Release 6 adds <i>PC</i>-relative address computation and load instructions. The <i>PC</i> value contains a full 32-bit address, all of which are significant during a memory reference.</p>						
ISA Mode	<p>In processors that implement the MIPS16e Application Specific Extension or the microMIPS base architectures, the <i>ISA Mode</i> is a single-bit register that determines in which mode the processor is executing, as follows:</p> <table border="1"> <thead> <tr> <th>Encoding</th><th>Meaning</th></tr> </thead> <tbody> <tr> <td>0</td><td>The processor is executing 32-bit MIPS instructions</td></tr> <tr> <td>1</td><td>The processor is executing MIPS16e or microMIPS instructions</td></tr> </tbody> </table> <p>In the MIPS Architecture, the <i>ISA Mode</i> value is only visible indirectly, such as when the processor stores a combined value of the upper bits of <i>PC</i> and the <i>ISA Mode</i> into a GPR on a jump-and-link or branch-and-link instruction, or into a Coprocessor 0 register on an exception.</p>	Encoding	Meaning	0	The processor is executing 32-bit MIPS instructions	1	The processor is executing MIPS16e or microMIPS instructions
Encoding	Meaning						
0	The processor is executing 32-bit MIPS instructions						
1	The processor is executing MIPS16e or microMIPS instructions						
PABITS	<p>The number of physical address bits implemented is represented by the symbol PABITS. As such, if 36 physical address bits were implemented, the size of the physical address space would be $2^{\text{PABITS}} = 2^{36}$ bytes.</p>						
FP32RegistersMode	<p>Indicates whether the FPU has 32-bit or 64-bit floating point registers (FPRs). It is optional if the FPU has 32 64-bit FPRs in which 64-bit data types are stored in any FPR.</p> <p>microMIPS64 implementations have a compatibility mode in which the processor references the FPRs as if it were a microMIPS32 implementation. In such a case FP32RegisterMode is computed from the FR bit in the <i>Status</i> register. If this bit is a 0, the processor operates as if it had 32, 32-bit FPRs. If this bit is a 1, the processor operates with 32 64-bit FPRs.</p> <p>The value of FP32RegistersMode is computed from the FR bit in the <i>Status</i> register.</p>						
InstructionInBranchDelaySlot	<p>Indicates whether the instruction at the Program Counter address was executed in the delay slot of a branch or jump. This condition reflects the <i>dynamic</i> state of the instruction, not the <i>static</i> state. That is, the value is false if a branch or jump occurs to an instruction whose <i>PC</i> immediately follows a branch or jump, but which is not executed in the delay slot of a branch or jump.</p>						

Table 1.1 Symbols Used in Instruction Operation Statements (Continued)

Symbol	Meaning
SignalException(exception, argument)	Causes an exception to be signaled, using the exception parameter as the type of exception and the argument parameter as an exception-specific argument). Control does not return from this pseudocode function—the exception is signaled at the point of the call.

1.4 Notation for Register Field Accessibility

In this document, the read/write properties of register fields use the notations shown in [Table 1.1](#).

Table 1.2 Read/Write Register Field Notation

Read/Write Notation	Hardware Interpretation	Software Interpretation
R/W	<p>A field in which all bits are readable and writable by software and, potentially, by hardware. Hardware updates of this field are visible by software read. Software updates of this field are visible by hardware read.</p> <p>If the Reset State of this field is “Undefined”, either software or hardware must initialize the value before the first read will return a predictable value. This should not be confused with the formal definition of UNDEFINED behavior.</p>	
R	<p>A field which is either static or is updated only by hardware.</p> <p>If the Reset State of this field is either “0”, “Preset”, or “Externally Set”, hardware initializes this field to zero or to the appropriate state, respectively, on powerup. The term “Preset” is used to suggest that the processor establishes the appropriate state, whereas the term “Externally Set” is used to suggest that the state is established via an external source (e.g., personality pins or initialization bit stream). These terms are suggestions only, and are not intended to act as a requirement on the implementation.</p> <p>If the Reset State of this field is “Undefined”, hardware updates this field only under those conditions specified in the description of the field.</p>	<p>A field to which the value written by software is ignored by hardware. Software may write any value to this field without affecting hardware behavior. Software reads of this field return the last value updated by hardware.</p> <p>If the Reset State of this field is “Undefined”, software reads of this field result in an UNPREDICTABLE value except after a hardware update done under the conditions specified in the description of the field.</p>

Table 1.2 Read/Write Register Field Notation (Continued)

Read/Write Notation	Hardware Interpretation	Software Interpretation				
R0	<p>R0 = reserved, read as zero, ignore writes by software.</p> <p>Hardware ignores software writes to an R0 field. Neither the occurrence of such writes, nor the values written, affects hardware behavior.</p> <p>Hardware always returns 0 to software reads of R0 fields.</p> <p>The Reset State of an R0 field must always be 0.</p> <p>If software performs an mtc0 instruction which writes a non-zero value to an R0 field, the write to the R0 field will be ignored, but permitted writes to other fields in the register will not be affected.</p>	<p>Architectural Compatibility: R0 fields are reserved, and may be used for not-yet-defined purposes in future revisions of the architecture.</p> <p>When writing an R0 field, current software should only write either all 0s, or, preferably, write back the same value that was read from the field.</p> <p>Current software should not assume that the value read from R0 fields is zero, because this may not be true on future hardware.</p> <p>Future revisions of the architecture may redefine an R0 field, but must do so in such a way that software which is unaware of the new definition and either writes zeros or writes back the value it has read from the field will continue to work correctly.</p> <p>Writing back the same value that was read is guaranteed to have no unexpected effects on current or future hardware behavior. (Except for non-atomicity of such read-writes.)</p> <p>Writing zeros to an R0 field may not be preferred because in the future this may interfere with the operation of other software which has been updated for the new field definition.</p>				
0	<p style="text-align: center;">Release 6</p> <p>Release 6 legacy “0” behaves like R0 - read as zero, nonzero writes ignored. Legacy “0” should not be defined for any new control register fields; R0 should be used instead.</p> <table><tr><td>HW returns 0 when read. HW ignores writes.</td><td>Only zero should be written, or, value read from register.</td></tr></table> <p style="text-align: center;">pre-Release 6</p> <p>pre-Release 6 legacy “0” - read as zero, nonzero writes UNDEFINED</p> <table><tr><td>A field which hardware does not update, and for which hardware can assume a zero value.</td><td>A field to which the value written by software must be zero. Software writes of non-zero values to this field may result in UNDEFINED behavior of the hardware. Software reads of this field return zero as long as all previous software writes are zero. If the Reset State of this field is “Undefined”, software must write this field with zero before it is guaranteed to read as zero.</td></tr></table>		HW returns 0 when read. HW ignores writes.	Only zero should be written, or, value read from register.	A field which hardware does not update, and for which hardware can assume a zero value.	A field to which the value written by software must be zero. Software writes of non-zero values to this field may result in UNDEFINED behavior of the hardware. Software reads of this field return zero as long as all previous software writes are zero. If the Reset State of this field is “Undefined”, software must write this field with zero before it is guaranteed to read as zero.
HW returns 0 when read. HW ignores writes.	Only zero should be written, or, value read from register.					
A field which hardware does not update, and for which hardware can assume a zero value.	A field to which the value written by software must be zero. Software writes of non-zero values to this field may result in UNDEFINED behavior of the hardware. Software reads of this field return zero as long as all previous software writes are zero. If the Reset State of this field is “Undefined”, software must write this field with zero before it is guaranteed to read as zero.					

Table 1.2 Read/Write Register Field Notation (Continued)

Read/Write Notation	Hardware Interpretation	Software Interpretation
R/W0	<p>Like R/W, except that writes of non-zero to a R/W0 field are ignored. E.g. Status.NMI</p> <p>Hardware may set or clear an R/W0 bit.</p> <p>Hardware ignores software writes of nonzero to an R/W0 field. Neither the occurrence of such writes, nor the values written, affects hardware behavior.</p> <p>Software writes of 0 to an R/W0 field may have an effect.</p> <p>Hardware may return 0 or nonzero to software reads of an R/W0 bit.</p> <p>If software performs an mtc0 instruction which writes a non-zero value to an R/W0 field, the write to the R/W0 field will be ignored, but permitted writes to other fields in the register will not be affected.</p>	<p>Software can only clear an R/W0 bit.</p> <p>Software writes 0 to an R/W0 field to clear the field.</p> <p>Software writes nonzero to an R/W0 bit in order to guarantee that the bit is not affected by the write.</p>

1.5 For More Information

MIPS processor manuals and additional information about MIPS products can be found at <http://www.o-k-u.com>.

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Guide to the Instruction Set

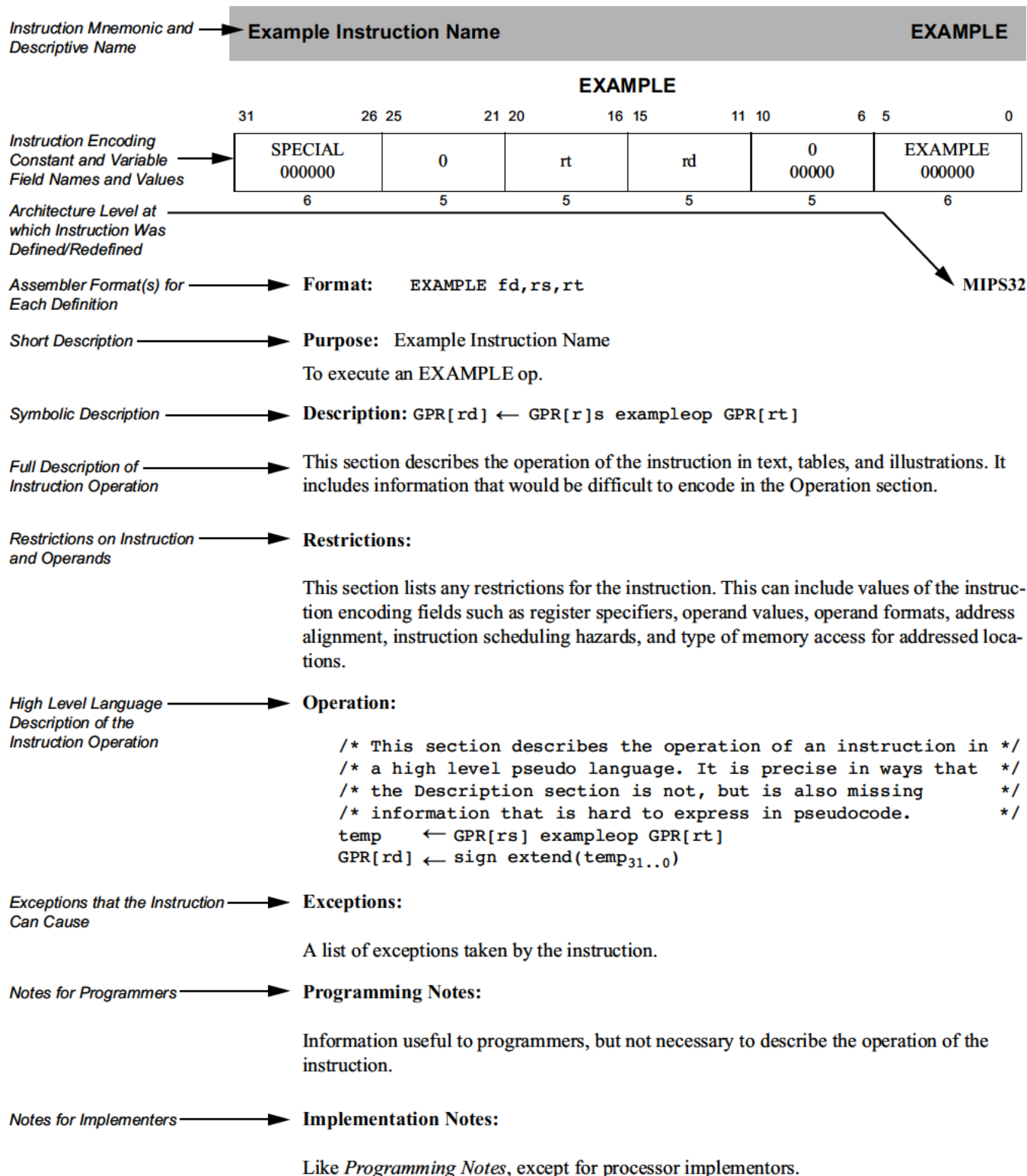
This chapter provides a detailed guide to understanding the instruction descriptions, which are listed in alphabetical order in the tables at the beginning of the next chapter.

2.1 Understanding the Instruction Fields

Figure 2.1 shows an example instruction. Following the figure are descriptions of the fields listed below:

- “Instruction Fields” on page 12
- “Instruction Descriptive Name and Mnemonic” on page 12
- “Format Field” on page 12
- “Purpose Field” on page 13
- “Description Field” on page 13
- “Restrictions Field” on page 13
- “Operation Field” on page 14
- “Exceptions Field” on page 15
- “Programming Notes and Implementation Notes Fields” on page 15

Figure 2.1 Example of Instruction Description

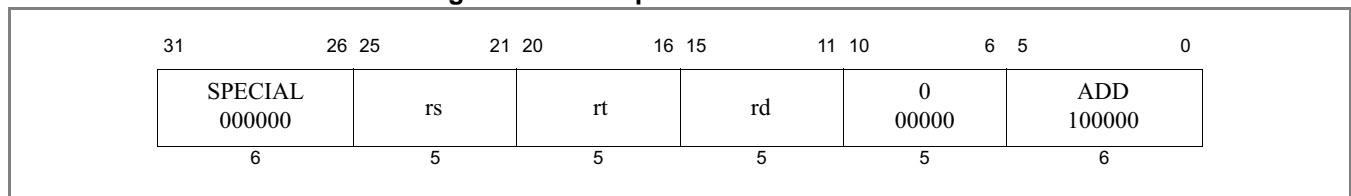


2.1.1 Instruction Fields

Fields encoding the instruction word are shown in register form at the top of the instruction description. The following rules are followed:

- The values of constant fields and the *opcode* names are listed in uppercase (SPECIAL and ADD in Figure 2.2). Constant values in a field are shown in binary below the symbolic or hexadecimal value.
- All variable fields are listed with the lowercase names used in the instruction description (*rs*, *rt*, and *rd* in Figure 2.2).
- Fields that contain zeros but are not named are unused fields that are required to be zero (bits 10:6 in Figure 2.2). If such fields are set to non-zero values, the operation of the processor is **UNPREDICTABLE**.

Figure 2.2 Example of Instruction Fields



2.1.2 Instruction Descriptive Name and Mnemonic

The instruction descriptive name and mnemonic are printed as page headings for each instruction, as shown in Figure 2.3.

Figure 2.3 Example of Instruction Descriptive Name and Mnemonic



2.1.3 Format Field

The assembler formats for the instruction and the architecture level at which the instruction was originally defined are given in the *Format* field. If the instruction definition was later extended, the architecture levels at which it was extended and the assembler formats for the extended definition are shown in their order of extension (for an example, see C.cond fmt). The MIPS architecture levels are inclusive; higher architecture levels include all instructions in previous levels. Extensions to instructions are backwards compatible. The original assembler formats are valid for the extended architecture.

Figure 2.4 Example of Instruction Format

Format:	ADD fd,rs,rt	MIPS32
----------------	--------------	---------------

The assembler format is shown with literal parts of the assembler instruction printed in uppercase characters. The variable parts, the operands, are shown as the lowercase names of the appropriate fields.

The architectural level at which the instruction was first defined, for example “MIPS32” is shown at the right side of the page. Instructions introduced at different times by different ISA family members, are indicated by markings such as “MIPS64, MIPS32 Release 2”. Instructions removed by particular architecture release are indicated in the Availability section.

There can be more than one assembler format for each architecture level. Floating point operations on formatted data show an assembly format with the actual assembler mnemonic for each valid value of the *fmt* field. For example, the ADD *fmt* instruction lists both ADD.S and ADD.D.

The assembler format lines sometimes include parenthetical comments to help explain variations in the formats (once again, see C.cond.fmt). These comments are not a part of the assembler format.

The term *decoded_immediate* is used if the immediate field is encoded within the binary format but the assembler format uses the decoded value. The term *left_shifted_offset* is used if the offset field is encoded within the binary format but the assembler format uses value after the appropriate amount of left shifting.

2.1.4 Purpose Field

The *Purpose* field gives a short description of the use of the instruction.

Figure 2.5 Example of Instruction Purpose

Purpose: Add Word

To add 32-bit integers. If an overflow occurs, then trap.

2.1.5 Description Field

If a one-line symbolic description of the instruction is feasible, it appears immediately to the right of the *Description* heading. The main purpose is to show how fields in the instruction are used in the arithmetic or logical operation.

Figure 2.6 Example of Instruction Description

Description: GPR[rd] ♦ GPR[rs] + GPR[rt]

The 32-bit word value in GPR *rt* is added to the 32-bit value in GPR *rs* to produce a 32-bit result.

- If the addition results in 32-bit 2's complement arithmetic overflow, the destination register is not modified and an Integer Overflow exception occurs.
- If the addition does not overflow, the 32-bit result is placed into GPR *rd*.

The body of the section is a description of the operation of the instruction in text, tables, and figures. This description complements the high-level language description in the *Operation* section.

This section uses acronyms for register descriptions. “GPR *rt*” is CPU general-purpose register specified by the instruction field *rt*. “FPR *fs*” is the floating point operand register specified by the instruction field *fs*. “CP1 register *fd*” is the coprocessor 1 general register specified by the instruction field *fd*. “FCSR” is the floating point *Control / Status* register.

2.1.6 Restrictions Field

The *Restrictions* field documents any possible restrictions that may affect the instruction. Most restrictions fall into one of the following six categories:

- Valid values for instruction fields (for example, see floating point ADD.fmt)
- ALIGNMENT requirements for memory addresses (for example, see LW)

- Valid values of operands (for example, see ALNV.PS)
- Valid operand formats (for example, see floating point ADD.fmt)
- Order of instructions necessary to guarantee correct execution. These ordering constraints avoid pipeline hazards for which some processors do not have hardware interlocks (for example, see MUL).
- Valid memory access types (for example, see LL/SC)

Figure 2.7 Example of Instruction Restrictions**Restrictions:**

None

2.1.7 Availability and Compatibility Fields

The *Availability* and *Compatibility* sections are not provided for all instructions. These sections list considerations relevant to whether and how an implementation may implement some instructions, when software may use such instructions, and how software can determine if an instruction or feature is present. Such considerations include:

- Some instructions are not present on all architecture releases. Sometimes the implementation is required to signal a Reserved Instruction exception, but sometimes executing such an instruction encoding is architecturally defined to give UNPREDICTABLE results.
- Some instructions are available for implementations of a particular architecture release, but may be provided only if an optional feature is implemented. Control register bits typically allow software to determine if the feature is present.
- Some instructions may not behave the same way on all implementations. Typically this involves behavior that was UNPREDICTABLE in some implementations, but which is made architectural and guaranteed consistent so that software can rely on it in subsequent architecture releases.
- Some instructions are prohibited for certain architecture releases and/or optional feature combinations.
- Some instructions may be removed for certain architecture releases. Implementations may then be required to signal a Reserved Instruction exception for the removed instruction encoding; but sometimes the instruction encoding is reused for other instructions.

All of these considerations may apply to the same instruction. If such considerations applicable to an instruction are simple, the architecture level in which an instruction was defined or redefined in the *Format* field, and/or the *Restrictions* section, may be sufficient; but if the set of such considerations applicable to an instruction is complicated, the *Availability* and *Compatibility* sections may be provided.

2.1.8 Operation Field

The *Operation* field describes the operation of the instruction as pseudocode in a high-level language notation resembling Pascal. This formal description complements the *Description* section; it is not complete in itself because many of the restrictions are either difficult to include in the pseudocode or are omitted for legibility.

Figure 2.8 Example of Instruction Operation**Operation:**

```

temp ← (GPR[rs]31 | GPR[rs]31..0) + (GPR[rt]31 | GPR[rt]31..0)
if temp32 ≠ temp31 then
    SignalException(IntegerOverflow)
else
    GPR[rd] ← temp
endif

```

See 2.2 “Operation Section Notation and Functions” on page 15 for more information on the formal notation used here.

2.1.9 Exceptions Field

The *Exceptions* field lists the exceptions that can be caused by *Operation* of the instruction. It omits exceptions that can be caused by the instruction fetch, for instance, TLB Refill, and also omits exceptions that can be caused by asynchronous external events such as an Interrupt. Although a Bus Error exception may be caused by the operation of a load or store instruction, this section does not list Bus Error for load and store instructions because the relationship between load and store instructions and external error indications, like Bus Error, are dependent upon the implementation.

Figure 2.9 Example of Instruction Exception**Exceptions:**

Integer Overflow

An instruction may cause implementation-dependent exceptions that are not present in the *Exceptions* section.

2.1.10 Programming Notes and Implementation Notes Fields

The *Notes* sections contain material that is useful for programmers and implementors, respectively, but that is not necessary to describe the instruction and does not belong in the description sections.

Figure 2.10 Example of Instruction Programming Notes**Programming Notes:**

ADDU performs the same arithmetic operation but does not trap on overflow.

2.2 Operation Section Notation and Functions

In an instruction description, the *Operation* section uses a high-level language notation to describe the operation performed by each instruction. Special symbols used in the pseudocode are described in the previous chapter. Specific pseudocode functions are described below.

This section presents information about the following topics:

- “Instruction Execution Ordering” on page 16
- “Pseudocode Functions” on page 16

2.2.1 Instruction Execution Ordering

Each of the high-level language statements in the *Operations* section are executed sequentially (except as constrained by conditional and loop constructs).

2.2.2 Pseudocode Functions

There are several functions used in the pseudocode descriptions. These are used either to make the pseudocode more readable, to abstract implementation-specific behavior, or both.

These functions are defined in this section, and include the following:

- “Coproprocessor General Register Access Functions” on page 16
- “Memory Operation Functions” on page 17
- “Floating Point Functions” on page 20
- “Instruction Mode Checking Functions” on page 23
- “Miscellaneous Functions” on page 27

2.2.2.1 Coprocessor General Register Access Functions

Defined coprocessors, except for CP0, have instructions to exchange words and doublewords between coprocessor general registers and the rest of the system. What a coprocessor does with a word or doubleword supplied to it and how a coprocessor supplies a word or doubleword is defined by the coprocessor itself. This behavior is abstracted into the functions described in this section.

2.2.2.1.1 COP_LW

The COP_LW function defines the action taken by coprocessor *z* when supplied with a word from memory during a load word operation. The action is coprocessor-specific. The typical action would be to store the contents of memword in coprocessor general register *rt*.

Figure 2.11 COP_LW Pseudocode Function

```
COP_LW (z, rt, memword)
  z: The coprocessor unit number
  rt: Coprocessor general register specifier
  memword: A 32-bit word value supplied to the coprocessor

  /* Coprocessor-dependent action */

endfunction COP_LW
```

2.2.2.1.2 COP_LD

The COP_LD function defines the action taken by coprocessor *z* when supplied with a doubleword from memory during a load doubleword operation. The action is coprocessor-specific. The typical action would be to store the contents of memdouble in coprocessor general register *rt*.

Figure 2.12 COP_LD Pseudocode Function

```
COP_LD (z, rt, memdouble)
  z: The coprocessor unit number
  rt: Coprocessor general register specifier
  memdouble: 64-bit doubleword value supplied to the coprocessor.
```



```

/* Coprocessor-dependent action */

endfunction COP_LD

```

2.2.2.1.3 COP_SW

The COP_SW function defines the action taken by coprocessor *z* to supply a word of data during a store word operation. The action is coprocessor-specific. The typical action would be to supply the contents of the low-order word in coprocessor general register *rt*.

Figure 2.13 COP_SW Pseudocode Function

```

dataword ← COP_SW (z, rt)
  z: The coprocessor unit number
  rt: Coprocessor general register specifier
  dataword: 32-bit word value

/* Coprocessor-dependent action */

endfunction COP_SW

```

2.2.2.1.4 COP_SD

The COP_SD function defines the action taken by coprocessor *z* to supply a doubleword of data during a store doubleword operation. The action is coprocessor-specific. The typical action would be to supply the contents of the low-order doubleword in coprocessor general register *rt*.

Figure 2.14 COP_SD Pseudocode Function

```

datadouble ← COP_SD (z, rt)
  z: The coprocessor unit number
  rt: Coprocessor general register specifier
  datadouble: 64-bit doubleword value

/* Coprocessor-dependent action */

endfunction COP_SD

```

2.2.2.1.5 CoprocessorOperation

The CoprocessorOperation function performs the specified Coprocessor operation.

Figure 2.15 CoprocessorOperation Pseudocode Function

```

CoprocessorOperation (z, cop_fun)

/* z:          Coprocessor unit number */
/* cop_fun:    Coprocessor function from function field of instruction */

/* Transmit the cop_fun value to coprocessor z */

endfunction CoprocessorOperation

```

2.2.2.2 Memory Operation Functions

Regardless of byte ordering (big- or little-endian), the address of a halfword, word, or doubleword is the smallest byte address of the bytes that form the object. For big-endian ordering this is the most-significant byte; for a little-endian ordering this is the least-significant byte.

In the *Operation* pseudocode for load and store operations, the following functions summarize the handling of virtual addresses and the access of physical memory. The size of the data item to be loaded or stored is passed in the *AccessLength* field. The valid constant names and values are shown in Table 2.1. The bytes within the addressed unit of memory (word for 32-bit processors or doubleword for 64-bit processors) that are used can be determined directly from the *AccessLength* and the two or three low-order bits of the address.

2.2.2.2.1 Misaligned Support

MIPS processors originally required all memory accesses to be naturally aligned. MSA (the MIPS SIMD Architecture) supported misaligned memory accesses for its 128 bit packed SIMD vector loads and stores, from its introduction in MIPS Release 5. Release 6 requires systems to provide support for misaligned memory accesses for all ordinary memory reference instructions: the system must provide a mechanism to complete a misaligned memory reference for this instruction, ranging from full execution in hardware to trap-and-emulate.

The pseudocode function `MisalignedSupport` encapsulates the version number check to determine if misalignment is supported for an ordinary memory access.

Figure 2.16 MisalignedSupport Pseudocode Function

```
predicate ← MisalignedSupport ()
    return Config.AR ≥ 2 // Architecture Revision 2 corresponds to MIPS Release 6.
end function
```

See Appendix B, “Misaligned Memory Accesses” on page 511 for a more detailed discussion of misalignment, including pseudocode functions for the actual misaligned memory access.

2.2.2.2.2 AddressTranslation

The `AddressTranslation` function translates a virtual address to a physical address and its cacheability and coherency attribute, describing the mechanism used to resolve the memory reference.

Given the virtual address *vAddr*, and whether the reference is to Instructions or Data (*IorD*), find the corresponding physical address (*pAddr*) and the cacheability and coherency attribute (*CCA*) used to resolve the reference. If the virtual address is in one of the unmapped address spaces, the physical address and *CCA* are determined directly by the virtual address. If the virtual address is in one of the mapped address spaces then the TLB or fixed mapping MMU determines the physical address and access type; if the required translation is not present in the TLB or the desired access is not permitted, the function fails and an exception is taken.

Figure 2.17 AddressTranslation Pseudocode Function

```
(pAddr, CCA) ← AddressTranslation (vAddr, IorD, LorS)

/* pAddr: physical address */
/* CCA: Cacheability&Coherency Attribute, the method used to access caches */
/*      and memory and resolve the reference */

/* vAddr: virtual address */
/* IorD: Indicates whether access is for INSTRUCTION or DATA */
/* LorS: Indicates whether access is for LOAD or STORE */

/* See the address translation description for the appropriate MMU */
/* type in Volume III of this book for the exact translation mechanism */

endfunction AddressTranslation
```

2.2.2.2.3 LoadMemory

The `LoadMemory` function loads a value from memory.

This action uses cache and main memory as specified in both the Cacheability and Coherency Attribute (*CCA*) and the access (*IorD*) to find the contents of *AccessLength* memory bytes, starting at physical location *pAddr*. The data is returned in a fixed-width naturally aligned memory element (*MemElem*). The low-order 2 (or 3) bits of the address and the *AccessLength* indicate which of the bytes within *MemElem* need to be passed to the processor. If the memory access type of the reference is *uncached*, only the referenced bytes are read from memory and marked as valid within the memory element. If the access type is *cached* but the data is not present in cache, an implementation-specific *size* and *alignment* block of memory is read and loaded into the cache to satisfy a load reference. At a minimum, this block is the entire memory element.

Figure 2.18 LoadMemory Pseudocode Function

```
MemElem ← LoadMemory (CCA, AccessLength, pAddr, vAddr, IorD)

/* MemElem:  Data is returned in a fixed width with a natural alignment. The */
/*           width is the same size as the CPU general-purpose register, */
/*           32 or 64 bits, aligned on a 32- or 64-bit boundary, */
/*           respectively. */
/* CCA:      Cacheability&CoherencyAttribute=method used to access caches */
/*           and memory and resolve the reference */

/* AccessLength: Length, in bytes, of access */
/* pAddr:      physical address */
/* vAddr:      virtual address */
/* IorD:       Indicates whether access is for Instructions or Data */

endfunction LoadMemory
```

2.2.2.2.4 StoreMemory

The StoreMemory function stores a value to memory.

The specified data is stored into the physical location *pAddr* using the memory hierarchy (data caches and main memory) as specified by the Cacheability and Coherency Attribute (*CCA*). The *MemElem* contains the data for an aligned, fixed-width memory element (a word for 32-bit processors, a doubleword for 64-bit processors), though only the bytes that are actually stored to memory need be valid. The low-order two (or three) bits of *pAddr* and the *AccessLength* field indicate which of the bytes within the *MemElem* data should be stored; only these bytes in memory will actually be changed.

Figure 2.19 StoreMemory Pseudocode Function

```
StoreMemory (CCA, AccessLength, MemElem, pAddr, vAddr)

/* CCA:      Cacheability&Coherency Attribute, the method used to access */
/*           caches and memory and resolve the reference. */
/* AccessLength: Length, in bytes, of access */
/* MemElem:  Data in the width and alignment of a memory element. */
/*           The width is the same size as the CPU general */
/*           purpose register, either 4 or 8 bytes, */
/*           aligned on a 4- or 8-byte boundary. For a */
/*           partial-memory-element store, only the bytes that will be */
/*           stored must be valid. */
/* pAddr:    physical address */
/* vAddr:    virtual address */

endfunction StoreMemory
```

2.2.2.2.5 Prefetch

The Prefetch function prefetches data from memory.

Prefetch is an advisory instruction for which an implementation-specific action is taken. The action taken may increase performance but must not change the meaning of the program or alter architecturally visible state.

Figure 2.20 Prefetch Pseudocode Function

```
Prefetch (CCA, pAddr, vAddr, DATA, hint)

/* CCA:   Cacheability&Coherency Attribute, the method used to access */
/*        caches and memory and resolve the reference. */
/* pAddr: physical address */
/* vAddr: virtual address */
/* DATA: Indicates that access is for DATA */
/* hint:  hint that indicates the possible use of the data */

endfunction Prefetch
```

Table 2.1 lists the data access lengths and their labels for loads and stores.

Table 2.1 AccessLength Specifications for Loads/Stores

AccessLength Name	Value	Meaning
DOUBLEWORD	7	8 bytes (64 bits)
SEPTIBYTE	6	7 bytes (56 bits)
SEXTIBYTE	5	6 bytes (48 bits)
QUINTIBYTE	4	5 bytes (40 bits)
WORD	3	4 bytes (32 bits)
TRIPLEBYTE	2	3 bytes (24 bits)
HALFWORD	1	2 bytes (16 bits)
BYTE	0	1 byte (8 bits)

2.2.2.2.6 SyncOperation

The SyncOperation function orders loads and stores to synchronize shared memory.

This action makes the effects of the synchronizable loads and stores indicated by *stype* occur in the same order for all processors.

Figure 2.21 SyncOperation Pseudocode Function

```
SyncOperation(stype)

/* stype: Type of load/store ordering to perform. */

/* Perform implementation-dependent operation to complete the */
/* required synchronization operation */

endfunction SyncOperation
```

2.2.2.3 Floating Point Functions

The pseudocode shown in below specifies how the unformatted contents loaded or moved to CP1 registers are interpreted to form a formatted value. If an FPR contains a value in some format, rather than unformatted contents from a load (uninterpreted), it is valid to interpret the value in that format (but not to interpret it in a different format).

2.2.2.3.1 ValueFPR

The ValueFPR function returns a formatted value from the floating point registers.

Figure 2.22 ValueFPR Pseudocode Function

```

value ← ValueFPR(fpr, fmt)

/* value: The formatted value from the FPR */

/* fpr:   The FPR number */
/* fmt:   The format of the data, one of: */
/*        S, D, W, L, PS, */
/*        OB, QH, */
/*        UNINTERPRETED_WORD, */
/*        UNINTERPRETED_DOUBLEWORD */
/* The UNINTERPRETED values are used to indicate that the datatype */
/* is not known as, for example, in SWC1 and SDC1 */

case fmt of
  S, W, UNINTERPRETED_WORD:
    valueFPR ← FPR[fpr]

  D, UNINTERPRETED_DOUBLEWORD:
    if (FP32RegistersMode = 0)
      if (fpr0 ≠ 0) then
        valueFPR ← UNPREDICTABLE
      else
        valueFPR ← FPR[fpr+1]31..0 || FPR[fpr]31..0
      endif
    else
      valueFPR ← FPR[fpr]
    endif

  L:
    if (FP32RegistersMode = 0) then
      valueFPR ← UNPREDICTABLE
    else
      valueFPR ← FPR[fpr]
    endif

  DEFAULT:
    valueFPR ← UNPREDICTABLE

endcase
endfunction ValueFPR

```

The pseudocode shown below specifies the way a binary encoding representing a formatted value is stored into CP1 registers by a computational or move operation. This binary representation is visible to store or move-from instructions. Once an FPR receives a value from the StoreFPR(), it is not valid to interpret the value with ValueFPR() in a different format.

2.2.2.3.2 StoreFPR

Figure 2.23 StoreFPR Pseudocode Function

```

StoreFPR (fpr, fmt, value)

/* fpr:   The FPR number */
/* fmt:   The format of the data, one of: */
/*        S, D, W, L, PS, */

```

```

/*      OB, QH, */
/*      UNINTERPRETED_WORD, */
/*      UNINTERPRETED_DOUBLEWORD */
/* value: The formatted value to be stored into the FPR */

/* The UNINTERPRETED values are used to indicate that the datatype */
/* is not known as, for example, in LWC1 and LDC1 */

case fmt of
  S, W, UNINTERPRETED_WORD:
    FPR[fpr] ← value

  D, UNINTERPRETED_DOUBLEWORD:
    if (FP32RegistersMode = 0)
      if (fpr0 ≠ 0) then
        UNPREDICTABLE
      else
        FPR[fpr] ← UNPREDICTABLE32 || value31..0
        FPR[fpr+1] ← UNPREDICTABLE32 || value63..32
      endif
    else
      FPR[fpr] ← value
    endif

  L:
    if (FP32RegistersMode = 0) then
      UNPREDICTABLE
    else
      FPR[fpr] ← value
    endif

endcase

endfunction StoreFPR

```

2.2.2.3.3 CheckFPEException

The pseudocode shown below checks for an enabled floating point exception and conditionally signals the exception.

Figure 2.24 CheckFPEException Pseudocode Function

```

CheckFPEException()

/* A floating point exception is signaled if the E bit of the Cause field is a 1 */
/* (Unimplemented Operations have no enable) or if any bit in the Cause field */
/* and the corresponding bit in the Enable field are both 1 */

if ( (FCSR17 = 1) or
      ((FCSR16..12 and FCSR11..7) ≠ 0) ) then
  SignalException(FloatingPointException)
endif

endfunction CheckFPEException

```

2.2.2.3.4 FPConditionCode

The FPConditionCode function returns the value of a specific floating point condition code.

Figure 2.25 FPConditionCode Pseudocode Function

```

tf ← FPConditionCode(cc)

/* tf: The value of the specified condition code */

/* cc: The Condition code number in the range 0..7 */

if cc = 0 then
    FPConditionCode ← FCSR23
else
    FPConditionCode ← FCSR24+cc
endif

endfunction FPConditionCode

```

2.2.2.3.5 SetFPConditionCode

The SetFPConditionCode function writes a new value to a specific floating point condition code.

Figure 2.26 SetFPConditionCode Pseudocode Function

```

SetFPConditionCode(cc, tf)
    if cc = 0 then
        FCSR ← FCSR31..24 || tf || FCSR22..0
    else
        FCSR ← FCSR31..25+cc || tf || FCSR23+cc..0
    endif

endfunction SetFPConditionCode

```

2.2.2.4 Instruction Mode Checking Functions**2.2.2.4.1 Are64BitFPOperationsEnabled**

The Are64BitFPOperationsEnabled function is used to determine if a 64-bit floating point instruction may be executed (and conversely, whether a Reserved Instruction exception should be signaled). On a Release 1 processor, such operations are never enabled and this function returns 0. On a Release 2 processor, which supports a 64-bit FPU on a 32-bit processors (and therefore, on a 64-bit processor running with 64-bit operations disabled), the function simply checks the *F64* bit in the *FIR* register.

Figure 2.27 Are64BitFPOperationsEnabled Pseudocode Function

```

enabled ← Are64BitFPOperationsEnabled()

/* enabled: true if 64-bit floating point operations are enabled; */
/* false if they are not */

if (ArchitectureRevision() ≥ 2) then
    Are64BitFPOperationsEnabled ← FIRF64
else
    Are64BitFPOperationsEnabled ← 0
endif

endfunction Are64FPBitOperationsEnabled

```

2.2.2.4.2 IsCoproprocessorEnabled

The IsCoproprocessorEnabled function is used to determine if access is available to one of the four coprocessors. This is primarily done by looking at the value of the appropriate *CU* bit in the *Status* register, but complicated by the fact that access to coprocessor 0 is also enabled if the processor is running in Kernel Mode or Debug Mode.

Figure 2.28 IsCoproprocessorEnabled PseudocodeFunction

```
enabled ← IsCoproprocessorEnabled(z)

/* enabled: true if the coprocessor is enabled; false if it is not */

/* z: The coprocessor unit number in the range 0..3 */

case z of
  0:
    IsCoproprocessorEnabled ←
      (StatusKSU = 0b00) or (DebugDM = 1) or
      (StatusEXL = 1) or (StatusERL = 1)
  1:
    IsCoproprocessorEnabled ← (StatusCU1 = 1)
  2:
    IsCoproprocessorEnabled ← (StatusCU2 = 1)
  3:
    IsCoproprocessorEnabled ← (StatusCU3 = 1)
endcase

endfunction IsCoproprocessorEnabled
```

2.2.2.4.3 IsCoproprocessor2Implemented

The IsCoproprocessor2Implemented function is used to determine if coprocessor 2 is implemented. This is determined by the state of the *C2* bit in the *Config1* register.

Figure 2.29 IsCoproprocessor2 Pseudocode Function

```
impl ← IsCoproprocessor2Implemented()

/* impl: true if coprocessor 2 is implemented; false if it is not */

IsCoproprocessor2Implemented ← Config1C2

endfunction IsCoproprocessor2Implemented
```

2.2.2.4.4 IsEJTAGImplemented

The IsEJTAGImplemented function is used to determine if EJTAG is implemented by the processor. This is determined by the state of the *EP* bit in the *Config1* register.

Figure 2.30 IsEJTAGImplemented Pseudocode Function

```
impl ← IsEJTAGImplemented()

/* impl: true if EJTAG is implemented; false if it is not */

IsEJTAGImplemented ← Config1EP

endfunction IsEJTAGImplemented
```


2.2.2.4.5 IsFloatingPointImplemented

The IsFloatingPointImplemented function is used to determine if floating point is implemented by the processor and, additionally, whether a particular floating point datatype is implemented. Whether floating point is implemented at all is determined by the state of the *FP* bit in the *Config1* register. The determination of whether a particular datatype is implemented is done by looking at the architecture of the chip (MIPS32 or MIPS64, as determined by the *AT* field in the *Config* register), and the state of the *S*, *D*, and *PS* bits in the *FIR* coprocessor 1 register.

Figure 2.31 IsFloatingPointImplemented Pseudocode Function

```
impl ← IsFloatingPointImplemented(fmt)

/* impl: true if floating point is implemented; false if it is not */

/* fmt: The floating point datatype to be checked: */
/*      0: Determine if any floating point datatype is implemented */
/*      S, D, W, L, PS: Determine if a specific datatype is */
/*                      implemented */

if Config1FP = 0 then
    IsFloatingPointImplemented ← 0
else
    case fmt of
        0:
            IsFloatingPointImplemented ← 1
        S:
            IsFloatingPointImplemented ← FIRS

        W:
            IsFloatingPointImplemented ←
                ( ((ArchitectureRevision() = 1) and FIRS)
                  or
                  ((ArchitectureRevision() ≥ 2) and FIRW) )

        D:
            IsFloatingPointImplemented ← FIRD

        L: /* L datatype is valid on a MIPS64 Release 1 implementation */
           /* or on a Release 2 implementation with the L bit set in FIR */
            IsFloatingPointImplemented ←
                ( ((ArchitectureRevision() = 1) and
                  ((ConfigAT = 1) or (Config1AT = 2)))
                  or
                  ((ArchitectureRevision() ≥ 2) and FIRL) )

        PS:
            IsFloatingPointImplemented ← FIRPS and
                ( ((ArchitectureRevision() = 1) and
                  ((ConfigAT = 1) or (Config1AT = 2)))
                  or
                  (ArchitectureRevision() ≥ 2) )
    endcase
endif

endfunction IsFloatingPointImplemented
```

2.2.2.5 Pseudocode Functions Related to Sign and Zero Extension

2.2.2.5.1 Sign extension and zero extension in pseudocode

Much pseudocode uses a generic function `sign_extend` without specifying from what bit position the extension is done, when the intention is obvious. E.g. `sign_extend(immediate16)` or `sign_extend(dispatch9)`.

However, sometimes it is necessary to specify the bit position. For example, `sign_extend(temp31..0)` or the more complicated `(offset15) GPRLen-(16+2) || offset || 02`.

The explicit notation `sign_extend.nbits(val)` or `sign_extend(val, nbits)` is suggested as a simplification. They say to sign extend as if an `nbits`-sized signed integer. The width to be sign extended to is usually apparent by context, and is usually `GPRLen`, 32 or 64 bits. The previous examples then become.

```
sign_extend(temp31..0)
= sign_extend.32(temp)
```

and

```
(offset15) GPRLen-(16+2) || offset || 02
= sign_extend.16(offset) << 2
```

Note that `sign_extend.N(value)` extends from bit position `N-1`, if the bits are numbered `0..N-1` as is typical.

The explicit notations `sign_extend.nbits(val)` or `sign_extend(val, nbits)` is used as a simplification. These notations say to sign extend as if an `nbits`-sized signed integer. The width to be sign extended to is usually apparent by context, and is usually `GPRLen`, 32 or 64 bits.

Figure 2.32 sign_extend Pseudocode Functions

```
sign_extend.nbits(val) = sign_extend(val, nbits) /* syntactic equivalents */

function sign_extend(val, nbits)
  return (val_nbits-1) GPRLen-nbits || val_nbits-1..0
end function
```

The earlier examples can be expressed as

```
(offset15) GPRLen-(16+2) || offset || 02
= sign_extend.16(offset) << 2)
```

and

```
sign_extend(temp31..0)
= sign_extend.32(temp)
```

Similarly for `zero_extension`, although zero extension is less common than sign extension in the MIPS ISA.

Floating point may use notations such as `zero_extend.fmt` corresponding to the format of the FPU instruction. E.g. `zero_extend.S` and `zero_extend.D` are equivalent to `zero_extend.32` and `zero_extend.64`.

Existing pseudocode may use any of these, or other, notations.

2.2.2.5.2 memory_address

The pseudocode function `memory_address` performs mode-dependent address space wrapping for compatibility between MIPS32 and MIPS64. It is applied to all memory references. It may be specified explicitly in some places, particularly for new memory reference instructions, but it is also declared to apply implicitly to all memory references as defined below. In addition, certain instructions that are used to calculate effective memory addresses but which are not themselves memory accesses specify `memory_address` explicitly in their pseudocode.

Figure 2.33 memory_address Pseudocode Function

```
function memory_address(ea)
    return ea
end function
```

On a 32-bit CPU, `memory_address` returns its 32-bit effective address argument unaffected.

In addition to the use of `memory_address` for all memory references (including load and store instructions, LL/SC), Release 6 extends this behavior to control transfers (branch and call instructions), and to the PC-relative address calculation instructions (ADDIUPC, AUIPC, ALUIPC). In newer instructions the function is explicit in the pseudocode.

Implicit address space wrapping for all instruction fetches is described by the following pseudocode fragment which should be considered part of instruction fetch:

Figure 2.34 Instruction Fetch Implicit memory_address Wrapping

```
PC ← memory_address( PC )
( instruction_data, length ) ← instruction_fetch( PC )
/* decode and execute instruction */
```

Implicit address space wrapping for all data memory accesses is described by the following pseudocode, which is inserted at the top of the AddressTranslation pseudocode function:

Figure 2.35 AddressTranslation implicit memory_address Wrapping

```
(pAddr, CCA) ← AddressTranslation( vAddr, IorD, LorS )
vAddr ← memory_address(vAddr)
```

In addition to its use in instruction pseudocode,

2.2.2.6 Miscellaneous Functions

This section lists miscellaneous functions not covered in previous sections.

2.2.2.6.1 SignalException

The `SignalException` function signals an exception condition.

This action results in an exception that aborts the instruction. The instruction operation pseudocode never sees a return from this function call.

Figure 2.36 SignalException Pseudocode Function

```
SignalException(Exception, argument)

/* Exception:    The exception condition that exists. */
/* argument:     A exception-dependent argument, if any */

endfunction SignalException
```

2.2.2.6.2 SignalDebugBreakpointException

The `SignalDebugBreakpointException` function signals a condition that causes entry into Debug Mode from non-Debug Mode.

This action results in an exception that aborts the instruction. The instruction operation pseudocode never sees a return from this function call.

Figure 2.37 SignalDebugBreakpointException Pseudocode Function

```
SignalDebugBreakpointException()

endfunction SignalDebugBreakpointException
```

2.2.2.6.3 SignalDebugModeBreakpointException

The SignalDebugModeBreakpointException function signals a condition that causes entry into Debug Mode from Debug Mode (i.e., an exception generated while already running in Debug Mode).

This action results in an exception that aborts the instruction. The instruction operation pseudocode never sees a return from this function call.

Figure 2.38 SignalDebugModeBreakpointException Pseudocode Function

```
SignalDebugModeBreakpointException()

endfunction SignalDebugModeBreakpointException
```

2.2.2.6.4 NullifyCurrentInstruction

The NullifyCurrentInstruction function nullifies the current instruction.

The instruction is aborted, inhibiting not only the functional effect of the instruction, but also inhibiting all exceptions detected during fetch, decode, or execution of the instruction in question. For branch-likely instructions, nullification kills the instruction in the delay slot of the branch likely instruction.

Figure 2.39 NullifyCurrentInstruction PseudoCode Function

```
NullifyCurrentInstruction()

endfunction NullifyCurrentInstruction
```

2.2.2.6.5 PolyMult

The PolyMult function multiplies two binary polynomial coefficients.

Figure 2.40 PolyMult Pseudocode Function

```
PolyMult(x, y)
    temp ← 0
    for i in 0 .. 31
        if xi = 1 then
            temp ← temp xor (y(31-i)..0 || 0i)
        endif
    endfor

    PolyMult ← temp

endfunction PolyMult
```

2.3 Op and Function Subfield Notation

In some instructions, the instruction subfields *op* and *function* can have constant 5- or 6-bit values. When reference is made to these instructions, uppercase mnemonics are used. For instance, in the floating point ADD instruction, *op*=COP1 and *function*=ADD. In other cases, a single field has both fixed and variable subfields, so the name contains both upper- and lowercase characters.

2.4 FPU Instructions

In the detailed description of each FPU instruction, all variable subfields in an instruction format (such as *fs*, *ft*, *immediate*, and so on) are shown in lowercase. The instruction name (such as ADD, SUB, and so on) is shown in uppercase.

For the sake of clarity, an alias is sometimes used for a variable subfield in the formats of specific instructions. For example, *rs=base* in the format for load and store instructions. Such an alias is always lowercase since it refers to a variable subfield.

Bit encodings for mnemonics are given in Volume I, in the chapters describing the CPU, FPU, MDMX, and MIPS16e instructions.

See “Op and Function Subfield Notation” on page 28 for a description of the *op* and *function* subfields.

The nanoMIPS® DSP Application Specific Extension to the nanoMIPS32® Architecture

3.1 Base Architecture Requirements

The Release 6 nanoMIPS DSP Module requires the implementation of the Release 6 nanoMIPS baseline architecture for support, specifically the Instruction Set and Privileged Resource Architectures.

3.2 Compliance and Subsetting

Instruction subsetting is not allowed for any version of the DSP Module.

3.3 Introduction to the nanoMIPS® DSP Module

This document contains a complete specification of the DSP Module for the nanoMIPS32™ architecture. Statements about DSP Module include MIPS DSP Rev1/2/3 and nanoMIPS DSP except where noted. The table entries in [Chapter 4, “nanoMIPS® DSP Module Instruction Summary” on page 41](#) contain notations which flag the Rev2 instructions, and changes related to nanoMIPS; this information is also available in the per instruction pages. The extensions comprises new integer instructions and new state that includes new HI-LO accumulator pairs and a *DSPControl* register. 32-bit and 64-bit versions of the DSP Module exist which can be included with 32-bit and 64-bit versions of the baseline architecture, respectively.

The Module has been designed to benefit a wide range of DSP, multimedia, and DSP-like algorithms. The performance increase from these extensions can be used to integrate DSP-like functionality into MIPS cores used in a SOC (System on Chip), potentially reducing overall system cost. The Module includes many of the typical features found in other integer-based DSP extensions, for example, support for operations on fractional data types and register SIMD (Single Instruction Multiple Data) operations such as add, subtract, multiply, shift, etc. In addition, the extensions includes some key features that efficiently address specific problems often encountered in DSP applications. These include, for example, support for complex multiplication, variable bit insertion and extraction, and the implementation and use of virtual circular buffers.

This chapter contains a basic overview of the principles behind DSP application processing and the data types and structures needed to efficiently process such applications. [Chapter 4, “nanoMIPS® DSP Module Instruction Summary” on page 41](#), contains a list of all the instructions in the DSP Module arranged by function type. [Chapter 5, “Instruction Encoding” on page 57](#), describes the position of the new instructions in the MIPS instruction opcode map. The rest of the specification contains a complete list of all the instructions that comprise the DSP Module, and serves as a quick reference guide to all the instructions. Finally, various Appendix chapters describe how to implement and use the DSP Module instructions in some common algorithms and inner loops.

3.4 DSP Applications and their Requirements

The DSP Module has been designed specifically to improve the performance of a set of DSP and DSP-like applications. Table 3.1 shows these application areas sorted by the size of the data operands typically preferred by that application for internal computations. For example, raw audio data is usually signed 16-bit, but 32-bit internal calculations are often necessary for high quality audio. (Typically, an internal precision of about 28 bits may be all that is required which can be achieved using a fractional data type of the appropriate width.) There is some cross-over in some cases, which are not explicitly listed here. For example, some hand-held consumer devices may use lower precision internal arithmetic for audio processing, that is, 16-bit internal data formats may be sufficient for the quality required for hand-held devices.

Table 3.1 Data Size of DSP Applications

In/Out Data Size	Internal Data Size	Applications
8 bits	8/16 bits	<ul style="list-style-type: none"> Printer image processing. Still JPEG processing. Moving video processing
16 bits	16 bits	<ul style="list-style-type: none"> Voice Processing. For example, G.723.1, G.729, G.726, echo cancellation, noise cancellation, channel equalization, etc. Soft modem processing. For example V.92. General DSP processing. For example, filters, correlation, convolution, etc.
16/24 bits	32 bits	<ul style="list-style-type: none"> Audio decoding and encoding. For example, MP3, AAC, SRS TruSurround, Dolby Digital Decoder, Pro Logic II, etc.

3.5 Fixed-Point Data Types

Typical implementations of DSP algorithms use fractional fixed-point arithmetic, for reasons of size, cost, and power efficiency. Unlike floating-point arithmetic, fractional fixed-point arithmetic assumes that the position of the decimal point is fixed with respect to the bits representing the fractional value in the operand. To understand this type of arithmetic further, please consult DSP textbooks or other references that are easily available on the internet.

Fractional fixed-point data types are often referred to using Q format notation. The general form for this notation is $Qm.n$, where Q designates that the data is in fractional fixed-point format, m is the number of bits used to designate the twos complement integer portion of the number, and n is the number of bits used to designate the twos complement fractional part of the number. Because the twos complement number is signed, the number of bits required to express a number is $m+n+1$, where the additional bit is required to denote the sign. In typical usage, it is very common for m to be zero. That is, only fractional bits are represented. In this case, a Q notation of the form $Q0.n$ is abbreviated to Qn .

For example, a 32-bit word can be used to represent data in Q31 format, which implies one (left-most) sign bit followed by the binary point and then 31 bits representing the fractional data value. The interpretation of the 32 bits of the Q31 representation is shown in Table 3.2. Negative values are represented using the twos-complement of the equivalent positive value. This format can represent numbers in the range of -1.0 to +0.99999999.... Similarly a 16-bit halfword can be used to represent data in Q15 format, which implies one sign bit followed by 15 fractional bits that represent a value between -1.0 and +0.9999....

Table 3.2 The Value of a Fixed-Point Q31 Number

+	2^{-1}	2^{-2}	2^{-3}	2^{-4}	2^{-5}	2^{-6}	2^{-7}	2^{-8}	2^{-9}	2^{-10}	2^{-11}	2^{-12}	2^{-13}	2^{-14}	2^{-15}	2^{-16}	2^{-17}	2^{-18}	2^{-19}	2^{-20}	2^{-21}	2^{-22}	2^{-23}	2^{-24}	2^{-25}	2^{-26}	2^{-27}	2^{-28}	2^{-29}	2^{-30}	2^{-31}
-																															

Table 3.3 shows the limits of the Q15 and the Q31 representations. Note that the value -1.0 can be represented exactly, but the value +1.0 cannot. For practical purposes, 0x7FFFFFFF is used to represent 1.0 inexactly. Thus, the multiplication of two values where both are -1 will result in an overflow since there is no representation for +1 in fixed-point format. Saturating instructions must check for this case and prevent the overflow by clamping the result to the maximal representable value. Instructions in the DSP Module that operate on fractional data types include a “Q” in the instruction mnemonic; the assumed size of the instruction operands is detailed in the instruction description.

Table 3.3 The Limits of Q15 and Q31 Representations

Fixed-Point Representation	Definition	Hexadecimal Representation	Decimal Equivalent
Q15 minimum	$-2^{15}/2^{15}$	0x8000	-1.0
Q15 maximum	$(2^{15}-1)/2^{15}$	0x7FFF	0.999969482421875
Q31 minimum	$-2^{31}/2^{31}$	0x80000000	-1.0
Q31 maximum	$(2^{31}-1)/2^{31}$	0x7FFFFFFF	0.999999995343387126922607421875

Given a fixed-point representation, we can compute the corresponding decimal value by using bit weights per position as shown in Figure 3.1 for a hypothetical Q7 format number representation with 8 total bits.

DSP applications often, but not always, prefer to saturate the result after an arithmetic operation that causes an overflow or underflow. For operations on signed values, saturation clamps the result to the smallest negative or largest positive value in the case of underflow and overflow, respectively. For operations on unsigned values, saturation clamps the result to either zero or the maximum positive value.

Figure 3.1 Computing the Value of a Fixed-Point (Q7) Number

bit weights	-2^0	2^{-1}	2^{-2}	2^{-3}	2^{-4}	2^{-5}	2^{-6}	2^{-7}	
Example binary value	0	1	1	0	0	1	0	0	decimal value is $2^{-1} + 2^{-2} + 2^{-5}$ $= 0.5 + 0.25 + 0.03125$ $= 0.78125$
Example binary value	0	0	1	1	0	0	0	0	decimal value is $2^{-2} + 2^{-3}$ $= 0.25 + 0.125$ $= 0.375$
Example binary value	0	1	1	1	1	1	1	1	maximum positive value decimal value is $2^{-1} + 2^{-2} + 2^{-3} + 2^{-4} + 2^{-5} + 2^{-6} + 2^{-7}$ $= 0.5 + 0.25 + 0.125 + 0.0625$ $+ 0.03125 + 0.01562 + 0.00781$ $= 0.99218$
Example binary value	1	0	1	0	1	0	0	0	decimal value is $-2^0 + 2^{-2} + 2^{-4}$ $= -1.0 + 0.25 + 0.0625$ $= -0.6875$
Example binary value	1	0	0	0	0	0	0	0	maximum negative value decimal value is -2^0 $= -1.0$

3.6 Saturating Math

Many of the DSP Module arithmetic instructions provide optional saturation of the results, as detailed in each instructions description.

Saturation of fixed-point addition, subtraction, or shift operations that result in an underflow or overflow requires clamping the result value to the closest available fixed-point value representable in the given number of result bits. For operations on unsigned values, underflow is clamped to zero, and overflow to the largest positive fixed-point value. For operations on signed values, underflow is clamped to the minimum negative fixed-point value and overflow to the maximum positive value.

Saturation of fractional fixed-point multiplication operations clamps the result to the maximum representable fixed-point value when both input multiplicands are equal to the minimum negative value of -1.0, which is independent of the Q format used.

3.7 Conventions Used in the Instruction Mnemonics

DSP Module instructions with a **Q** in the mnemonic assume the input operands to be in fractional fixed-point format. Multiplication instructions that operate on fractional fixed-point data will not produce correct results when used with integer fixed-point data. However, addition and subtraction instructions will work correctly with either fractional fixed-point or signed integer fixed-point data.

Instructions that use unsigned data are indicated with the letter **U**. This letter appears after the letter **Q** for fractional in the instruction mnemonic. For example, the **ADDQU** instruction performs an unsigned addition of fractional data. In the MIPS base instruction set, the overflow trap distinguishes signed and unsigned arithmetic instructions. In the DSP Module, the results of saturation distinguish signed and unsigned arithmetic instructions.

Some instructions provide optional rounding up, saturation, or rounding up and saturation of the result(s). These instructions use one of the modifiers **_RS**, **_R**, **_S**, or **_SA** in their mnemonic. For example, **MULQ_RS** is a multiply instruction (**MUL**) where the result is the same size as the input operands (indicated by the absence of **E** for expanded result in the mnemonic) that assumes fractional (**Q**) input data operands, and where the result is rounded up and saturated (**_RS**) before writing the result in the destination register. (For fractional multiplication, saturation clamps the result to the maximum positive representable value if both multiplicands are equal to -1.0.) Several multiply-accumulate (dot product) instructions use a variant of the saturation flag, **_SA**, indicating that the accumulated value is saturated in addition to the regular fractional multiplication saturation check.

The DSP Module instructions provide support for single-instruction, multiple data (SIMD) operations where a single instruction can invoke multiple operation on multiple data operands. As noted previously, DSP applications typically use data types that are 8, 16, or 32 bits wide. In the nanoMIPS32 architecture a general-purpose register (GPR) is 32 bits wide, and in the nanoMIPS64 architecture, 64 bits wide. Thus, each GPR can be used to hold one or more operands of each size. For example, a 64-bit GPR can store eight 8-bit operands, a 32-bit GPR can store two 16-bit operands, and so on. A GPR containing multiple data operands is referred to as a *vector*.

nanoMIPS32 implementations of the DSP Module support three basic formats for data operands: 32 bit, 16 bit, and 8 bit. The latter format is motivated by the fact that video applications typically operate on 8-bit data. The instruction mnemonics indicate the supported data types as follows:

- **W** = “Word”, 1×32 -bit
- **PH** = “Paired Halfword”, 2×16 -bit. See [Figure 3.2](#).
- **QB** = “Quad Byte”, 4×8 -bit. See [Figure 3.3](#).

Figure 3.2 A Paired-Half (PH) Representation in a GPR for the microMIPS32 Architecture

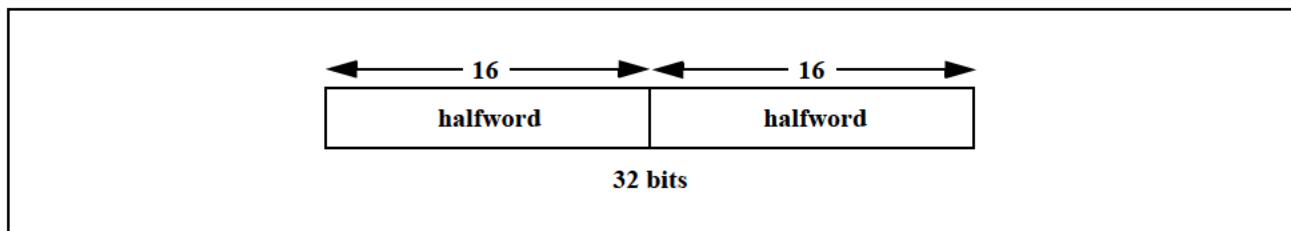
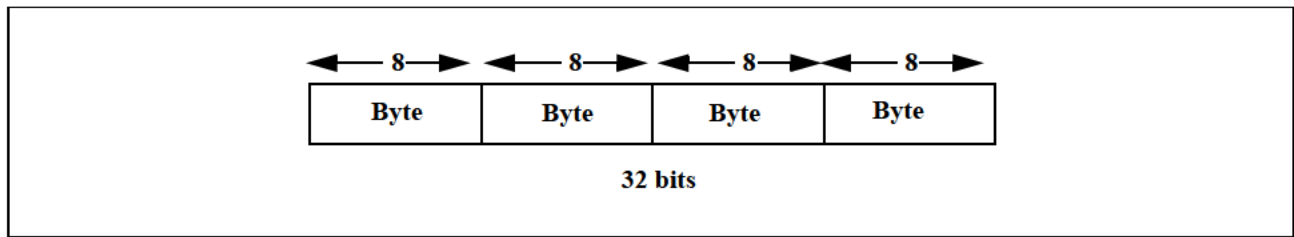
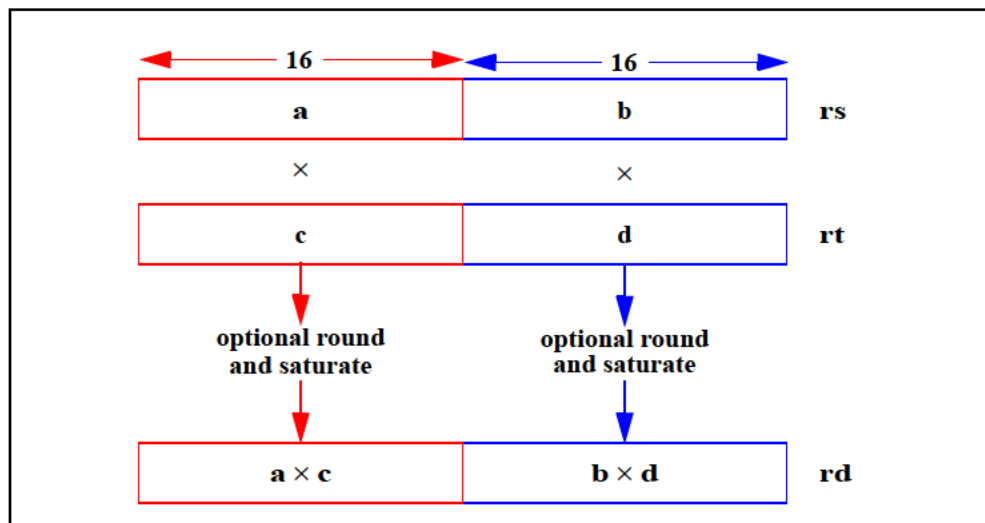


Figure 3.3 A Quad-Byte (QB) Representation in a GPR for the nanoMIPS32 Architecture

For example, **MULQ_RS.PH rd, rs, rt** refers to the multiply instruction (**MUL**) that multiplies two vector elements of type fractional (**Q**) 16 bit (Halfword) data (**PH**) with rounding and saturation (**_RS**). Each source register supplies two data elements and the two results are written into the destination register in the corresponding vector position as shown in [Figure 3.4](#).

When an instruction shows two format types, then the first is the output size and the second is the input size. For example, **PREC_RQ.PH.W** is the (fractional) precision reduction instruction that creates a **PH** output format and uses **W** format as input from the two source registers. When the instruction only shows one format then this implies the same source and destination format.

Figure 3.4 Operation of MULQ_RS.PH rd, rs, rt

3.8 Effect of Endian-ness on Register SIMD Data

The order of data in memory and therefore in the register has a direct impact on the algorithm being executed. To reduce the effort required by the programmer and the development tools to take endian-ness into account, many of the instructions operate on pre-defined bits of a given register. The assembler can be used to map the endian-agnostic names to the actual instructions based on the endian-ness of the processor during the compilation and assembling of the instructions.

When a SIMD vector is loaded into a register or stored back to memory from a register, the endian-ness of the processor and memory has an impact on the view of the data. For example, consider a vector of eight byte values aligned in memory on a 64-bit boundary and loaded into a 64-bit register using the load double instruction: the order of the eight byte values within the register depends on the processor endian-ness. In a big-endian processor, the byte value stored

at the lowest memory address is loaded into the left-most (most-significant) 8 bits of the 64-bit register. In a little-endian processor, the same byte value is loaded into the right-most (least-significant) 8 bits of the register.

In general, if the byte elements are numbered 0-7 according to their order in memory, in a big-endian configuration, element 0 is at the most-significant end and element 7 is at the least-significant end. In a little-endian configuration, the order is reversed. This effect applies to all the sizes of data when they are in SIMD format.

To avoid dealing with the endian-ness issue directly, the instructions in the DSP Module simply refer to the left and right elements of the register when it is required to specify a subset of the elements. This issue can quite easily be dealt with in the assembler or user code using suitably defined mnemonics that use the appropriate instruction for a given endian-ness of the processor. A description of how to do this is specified in [Appendix A](#).

3.9 Additional Register State for the DSP Module

The DSP Module adds four new registers. The operating system is required to recognize the presence of the DSP Module and to include these additional registers in context save and restore operations.

- Three additional *HI-LO* registers to create a total of four accumulator registers. Many common DSP computations involve accumulation, e.g., convolution. DSP Module instructions that target the accumulators use two bits to specify the destination accumulator, with the zero value referring to the original accumulator of the MIPS architecture.

Release 6 of the MIPS Architecture moves the accumulators into the DSP Module for use as a DSP resource exclusively.

- A new control register, *DSPControl*, is used to hold extra state bits needed for efficient support of the new instructions. [Figure 3.5](#) illustrates the bits in this register. [Table 3.4](#) describes the use of the various bits and the instructions that refer to the fields. [Table 3.5](#) lists the instructions that affect the *DSPControl* register *ouflag* field.

Figure 3.5 MIPS® DSP Module Control Register (DSPControl) Format

31	28	27	24	23	16	15	14	13	12	7	6	5	0
0	ccond			ouflag			0	EFI	c	scount			pos

Table 3.4 MIPS® DSP Module Control Register (DSPControl) Field Descriptions

Fields		Description	Read / Write	Reset State	Compliance
Name	Bits				
0	31:28, 15	Not used in the nanoMIPS32 architecture, but these are reserved bits since they are used in the nanoMIPS64 architecture. Must be written as zero; returns zero on read.	0	0	Required
ccond	27:24	Condition code bits set by vector comparison instructions and used as source selectors by PICK instructions. The vector element size determines the number of bits set by a comparison (1, 2, or 4); bits not set are UNPRE-DICTABLE after the comparison.	R/W	0	Required

Table 3.4 MIPS® DSP Module Control Register (DSPControl) Field Descriptions

Fields		Description	Read / Write	Reset State	Compliance
Name	Bits				
ouflag	23:16	Overflow/underflow indication bits set when the result(s) of specific instructions (listed in Table 3.5) caused, or, if optional saturation has been used, would have caused overflow or underflow.	R/W	0	Required
EFI	14	Extract Fail Indicator. This bit is set to 1 when one of the extraction instructions (EXTP, EXTPV, EXTPDP, or EXTPDP) fails. Failure occurs when there are insufficient bits to extract, i.e., when the value of the <i>pos</i> field in the <i>DSPControl</i> register is less than the <i>size</i> argument specified in the instruction. This bit is not sticky—the bit is set or reset after each extraction operation.	R/W	0	Required
c	13	Carry bit set and used by a special add instruction used to implement a 64-bit addition across two GPRs in a nanoMIPS32 implementation. Instruction ADDSC sets the bit and instruction ADDWC uses this bit.	R/W	0	Required
scount	12:7	This field is used by the INSV instruction to specify the size of the bit field to be inserted.	R/W	0	Required
pos	5:0	This field is used by the variable insert instruction INSV to specify the position to insert bits. It is also used to indicate the extract position for the EXTP, EXTPV, EXTPDP, and EXTPDPV instructions. The <i>decrement pos</i> (DP) variants of these instructions decrement the value of the <i>pos</i> field by the amount <i>size</i> +1 after the extraction completes successfully. The MTHLIP instruction increments the value of <i>pos</i> by 32 after copying the value of LO to HI.	R/W	0	Required

The bits of the overflow flag (*ouflag*) field in the *DSPControl* register are set by a number of instructions. These bits are sticky and can be reset only by an explicit write to these bits in the register (using the **WRDSP** instruction). The table below shows which bits can be set by which instructions and under what conditions.

Table 3.5 Instructions that set the ouflag bits in DSPControl

Bit Number	Instructions That Set This Bit
16	Instructions that set this bit when the destination is accumulator (<i>HI-LO</i> pair) zero and an operation overflow or underflow occurs are: DPAQ_S, DPAQ_SA, DPSQ_S, DPSQ_SA, MAQ_S, MAQ_SA, and MULSAQ_S, DPAQX_S, DPAQX_SA, DPSQX_S, DPSQX_SA.
17	Instructions as above, when the destination is accumulator (<i>HI-LO</i> pair) one.
18	Instructions as above, when the destination is accumulator (<i>HI-LO</i> pair) two.

Table 3.5 Instructions that set the ouflag bits in DSPControl

Bit Number	Instructions That Set This Bit
19	Instructions as above, when the destination is accumulator (<i>HI-LO</i> pair) three.
20	Instructions that on an overflow/underflow will set this bit are: ABSQ_S, ADD, ADD_S, ADDQ, ADDQ_S, ADDU, ADDU_S, ADDWC, SUB, SUB_S, SUBQ, SUBQ_S, SUBU, and SUBU_S.
21	Instructions that on an overflow/underflow will set this bit are: MUL, MUL_S, MULEQ_S, MULEU_S, MULQ_RS, and MULQ_S.
22	Instructions that on an overflow/underflow will set this bit are: PRECRQ_RS, PRECRQU_RS, SHLL, SHLL_S, SHLLV, and SHLLV_S.
23	Instructions that on an overflow/underflow will set this bit are: EXTR, EXTR_S, EXTR_RS, EXTRV, EXTRV_RS

3.10 Software Detection of the DSP Module

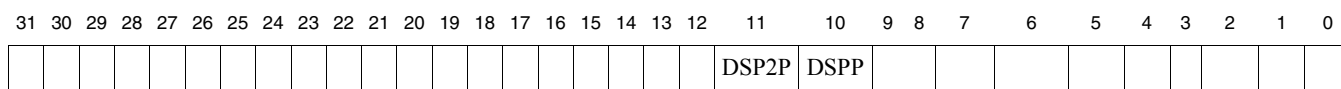
Bit 10 in the *config3* CP0 register, “DSP Present” (DSPP), is used to indicate the presence of the DSP Module Rev1, and bit 11, “DSP Rev2 Present,” (DSP2P), the presence of the DSP Module Rev2, as shown in [Figure 3.6](#). Valid DSP Module Rev2 implementations set both DSPP and DSP2P bits: the condition of DSP2P set and DSPP unset is invalid. Software may read the DSPP, DSP2P bits of the *Config3* CP0 register to check whether this processor has implemented the DSP Module Rev1 and DSP Module Rev2.

Release 6 of the MIPS Architecture moves the accumulators into the DSP Module for use as a DSP resource exclusively, and introduces the compact branch BPOSGE32C, for which DSP Module Rev3 is required. An implementation supports Rev3 if `CP0 Config3DSP=1` and `Config3DSP2=1` and `ConfigAR>=2`.

Software must read `Config3MMAR` to determine if Release 6 nanoMIPS is supported. If `CP0 Config3DSPB=1` and `Config3DSPB=1` and `Config3MMAR>=3`, then Release 6 nanoMIPS DSP is supported.

Any attempt to execute DSP Module instructions must cause a Reserved Instruction Exception if DSPP, and DSP2P are not indicating the presence of the appropriate DSP Module implementation. The DSPP and DSP2P bits are fixed by the hardware implementation and are read-only for software.

Figure 3.6 Config3 Register Format



The “DSP Module Enable” (DSPEn) bit—the MX bit, bit 24 in the CP0 *Status* register as shown in [Figure 3.7](#)—is used to enable access to the extra instructions defined by the DSP Module as well as enabling four modified move instructions (MTLO/HI and MFLO/HI) that provide access to the three additional accumulators *ac1*, *ac2*, and *ac3*. Executing a DSP Module instruction or one of the four modified move instructions when DSPEn is set to zero causes a DSP State Disabled Exception and results in exception code 26 in the CP0 *Cause* register. This allows the OS to do lazy context-switching. [Table 3.6](#) shows the *Cause* Register exception code fields.

Figure 3.7 CP0 Status Register Format

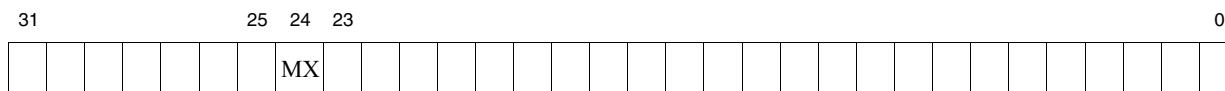


Table 3.6 Cause Register ExcCode Field

Exception Code Value		Mnemonic	Description
Decimal	Hexadecimal		
26	16#1a	DSPDis	DSP Module State Disabled Exception

3.11 Exception Table for the DSP Module

[Table 3.7](#) shows the exceptions caused when a DSP Module or DSP Module Rev2 instruction, MTLO/HI or MFLO/HI, or any other instruction such as an CorExtend instruction attempts to access the new DSP Module state, that is, *ac1*, *ac2*, or *ac3*, or the *DSPControl* register, and all other possible exceptions that relate to the DSP Module.

Implementation Note: Any implementation of the DSP Module must not read or write *ac1*, *ac2*, or *ac3* if *Status_{MX}*=0 for any instruction which might be interpreted as having a field which encodes an accumulator number. Such instructions include:

- DSP Module Rev1, Rev2 instructions
- MADD, MADDU, MSUB, MSUBU, MULT OR MULTU from the base instruction set.
- MADDP, MFLHXU, MTLHX, MULTP, or PPERM from the SmartMIPS® ASE instruction set.

Table 3.7 Exception Table for the DSP Module

<i>Config3_{DSP2P}</i>	<i>Config3_{DSP}</i>	<i>Status_{MX}</i>	Exception for DSP Module Rev2 (or Greater) Instructions	Exception for DSP Module Rev1 Instructions
0	0	×	Reserved Instruction	
0	1	0	Reserved Instruction	DSP Module State Disabled
0	1	1	Reserved Instruction	None
1	1	0	DSP Module State Disabled	
1	1	1	None	
1	1	0	DSP Module State Disabled	
1	1	1	None	

3.12 DSP Module Instructions that Read and Write the DSPControl Register

Many DSP Module instructions read and write the *DSPControl* register, some explicitly and some implicitly. Like other register resource in the architecture, it is the responsibility of the hardware implementation to ensure that appropriate execution dependency barriers are inserted and the pipeline stalled for read-after-write dependencies and other data dependencies that may occur. [Table 3.8](#) lists the DSP Module instructions that can read and write the *DSPControl*

register and the bits or fields in the register that they read or write.

Table 3.8 Instructions that Read/Write Fields in DSPControl

Instruction	Read/Write	DSPControl Field (Bits)
WRDSP	W	All (31:0)
EXTDPDP, EXTPDPV, MTHLIP	W	pos (5:0)
ADDSC	W	c (13)
EXTP, EXTPV, EXTPDP, EXTPDPV	W	EFI (14)
See Table 3.5	W	ouflag (23:16)
CMP, CMPU, and CMPGDU variants	W	ccond (27:24)
RDDSP	R	All (31:0)
BPOSGE32C, EXTP, EXTPV, EXTPDP, EXTPDPV, INSV	R	pos (5:0)
INSV	R	scount (12:7)
ADDWC	R	c (13)
PICK variants	R	ccond (27:24)

3.13 Arithmetic Exceptions

Under no circumstances do any of the DSP Module instructions cause an arithmetic exception. Other exceptions are possible, for example, the indexed load instruction can cause an address exception. The specific exceptions caused by the different instructions are listed in the per-instruction description pages.

nanoMIPS® DSP Module Instruction Summary

4.1 The nanoMIPS® DSP Module Instruction Summary

The tables in this chapter list all the instructions in the DSP Module. For operation details about each instruction, refer to the per-page descriptions. In each table, the column entitled “Writes GPR / ac / *DSPControl*”, indicates the explicit write performed by each instruction. This column indicates the writing of a field in the *DSPControl* register other than the *ouflag* field (which is written by a large number of instructions as a side-effect).

All instructions from the first version of the MIPS® DSP Module onwards are included in the nanoMIPS DSP Module unless explicitly stated otherwise. Release 6 nanoMIPS deprecates BPOSGE32, and replaces PREPEND, BALIGN, LBUX, LHX, LWX by instructions in the baseline nanoMIPS Instruction Set as indicated in the table below.

Table 4.1 List of Instructions in nanoMIPS® DSP Module in Arithmetic Sub-class

Instruction Mnemonics	Input Data Type	Output Data Type	Writes GPR / ac / <i>DSPControl</i>	App	Description
ADDQ.PH rd,rs,rt ADDQ_S.PH rd,rs,rt	Pair Q15	Pair Q15	GPR	VoIP SoftM	Element-wise addition of two vectors of Q15 fractional values, with optional saturation.
ADDQ_S.W rd,rs,rt	Q31	Q31	GPR	Audio	Add two Q31 fractional values with saturation.
ADDU.QB rd,rs,rt ADDU_S.QB rd,rs,rt	Quad Unsigned Byte	Quad Unsigned Byte	GPR	Video	Element-wise addition of unsigned byte values, with optional unsigned saturation.
ADDUH.QB rd,rs,rt ADDUH_R.QB rd,rs,rt Introduced in DSP-R2.	Quad Unsigned Byte	Quad Unsigned Byte	GPR	Video	Element-wise addition of vectors of four unsigned byte values, halving each result by right-shifting by one bit position. Results may be optionally rounded up in the least-significant bit.
ADDU.PH rd,rs,rt ADDU_S.PH rd,rs,rt Introduced in DSP-R2.	Pair Unsigned Halfword	Pair Unsigned Halfword	GPR	Video	Element-wise addition of vectors of two unsigned halfword values, with optional saturation on overflow.
ADDQH.PH rd,rs,rt ADDQH_R.PH rd,rs,rt Introduced in DSP-R2.	Pair Signed Halfword	Pair Signed Halfword	GPR	Misc	Element-wise addition of vectors of two signed halfword values, halving each result with right-shifting by one bit position. Results may be optionally rounded up in the least-significant bit.

Table 4.1 List of Instructions in nanoMIPS® DSP Module in Arithmetic Sub-class (Continued)

Instruction Mnemonics	Input Data Type	Output Data Type	Writes GPR / ac / DSPControl	App	Description
ADDQH.W rd,rs,rt ADDQH_R.W rd,rs,rt Introduced in DSP-R2.	Signed Word	Signed Word	GPR	Misc	Add two signed word values, halving the result with right-shifting by one bit position. Result may be optionally rounded up in the least-significant bit.
SUBQ.PH rd,rs,rt SUBQ_S.PH rd,rs,rt	Pair Q15	Pair Q15	GPR	VoIP	Element-wise subtraction of two vectors of Q15 fractional values, with optional saturation.
SUBQ_S.W rd,rs,rt	Q31	Q31	GPR	Audio	Subtraction with Q31 fractional values, with saturation.
SUBU.QB rd,rs,rt SUBU_S.QB rd,rs,rt	Quad Unsigned Byte	Quad Unsigned Byte	GPR	Video	Element-wise subtraction of unsigned byte values, with optional unsigned saturation.
SUBUH.QB rd,rs,rt SUBUH_R.QB rd,rs,rt Introduced in DSP-R2.	Quad Unsigned Byte	Quad Unsigned Byte	GPR	Video	Element-wise subtraction of unsigned byte values, shifting the results right one bit position (halving). The results may be optionally rounded up by adding 1 to each result at the most-significant discarded bit position before shifting.
SUBU.PH rd,rs,rt SUBU_S.PH rd,rs,rt Introduced in DSP-R2.	Pair Unsigned Halfword	Pair Unsigned Halfword	GPR	Video	Element-wise subtraction of vectors of two unsigned halfword values, with optional saturation on overflow.
SUBQH.PH rd,rs,rt SUBQH_R.PH rd,rs,rt Introduced in DSP-R2.	Pair Signed Halfword	Pair Signed Halfword	GPR	Misc	Element-wise subtraction of vectors of two signed halfword values, halving each result with right-shifting by one bit position. Results may be optionally rounded up in the least-significant bit.
SUBQH.W rd,rs,rt SUBQH_R.W rd,rs,rt Introduced in DSP-R2.	Signed Word	Signed Word	GPR	Misc	Subtract two signed word values, halving the result with right-shifting by one bit position. Result may be optionally rounded up in the least-significant bit.
ADDSC rd,rs,rt	Signed Word	Signed Word	GPR & DSPControl	Audio	Add two signed words and set the carry bit in the DSPControl register.
ADDWC rd,rs,rt	Signed Word	Signed Word	GPR	Audio	Add two signed words with the carry bit from the DSPControl register.
MODSUB rd,rs,rt	Signed Word	Signed Word	GPR	Misc	Modulo addressing support: update a byte index into a circular buffer by subtracting a specified decrement (in bytes) from the index, resetting the index to a specified value if the subtraction results in underflow.

Table 4.1 List of Instructions in nanoMIPS® DSP Module in Arithmetic Sub-class (Continued)

Instruction Mnemonics	Input Data Type	Output Data Type	Writes GPR / ac / DSPControl	App	Description
RADDU.W.QB rd,rs	Quad Unsigned Byte	Unsigned Word	GPR	Misc	Reduce (add together) the 4 unsigned byte values in <i>rs</i> , zero-extending the sum to 32 bits before writing to the destination register. For example, if all 4 input values are 0x80 (decimal 128), then the result in <i>rd</i> is 0x200 (decimal 512).
ABSQ_S.QB rd,rt Introduced in DSP-R2.	Quad Q7	Quad Q7	GPR	Misc	Find the absolute value of each of four Q7 fractional byte elements in the source register, saturating values of -1.0 to the maximum positive Q7 fractional value.
ABSQ_S.PH rd,rt	Pair Q15	Pair Q15	GPR	Misc	Find the absolute value of each of two Q15 fractional halfword elements in the source register, saturating values of -1.0 to the maximum positive Q15 fractional value.
ABSQ_S.W rd,rt	Q31	Q31	GPR	Misc	Find the absolute value of the Q31 fractional element in the source register, saturating the value -1.0 to the maximum positive Q31 fractional value.
PRECR.QB.PH rd,rs,rt Introduced in DSP-R2.	Two Pair Integer Halfwords	Four Integer Bytes	GPR	Misc	Reduce the precision of four signed integer halfword input values by discarding the eight most-significant bits from each to create four signed integer byte output values. The two halfword values from register <i>rs</i> are used to create the two left-most byte results, allowing an endian-agnostic implementation.
PRECRQ.QB.PH rd,rs,rt	2 Pair Q15	Quad Byte	GPR	Misc	Reduce the precision of four Q15 fractional input values by truncation to create four Q7 fractional output values. The two Q15 values from register <i>rs</i> are written to the two left-most byte results, allowing an endian-agnostic implementation.
PRECR_SRA.PH.W rt,rs,sa PRECR_SRA_R.PH.W rt,rs,sa Introduced in DSP-R2.	Two Integer Words	Pair Integer Halfword	GPR	Misc	Reduce the precision of two integer word values to create a pair of integer halfword values. Each word value is first shifted right arithmetically by <i>sa</i> bit positions, and optionally rounded up by adding 1 at the most-significant discard bit position. The 16 least-significant bits of each word are then written to the corresponding halfword elements of destination register <i>rt</i> .

Table 4.1 List of Instructions in nanoMIPS® DSP Module in Arithmetic Sub-class (Continued)

Instruction Mnemonics	Input Data Type	Output Data Type	Writes GPR / ac / DSPControl	App	Description
PRECRQ.PH.W rd,rs,rt PRECRQ_RS.PH.W rd,rs,rt	2 Q31	Pair half-word	GPR	Misc	Reduce the precision of two Q31 fractional input values by truncation to create two Q15 fractional output values. The Q15 value obtained from register <i>rs</i> creates the left-most result, allowing an endian-agnostic implementation. Results may be optionally rounded up and saturated before being written to the destination.
PRECRQU_S.QB.PH rd,rs,rt	2 Pair Q15	Quad Unsigned Byte	GPR	Misc	Reduce the precision of four Q15 fractional values by saturating and truncating to create four unsigned byte values.
PRECEQ.W.PHL rd,rt PRECEQ.W.PHR rd,rt	Q15	Q31	GPR	Misc	Expand the precision of a Q15 fractional value to create a Q31 fractional value by adding 16 least-significant bits to the input value.
PRECEQU.PH.QBL rd,rt PRECEQU.PH.QBR rd,rt PRECEQU.PH.QBLA rd,rt PRECEQU.PH.QBRA rd,rt	Unsigned Byte	Q15	GPR	Video	Expand the precision of two unsigned byte values by prepending a sign bit and adding seven least-significant bits to each to create two Q15 fractional values.
PRECEU.PH.QBL rd,rt PRECEU.PH.QBR rd,rt PRECEU.PH.QBLA rd,rt PRECEU.PH.QBRA rd,rt	Unsigned Byte	Unsigned halfword	GPR	Video	Expand the precision of two unsigned byte values by adding eight least-significant bits to each to create two unsigned halfword values.

Table 4.2 List of Instructions in nanoMIPS® DSP Module in GPR-Based Shift Sub-class

Instruction Mnemonics	Input Data Type	Output Data Type	Writes GPR / ac / DSPControl	App	Description
SHLL.QB rd, rt, sa SHLLV.QB rd, rt, rs	Quad Unsigned Byte	Quad Unsigned Byte	GPR	Misc	Element-wise left shift of eight signed bytes. Zeros are inserted into the bits emptied by the shift. The shift amount is specified by the three least-significant bits of <i>sa</i> or <i>rs</i> .
SHLL.PH rd, rt, sa SHLLV.PH rd, rt, rs SHLL_S.PH rd, rt, sa SHLLV_S.PH rd, rt, rs	Pair Signed halfword	Pair Signed halfword	GPR	Misc	Element-wise left shift of two signed halfwords, with optional saturation on overflow. Zeros are inserted into the bits emptied by the shift. The shift amount is specified by the four least-significant bits of <i>sa</i> or <i>rs</i> .

Table 4.2 List of Instructions in nanoMIPS® DSP Module in GPR-Based Shift Sub-class (Continued)

Instruction Mnemonics	Input Data Type	Output Data Type	Writes GPR / ac / DSPControl	App	Description
SHLL.S.W rd, rt, sa SHLLV.S.W rd, rt, rs	Signed Word	Signed Word	GPR	Misc	Left shift of a signed word, with saturation on overflow. Zeros are inserted into the bits emptied by the shift. The shift amount is specified by the five least-significant bits of <i>sa</i> or <i>rs</i> . Use the microMIPS32 instructions SLL or SLLV for non-saturating shift operations.
SHRL.QB rd, rt, sa SHRLV.QB rd, rt, rs	Quad Unsigned Byte	Quad Unsigned Byte	GPR	Video	Element-wise logical right shift of four byte values. Zeros are inserted into the bits emptied by the shift. The shift amount is specified by the three least-significant bits of <i>sa</i> or <i>rs</i> .
SHRL.PH rd, rt, sa SHRLV.PH rd, rt, rs Introduced in DSP-R2.	Pair Half-words	Pair Half-words	GPR	Video	Element-wise logical right shift of two half-word values. Zeros are inserted into the bits emptied by the shift. The shift amount is specified by the four least-significant bits of <i>rs</i> or the <i>sa</i> argument.
SHRA.QB rd,rt,sa SHRA_R.QB rd,rt,sa SHRAV.QB rd,rt,rs SHRAV_R.QB rd,rt,rs Introduced in DSP-R2.	Quad Byte	Quad Byte	GPR	Misc	Element-wise arithmetic (sign preserving) right shift of four byte values. Optional rounding may be performed, adding 1 at the most-significant discard bit position. The shift amount is specified by the three least-significant bits of <i>rs</i> or by the argument <i>sa</i> .
SHRA.PH rd, rt, sa SHRAV.PH rd, rt, rs SHRA_R.PH rd, rt, sa SHRAV_R.PH rd, rt, rs	Pair Signed halfword	Pair Signed halfword	GPR	Misc	Element-wise arithmetic (sign preserving) right shift of two halfword values. Optionally, rounding may be performed, adding 1 at the most-significant discard bit position. The shift amount is specified by the four least-significant bits of <i>rs</i> or by the argument <i>sa</i> .
SHRA_R.W rd, rt, sa SHRAV_R.W rd, rt, rs	Signed Word	Signed Word	GPR	Video	Arithmetic (sign preserving) right shift of a word value. Optionally, rounding may be performed, adding 1 at the most-significant discard bit position. The shift amount is specified by the five least-significant bits of <i>rs</i> or the argument <i>sa</i> .

Table 4.3 List of Instructions in nanoMIPS® DSP Module in Multiply Sub-class

Instruction Mnemonics	Input Data Type	Output Data Type	Writes GPR / ac / DSPControl	App	Description
MULEU_S.PH.QBL rd,rs,rt MULEU_S.PH.QBR rd,rs,rt	Pair Unsigned Byte, Pair Unsigned Halfword,	Pair Unsigned Halfword	GPR	Still Image	Element-wise multiplication of two unsigned byte values from register <i>rs</i> with two unsigned halfword values from register <i>rt</i> . Each 24-bit product is truncated to 16 bits, with saturation if the product exceeds 0xFFFF, and written to the corresponding element in the destination register.
MULQ_RS.PH rd,rs,rt	Pair Q15	Pair Q15	GPR	Misc	Element-wise multiplication of two Q15 fractional values to create two Q15 fractional results, with rounding and saturation. After multiplication, each 32-bit product is rounded up by adding 0x00008000, then truncated to create a Q15 fractional value that is written to the destination register. If both multiplicands are -1.0, the result is saturated to the maximum positive Q15 fractional value. To stay compliant with the base architecture, this instruction leaves the base <i>HI-LO</i> pair UNPREDICTABLE after the operation. The other DSP Module accumulators <i>ac1-ac3</i> are untouched.
MULEQ_S.W.PHL rd,rs,rt MULEQ_S.W.PHR rd,rs,rt	Pair Q15	Q31	GPR	VoIP	Multiplication of two Q15 fractional values, shifting the product left by 1 bit to create a Q31 fractional result. If both multiplicands are -1.0 the result is saturated to the maximum positive Q31 value. To stay compliant with the base architecture, this instruction leaves the base <i>HI-LO</i> pair UNPREDICTABLE after the operation. The other DSP Module accumulators <i>ac1-ac3</i> must be untouched.
DPAU.H.QBL DPAU.H.QBR	Pair Bytes	Halfword	Acc	Image	Dot-product accumulation. Two pairs of corresponding unsigned byte elements from source registers <i>rt</i> and <i>rs</i> are separately multiplied, and the two 16-bit products are then summed together. The summed products are then added to the accumulator.
DPSU.H.QBL DPSU.H.QBR	Pair Bytes	Halfword	Acc	Image	Dot-product subtraction. Two pairs of corresponding unsigned byte elements from source registers <i>rt</i> and <i>rs</i> are separately multiplied, and the two 16-bit products are then summed together. The summed products are then subtracted from the accumulator.

Table 4.3 List of Instructions in nanoMIPS® DSP Module in Multiply Sub-class (Continued)

Instruction Mnemonics	Input Data Type	Output Data Type	Writes GPR / ac / DSPControl	App	Description
DPA.W.PH ac,rs,rt Introduced in DSP-R2.	Pair Signed Halfword	Pair Signed Halfword	ac	VoIP / SoftM	Dot-product accumulation. The two pairs of corresponding signed integer halfword values from source registers <i>rt</i> and <i>rs</i> are separately multiplied to create two separate integer word products. The products are then summed and accumulated into the specified accumulator.
DPAX.W.PH ac,rs,rt Introduced in DSP-R2.	Pair Signed Halfword	Double-word	ac	VoIP	Dot-product with crossed operands and accumulation. The two crossed pairs of signed integer halfword values from source registers <i>rt</i> and <i>rs</i> are separately multiplied to create two separate integer word products. The products are then summed and accumulated into the specified accumulator.
DPAQ_S.W.PH ac,rs,rt	Pair Q15	Q32.31	ac	VoIP / SoftM	Dot-product accumulation. Two pairs of corresponding Q15 fractional values from source registers <i>rt</i> and <i>rs</i> are separately multiplied and left-shifted 1 bit to create two Q31 fractional products. For each product, if both multipliers are equal to -1.0 the product is clamped to the maximum positive Q31 fractional value. The products are then summed, and the sum is then sign extended to the width of the accumulator and accumulated into the specified accumulator. This instruction may be used to compute the imaginary component of a 16-bit complex multiplication operation after first swapping the operands to place them in the correct order.
DPAQX_S.W.PH ac,rs,rt Introduced in DSP-R2.	Pair Signed Halfword	Q32.31	ac	VoIP	Dot-product with saturating fractional multiplication and using crossed operands, with a final accumulation. The two crossed pairs of signed fractional halfword values from source registers <i>rt</i> and <i>rs</i> are separately multiplied to create two separate fractional word products. The products are then summed and accumulated into the specified accumulator.
DPAQX_SA.W.PH ac,rs,rt Introduced in DSP-R2.	Pair Signed Halfword	Q32.31	ac	VoIP	Dot-product with saturating fractional multiplication and using crossed operands, with a final saturating accumulation. The two crossed pairs of signed fractional halfword values from source registers <i>rt</i> and <i>rs</i> are separately multiplied to create two separate fractional word products. The products are then summed and accumulated with saturation into the specified accumulator.

Table 4.3 List of Instructions in nanoMIPS® DSP Module in Multiply Sub-class (Continued)

Instruction Mnemonics	Input Data Type	Output Data Type	Writes GPR / ac / DSPControl	App	Description
DPS.W.PH ac,rs,rt Introduced in DSP-R2.	Pair Signed Halfword	Double-word	ac	VoIP / SoftM	Dot-product subtraction. The two pairs of corresponding signed integer halfword values from source registers <i>rt</i> and <i>rs</i> are separately multiplied to create two separate integer word products. The products are then summed and subtracted from the specified accumulator.
DPSX.W.PH ac,rs,rt Introduced in DSP-R2.	Pair Signed Halfword	Q32.31	ac	VoIP	Dot-product with crossed operands and subtraction. The two crossed pairs of signed integer halfword values from source registers <i>rt</i> and <i>rs</i> are separately multiplied to create two separate integer word products. The products are then summed and subtracted into the specified accumulator.
DPSQ_S.W.PH ac,rs,rt	Pair Q15	Q32.31	ac	VoIP / SoftM	Dot-product subtraction. Two pairs of corresponding Q15 fractional values from source registers <i>rt</i> and <i>rs</i> are separately multiplied and left-shifted 1 bit to create two Q31 fractional products. For each product, if both multipliers are equal to -1.0 the product is clamped to the maximum positive Q31 fractional value. The products are then summed, and the sum is then sign extended to the width of the accumulator and subtracted from the specified accumulator. This instruction may be used to compute the imaginary component of a 16-bit complex multiplication operation after first swapping the operands to place them in the correct order.
DPSQX_S.W.PH ac,rs,rt Introduced in DSP-R2.	Pair Signed Halfword	Q32.31	ac	VoIP	Dot-product with saturating fractional multiplication and using crossed operands, with a final subtraction. The two crossed pairs of signed fractional halfword values from source registers <i>rt</i> and <i>rs</i> are separately multiplied to create two separate fractional word products. The products are then summed and subtracted from the specified accumulator.
DPSQX_SA.W.PH ac,rs,rt Introduced in DSP-R2.	Pair Signed Halfword	Q32.31	ac	VoIP	Dot-product with saturating fractional multiplication and using crossed operands, with a final saturating subtraction. The two crossed pairs of signed fractional halfword values from source registers <i>rt</i> and <i>rs</i> are separately multiplied to create two separate fractional word products. The products are then summed and subtracted with saturation into the specified accumulator.

Table 4.3 List of Instructions in nanoMIPS® DSP Module in Multiply Sub-class (Continued)

Instruction Mnemonics	Input Data Type	Output Data Type	Writes GPR / ac / DSPControl	App	Description
MULSAQ_S.W.PH ac,rs,rt	Pair Q15	Q32.31	ac	SoftM	Complex multiplication step. Performs element-wise fractional multiplication of the two Q15 fractional values from registers <i>rt</i> and <i>rs</i> , subtracting one product from the other to create a Q31 fractional result that is added to accumulator <i>ac</i> . The intermediate products are saturated to the maximum positive Q31 fractional value if both multiplicands are equal to -1.0.
DPAQ_SA.L.W ac,rs,rt	Q31	Q63	ac	Audio	Fractional multiplication of two Q31 fractional values to produce a Q63 fractional product. If both multiplicands are -1.0 the product is saturated to the maximum positive Q63 fractional value. The product is then added to accumulator <i>ac</i> . If the addition results in overflow or underflow, the accumulator is saturated to the maximum positive or minimum negative value.
DPSQ_SA.L.W ac,rs,rt	Q31	Q63	ac	Audio	Fractional multiplication of two Q31 fractional values to produce a Q63 fractional product. If both multiplicands are -1.0 the product is saturated to the maximum positive Q63 fractional value. The product is then subtracted from accumulator <i>ac</i> . If the addition results in overflow or underflow, the accumulator is saturated to the maximum positive or minimum negative value.
MAQ_S.W.PHL ac,rs,rt MAQ_S.W.PHR ac,rs,rt	Q15	Q32.31	ac	SoftM	Fractional multiply-accumulate. The product of two Q15 fractional values is sign extended to the width of the accumulator and added to accumulator <i>ac</i> . The intermediate product is saturated to the maximum positive Q31 fractional value if both multiplicands are equal to -1.0.
MAQ_SA.W.PHL ac,rs,rt MAQ_SA.W.PHR ac,rs,rt	Q15	Q31	ac	speech	Fractional multiply-accumulate with saturation after accumulation. The product of two Q15 fractional values is sign extended to the width of the accumulator and added to accumulator <i>ac</i> . The intermediate product is saturated to the maximum positive Q31 fractional value if both multiplicands are equal to -1.0. If the accumulation results in overflow or underflow, the accumulator value is saturated to the maximum positive or minimum negative Q31 fractional value.

Table 4.3 List of Instructions in nanoMIPS® DSP Module in Multiply Sub-class (Continued)

Instruction Mnemonics	Input Data Type	Output Data Type	Writes GPR / ac / DSPControl	App	Description
MUL.PH rd,rs,rt MUL_S.PH rd,rs,rt Introduced in DSP-R2.	Pair Signed Halfword	Pair Signed Halfword	GPR	speech	Element-wise multiplication of two vectors of signed integer halfwords, writing the 16 least-significant bits of each 32-bit product to the corresponding element of the destination register. Optional saturation clamps each 16-bit result to the maximum positive or minimum negative value if the product cannot be accurately represented in 16 bits.
MULQ_S.PH rd,rs,rt Introduced in DSP-R2.	Pair Q15	Pair Q15	GPR	speech	Element-wise multiplication of two vectors of Q15 fractional halfwords, writing the 16 most-significant bits of each Q31-format product to the corresponding element of the destination register. Each result is saturated to the maximum positive Q15 value if both multiplicands were equal to -1.0 (0x8000 hexadecimal).
MULQ_S.W rd,rs,rt Introduced in DSP-R2.	Q31	Q31	GPR	speech	Fractional multiplication of two Q31 format words to create a Q63 format result that is truncated by discarding the 32 least-significant bits before being written to the destination register. The result is saturated to the maximum positive Q31 value if both multiplicands were equal to -1.0 (0x80000000 hexadecimal).
MULQ_RS.W rd,rs,rt Introduced in DSP-R2.	Q31	Q31	GPR	speech	Multiplication of two Q31 fractional words to create a Q63-format intermediate product that is rounded up by adding a 1 at bit position 31. The 32 most-significant bits of the rounded result are then written to the destination register. If both multiplicands were equal to -1.0 (0x80000000 hexadecimal), rounding is not performed and the result is clamped to the maximum positive Q31 value before being written to the destination.
MULSA.W.PH ac,rs,rt Introduced in DSP-R2.	Pair Signed Halfword	Double-word	ac	speech	Element-wise multiplication of two vectors of signed integer halfwords to create two 32-bit word intermediate results. The right intermediate result is subtracted from the left intermediate result, and the resulting sum is accumulated into the specified accumulator.
MADD ac,rs,rt MADDU ac,rs,rt MSUB ac,rs,rt MSUBU ac,rs,rt MULT ac,rs,rt MULTU ac,rs,rt	Word	Double-word	ac	Misc	Allows these instructions to target accumulators <i>ac1</i> , <i>ac2</i> , and <i>ac3</i> (in addition to the original <i>ac0</i> destination).

Table 4.4 List of Instructions in MIPS® DSP Module in Bit/ Manipulation Sub-class

Instruction Mnemonics	Input Data Type	Output Data Type	Writes GPR / ac / DSPControl	App	Description
BITREV rd,rt	Unsigned Word	Unsigned Word	GPR	Audio / FFT	Reverse the order of the 16 least-significant bits of register <i>rt</i> , writing the result to register <i>rd</i> . The 16 most-significant bits are set to zero.
INSV rt,rs	Unsigned Word	Unsigned Word	GPR	Misc	Like the Release 2 INS instruction, except that the 5 bits for <i>pos</i> and <i>size</i> values are obtained from the <i>DSPControl</i> register. <i>size</i> = <i>scount</i> [14:10], and <i>pos</i> = <i>pos</i> [20:16].
REPL.QB rd,imm REPLV.QB rd,rt	Byte	Quad Byte	GPR	Video / Misc	Replicate a signed byte value into the four byte elements of register <i>rd</i> . The byte value is given by the 8 least-significant bits of the specified 10-bit immediate constant or by the 8 least-significant bits of register <i>rt</i> .
REPL.PH rd,imm REPLV.PH rd,rt	Signed halfword	Pair Signed halfword	GPR	Misc	Replicate a signed halfword value into the two halfword elements of register <i>rd</i> . The halfword value is given by the 16 least-significant bits of register <i>rt</i> , or by the value of the 10-bit immediate constant, sign-extended to 16 bits.

Table 4.5 List of Instructions in MIPS® DSP Module in Compare-Pick Sub-class

Instruction Mnemonics	Input Data Type	Output Data Type	Writes GPR / ac / DSPControl	App	Description
CMPU.EQ.QB rs,rt CMPU.LT.QB rs,rt CMPU.LE.QB rs,rt	Quad Unsigned Byte	Quad Unsigned Byte	DSPControl	Video	Element-wise unsigned comparison of the four unsigned byte elements of <i>rs</i> and <i>rt</i> , recording the boolean comparison results to the four right-most bits in the <i>ccond</i> field of the <i>DSPControl</i> register.
CMPGDU.EQ.QB rd,rs,rt CMPGDU.LT.QB rd,rs,rt CMPGDU.LE.QB rd,rs,rt Introduced in DSP-R2.	Quad Unsigned Byte	Quad Unsigned Byte	GPR DSPControl	Video	Element-wise unsigned comparison of the four right-most unsigned byte elements of <i>rs</i> and <i>rt</i> , recording the boolean comparison results to the four least-significant bits of register <i>rd</i> and to the four right-most bits in the <i>ccond</i> field of the <i>DSPControl</i> register.
CMPGU.EQ.QB rd,rs,rt CMPGU.LT.QB rd,rs,rt CMPGU.LE.QB rd,rs,rt	Quad Unsigned Byte	Quad Unsigned Byte	GPR	Video	Element-wise unsigned comparison of the four right-most unsigned byte elements of <i>rs</i> and <i>rt</i> , recording the boolean comparison results to the four least-significant bits of register <i>rd</i> .

Table 4.5 List of Instructions in MIPS® DSP Module in Compare-Pick Sub-class (Continued)

Instruction Mnemonics	Input Data Type	Output Data Type	Writes GPR / ac / DSPControl	App	Description
CMP.EQ.PH <i>rs</i> , <i>rt</i> CMP.LT.PH <i>rs</i> , <i>rt</i> CMP.LE.PH <i>rs</i> , <i>rt</i>	Pair Signed halfword	Pair Signed halfword	DSPControl	Misc	Element-wise signed comparison of the two halfword elements of <i>rs</i> and <i>rt</i> , recording the boolean comparison results to the two right-most bits in the <i>ccond</i> field of the <i>DSPControl</i> register.
PICK.QB <i>rd</i> , <i>rs</i> , <i>rt</i>	Quad Unsigned Byte	Quad Unsigned Byte	GPR	Video	Element-wise selection of unsigned bytes from the four bytes of registers <i>rs</i> and <i>rt</i> into the corresponding elements of register <i>rd</i> , based on the value of the four right-most bits of the <i>ccond</i> field in the <i>DSPControl</i> register. If the corresponding <i>ccond</i> bit is 1, the byte value is copied from register <i>rs</i> , otherwise it is copied from <i>rt</i> .
PICK.PH <i>rd</i> , <i>rs</i> , <i>rt</i>	Pair Signed halfword	Pair Signed halfword	GPR	Misc	Element-wise selection of signed halfwords from the two halfwords in registers <i>rs</i> and <i>rt</i> into the corresponding elements of register <i>rd</i> , based on the value of the two right-most bits of the <i>ccond</i> field in the <i>DSPControl</i> register. If the corresponding <i>ccond</i> bit is 1, the halfword value is copied from register <i>rs</i> , otherwise it is copied from <i>rt</i> .
APPEND <i>rt</i> , <i>rs</i> , <i>sa</i> Introduced in DSP-R2.	Two Words	Word	GPR	Misc	Shifts the 32-bit word in register <i>rt</i> left by <i>sa</i> bits, inserting the <i>sa</i> least-significant bits from register <i>rs</i> into the bit positions emptied by the shift. The 32-bit result is then written to register <i>rt</i> .
PREPEND <i>rt</i> , <i>rs</i> , <i>sa</i> Introduced in DSP-R2. Replaced by EXTW in baseline nanoMIPS ISA.	Two Words	Word	GPR	Misc	Shifts the 32-bit word in register <i>rt</i> right by <i>sa</i> bits, inserting the <i>sa</i> least-significant bits from register <i>rs</i> into the bit positions emptied by the shift. The 32-bit result is then written to register <i>rt</i> .
BALIGN <i>rt</i> , <i>rs</i> , <i>bp</i> Introduced in DSP-R2. Replaced by EXTW in baseline nanoMIPS ISA.	Two Words	Word	GPR	Misc	Packs <i>bp</i> bytes from register <i>rt</i> and (4- <i>bp</i>) bytes from register <i>rs</i> into a 32-bit word and writes it to register <i>rt</i> .
PACKRL.PH <i>rd</i> , <i>rs</i> , <i>rt</i>	Pair Signed Halfwords	Pair Signed Halfword	GPR	Misc	Pack two halfwords taken from registers <i>rs</i> and <i>rt</i> into destination register <i>rd</i> .

Table 4.6 List of Instructions in MIPS® DSP Module in Accumulator and DSPControl Access Sub-class

Instruction Mnemonics	Input Data Type	Output Data Type	Writes GPR / ac / DSPControl	App	Description
EXTR.W rt,ac,shift EXTR_R.W rt,ac,shift EXTR_RS.W rt,ac,shift	Q63	Q31	GPR	Misc	Extract a Q31 fractional value from the 32 least-significant bits of 64-bit accumulator <i>ac</i> . The accumulator value may be shifted right logically by <i>shift</i> bits prior to the extraction, and the extracted value may be optionally rounded or rounded and saturated before being written to register <i>rt</i> . The <i>shift</i> argument value ranges from 0 to 31. The optional rounding step adds 1 at the most-significant bit position discarded by the shift. The optional saturation clamps the extracted value to the maximum positive Q31 value if the rounding step results in overflow.
EXTR_S.H rt,ac,shift	Q63	Q15	GPR	Misc	Extract a Q15 fractional value from the 16 least-significant bits of 64-bit accumulator <i>ac</i> . The accumulator value may be shifted right logically by <i>shift</i> bits prior to the extraction, and the extracted value is saturated before being written to register <i>rt</i> . The <i>shift</i> argument value ranges from 0 to 31. The saturation clamps the extracted value to the maximum positive or minimum negative Q15 value if the shifted accumulator value cannot be represented accurately as a Q15 format value.
EXTRV_S.H rt,ac,rs	Q63	Q15	GPR	Misc	Extract a Q15 fractional value from the 16 least-significant bits of 64-bit accumulator <i>ac</i> . The accumulator value may be shifted right logically by <i>shift</i> bits prior to the extraction, and the extracted value is saturated before being written to register <i>rt</i> . The <i>shift</i> argument ranges from 0 to 31 and is given by the five least-significant bits of register <i>rs</i> . The saturation clamps the extracted value to the maximum positive or minimum negative Q15 value if the shifted accumulator value cannot be represented accurately as a Q15 format value.

Table 4.6 List of Instructions in MIPS® DSP Module in Accumulator and DSPControl Access Sub-class

Instruction Mnemonics	Input Data Type	Output Data Type	Writes GPR / ac / DSPControl	App	Description
EXTRV.W rt,ac,rs EXTRV_R.W rt,ac,rs EXTRV_RS.W rt,ac,rs	Q63	Q31	GPR	Misc	Extract a Q31 fractional value from the 32 least-significant bits of 64-bit accumulator <i>ac</i> . The accumulator value may be shifted right logically by <i>shift</i> bits prior to the extraction, and the extracted value may be optionally rounded or rounded and saturated before being written to register <i>rt</i> . The <i>shift</i> argument value is provided by the five least-significant bits of <i>rs</i> and ranges from 0 to 31. The optional rounding step adds 1 at the most-significant bit position discarded by the shift. The optional saturation clamps the extracted value to the maximum positive Q31 value if the rounding step results in overflow.
EXTP rt,ac,size EXTPV rt,ac,rs EXTPDP rt,ac,size EXTPDPV rt,ac,rs	Unsigned DWord	Unsigned Word	GPR / DSPControl	Audio / Video	Extract a set of <i>size</i> +1 contiguous bits from accumulator <i>ac</i> , right-justifying and sign-extending the result to 32 bits before writing the result to register <i>rt</i> . The position of the left-most bit to extract is given by the value of the <i>pos</i> field in the <i>DSPControl</i> register (see Appendix A for details). The number of bits (less one) to extract is provided either by the <i>size</i> immediate operand or by the five least-significant bits of <i>rs</i> . The EXTPDP and EXTPDPV instructions also decrement the <i>pos</i> field by <i>size</i> +1 to facilitate sequential bit field extraction operations.
SHILO ac,shift SHILOV ac,rs	Unsigned DWord	Unsigned DWord	ac	Misc	Shift accumulator <i>ac</i> left or right by the specified number of bits, writing the shifted value back to the accumulator. The signed shift argument is specified either by the immediate operand <i>shift</i> or by the six least-significant bits of register <i>rs</i> . A negative shift argument results in a right shift of up to 32 bits, and a positive shift argument results in a left shift of up to 31 bits.
MTHLIP rs, ac	Unsigned Word	Unsigned Word	ac / DSPControl	Audio / Video	Copy the <i>LO</i> register of the specified accumulator to the <i>HI</i> register, copy <i>rs</i> to <i>LO</i> , and increment the <i>pos</i> field in <i>DSPControl</i> by 32.
MFHI/MFLO/MTHI/MTLO	Unsigned Word	Unsigned Word	GPR/ac	Misc	Copy an unsigned word to or from the specified accumulator <i>HI</i> or <i>LO</i> register to the specified GPR.

Table 4.6 List of Instructions in MIPS® DSP Module in Accumulator and DSPControl Access Sub-class

Instruction Mnemonics	Input Data Type	Output Data Type	Writes GPR / ac / DSPControl	App	Description
WRDSP <i>rt</i> , <i>mask</i>	Unsigned Word	Unsigned Word	<i>DSPControl</i>	Misc	Overwrite specific fields in the <i>DSPControl</i> register using the corresponding bits from the specified GPR. Bits in the <i>mask</i> argument correspond to specific fields in <i>DSPControl</i> ; a value of 1 causes the corresponding <i>DSPControl</i> field to be overwritten using the corresponding bits in <i>rt</i> , otherwise the field is unchanged.
RDDSP <i>rt</i> , <i>mask</i>	Unsigned Word	Unsigned Word	GPR	Misc	Copy the values of specific fields in the <i>DSPControl</i> register to the specified GPR. Bits in the <i>mask</i> argument correspond to specific fields in <i>DSPControl</i> ; a value of 1 causes the corresponding <i>DSPControl</i> field to be copied to the corresponding bits in <i>rt</i> , otherwise the bits in <i>rt</i> are unchanged.

Table 4.7 List of Instructions in MIPS® DSP Module in Indexed-Load Sub-class

Instruction Mnemonics	Input Data Type	Output Data Type	Writes GPR / ac / DSPControl	App	Description
LBUX <i>rd</i> , <i>index</i> (<i>base</i>) Replaced by LBUX in baseline nanoMIPS ISA.	-	Unsigned byte	GPR	Misc	Index byte load from address <i>base</i> +(<i>index</i>). Loads the byte in the low-order bits of the destination register and zero-extends the result.
LHX <i>rd</i> , <i>index</i> (<i>base</i>) Replaced by LHX in baseline nanoMIPS ISA.	-	Signed halfword	GPR	Misc	Index halfword load from address <i>base</i> +(<i>index</i>). Loads the halfword in the low-order bits of the register and sign-extends the result.
LWX <i>rd</i> , <i>index</i> (<i>base</i>) Replaced by LWX in baseline nanoMIPS ISA.	-	Signed Word	GPR	Misc	Indexed word load from address <i>base</i> +(<i>index</i>).

Table 4.8 List of Instructions in MIPS® DSP Module in Branch Sub-class

Instruction Mnemonics	Input Data Type	Output Data Type	Writes GPR / ac / DSPControl	App	Description
BPOSGE32 offset Deprecated in nanoMIPS DSP. BPOSGE32C offset Introduced in DSP-R3.	-	-	-	Audio / Video	Branch if the <i>pos</i> value is greater than or equal to integer 32.

Instruction Encoding

The opcode map for DSP instructions is under development and will be made available in a subsequent non-preliminary release.

5.1 Instruction Bit Encoding

This chapter describes the bit encoding tables used for the MIPS DSP ASE. Table 5.1 describes the meaning of the symbols used in the tables. These tables only list the instruction encoding for the MIPS DSP ASE instructions. See Volumes I and II of this multi-volume set for a full encoding of all instructions.

Table 5.1 Symbols Used in the Instruction Encoding Tables

Symbol	Meaning
*	Operation or field codes marked with this symbol are reserved for future use. Executing such an instruction must cause a Reserved Instruction Exception.
δ	(Also <i>italic</i> field name.) Operation or field codes marked with this symbol denotes a field class. The instruction word must be further decoded by examining additional tables that show values for another instruction field.
β	Operation or field codes marked with this symbol represent a valid encoding for a higher-order MIPS ISA level. Executing such an instruction must cause a Reserved Instruction Exception.
θ	Operation or field codes marked with this symbol are available to licensed MIPS partners. To avoid multiple conflicting instruction definitions, MIPS Technologies will assist the partner in selecting appropriate encoding if requested by the partner. The partner is not required to consult with MIPS Technologies when one of these encoding is used. If no instruction is encoded with this value, executing such an instruction must cause a Reserved Instruction Exception (<i>SPECIAL2</i> encoding or coprocessor instruction encoding for a coprocessor to which access is allowed) or a Coprocessor Unusable Exception (coprocessor instruction encoding for a coprocessor to which access is not allowed).
σ	Field codes marked with this symbol represent an EJTAG support instruction and implementation of this encoding is optional for each implementation. If the encoding is not implemented, executing such an instruction must cause a Reserved Instruction Exception. If the encoding is implemented, it must match the instruction encoding as shown in the table.
ε	Operation or field codes marked with this symbol are reserved for MIPS Application Specific Extensions. If the ASE is not implemented, executing such an instruction must cause a Reserved Instruction Exception.
ϕ	Operation or field codes marked with this symbol are obsolete and will be removed from a future revision of the MIPS32 ISA. Software should avoid using these operation or field codes.
\oplus	Operation or field codes marked with this symbol are valid for Release 2 implementations of the architecture. Executing such an instruction in a Release 1 implementation must cause a Reserved Instruction Exception.

The MIPS® DSP Module Instruction Set

6.1 Compliance and Subsetting

There are no instruction subsets allowed for the MIPS DSP Module—all instructions must be implemented with all data format types as shown. Instructions are listed in alphabetical order, with a secondary sort on data type format from narrowest to widest, i.e., quad byte, paired halfword, and word.

6.2 DSP Module Specific Pseudocode Functions

This section defines the pseudocode functions that are specific to the DSP Module and DSP Module Rev2. These functions are used in the Operation section of each DSP Module instruction description.

6.2.1 ValidateAccessToDSPResources()

The ValidateAccessToDSPResources function is used to determine if access is available to the DSP Module resources. This is done by looking at the state of the DSPP bit in *Config3* and MX bit in the *Status* register.

Figure 6.1 ValidateAccessToDSPResource Pseudocode Function

```
ValidateAccessToDSPResources()

/* The function does not return if an exception is signaled */

/* If DSP is not implemented by the processor, a Reserved */
/* Instruction exception is signaled */
if (Config3_DSPP = 0) then
    SignalException(ReservedInstruction)
endif

case Status_MX of

    /* MX off */
    1#0:
        SignalException(DSPDisabled)

    /* MX on */
    1#1:
        /* Access allowed to DSP Module resources */
endcase

endfunction ValidateAccessToDSPResources
```

6.2.2 ValidateAccessToDSP2Resources()

The `ValidateAccessToDSP2Resources` function is used to determine if access is available to the DSP Module Rev2 resources. This is done by checking the state of the DSP2P bit (DSP Rev2 Present, bit 11 in the *Config3* CPO register), and the MX bit in the *Status* register.

Figure 6.2 ValidateAccessToDSP2Resources Pseudocode Function

```

ValidateAccessToDSP2Resources()

/* The function does not return if an exception is signaled */

/* If DSP Module Rev2 is not implemented by the processor, a */
/* Reserved Instruction exception is signaled */
if ((Config3_DSP2P = 0) or (Config3_DSPP = 0)) then
    SignalException(ReservedInstruction)
endif

case Status_MX of

    /* MX off */
    1#0:
        SignalException(DSPDisabled)

    /* MX on */
    1#1:
        /* Access allowed to DSP Module Rev2 resources */
endcase

endfunction ValidateAccessToDSP2Resources

```

ABSQ_S.PH**Find Absolute Value of Two Fractional Halfwords**

31	26	25	21	20	16	15	9	8	6	5	3	2	0		
P32A 001000			rt		rs		0001000			100		111		111	
6			5		5		7			3		3		3	

Format: ABSQ_S.PH rt, rs**DSP****Purpose:** Find Absolute Value of Two Fractional Halfwords

Find the absolute value of each of a pair of Q15 fractional halfword values with 16-bit saturation.

Description: $rt \leftarrow \text{sat16}(\text{abs}(rs_{31..16})) \parallel \text{sat16}(\text{abs}(rs_{15..0}))$

For each value in the pair of Q15 fractional halfword values in register *rs*, the absolute value is found and written to the corresponding Q15 halfword in register *rt*. If either input value is the minimum Q15 value (-1.0 in decimal, 0x8000 in hexadecimal), the corresponding result is saturated to 0x7FFF.

This instruction sets bit 20 in the *DSPControl* register in the *ouflag* field if either input value was saturated.

Restrictions:

No data-dependent exceptions are possible.

The operands must be values in the specified format. If they are not, the results are **UNPREDICTABLE** and the values of the operand vectors become **UNPREDICTABLE**.

Operation:

```

ValidateAccessToDSPResources()
tempB15..0 ← satAbs16( GPR[rs]31..16 )
tempA15..0 ← satAbs16( GPR[rs]15..0 )
GPR[rt]31..0 ← tempB15..0 || tempA15..0

```

```

function satAbs16( a15..0 )
  if ( a15..0 = 0x8000 ) then
    DSPControlouflag:20 ← 1
    temp15..0 ← 0x7FFF
  else
    if ( a15 = 1 ) then
      temp15..0 ← -a15..0
    else
      temp15..0 ← a15..0
    endif
  endif
  return temp15..0
endfunction satAbs16

```

Exceptions:

Reserved Instruction, DSP Disabled

ABSQ_S.PH

Find Absolute Value of Two Fractional Halfwords

ABSQ_S.QB**Find Absolute Value of Four Fractional Byte Values**

31	26	25	21	20	16	15	9	8	6	5	3	2	0	
P32A 001000			rt		rs		0000000		100		111		111	
6			5		5		7		3		3		3	

Format: ABSQ_S.QB rt, rs**DSP-R2****Purpose:** Find Absolute Value of Four Fractional Byte Values

Find the absolute value of four fractional byte vector elements with saturation.

Description: $rt \leftarrow \text{sat8}(\text{abs}(rs_{31..24})) \parallel \text{sat8}(\text{abs}(rs_{23..16})) \parallel \text{sat8}(\text{abs}(rs_{15..8})) \parallel \text{sat8}(\text{abs}(rs_{7..0}))$

For each value in the four Q7 fractional byte elements in register *rs*, the absolute value is found and written to the corresponding byte in register *rt*. If either input value is the minimum Q7 value (-1.0 in decimal, 0x80 in hexadecimal), the corresponding result is saturated to 0x7F.s

This instruction sets bit 20 in *ouflag* field of the *DSPControl* register if any input value was saturated.

Restrictions:

No data-dependent exceptions are possible.

The operands must be a value in the specified format. If they are not, the results are **UNPREDICTABLE** and the values of the operand vectors become **UNPREDICTABLE**.

Operation:

```

ValidateAccessToDSP2Resources()
tempD7..0 ← abs8( GPR[rs]31..24 )
tempC7..0 ← abs8( GPR[rs]23..16 )
tempB7..0 ← abs8( GPR[rs]15..8 )
tempA7..0 ← abs8( GPR[rs]7..0 )
GPR[rt]31..0 ← tempD7..0 || tempC7..0 || tempB7..0 || tempA7..0

function abs8( a7..0 )
  if ( a7..0 = 0x80 ) then
    DSPControl_ouflag:20 ← 1
    temp7..0 ← 0x7F
  else
    if ( a7 = 1 ) then
      temp7..0 ← -a7..0
    else
      temp7..0 ← a7..0
    endif
  endif
  return temp7..0
endfunction abs8

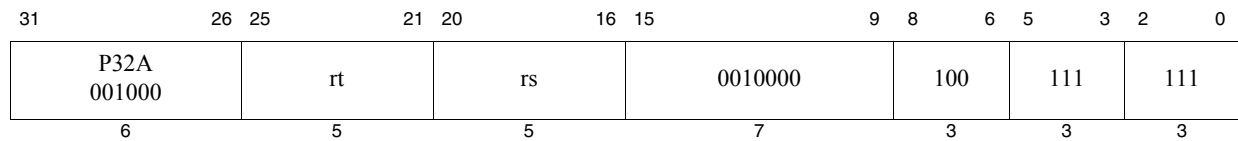
```

Exceptions:

Reserved Instruction, DSP Disabled

ABSQ_S.QB

Find Absolute Value of Four Fractional Byte Values

ABSQ_S.W**Find Absolute Value of Fractional Word****Format:** ABSQ_S.W rt, rs**DSP****Purpose:** Find Absolute Value of Fractional Word

Find the absolute value of a fractional Q31 value with 32-bit saturation.

Description: $rt \leftarrow \text{sat32}(\text{abs}(rs_{31..0}))$

The absolute value of the Q31 fractional value in register *rs* is found and written to destination register *rt*. If the input value is the minimum Q31 value (-1.0 in decimal, 0x80000000 in hexadecimal), the result is saturated to 0x7FFFFFFF before being sign-extended and written to register *rt*.

This instruction sets bit 20 in the *DSPControl* register in the *ouflag* field if the input value was saturated.

Restrictions:

No data-dependent exceptions are possible.

The operands must be values in the specified format. If they are not, the results are **UNPREDICTABLE** and the values of the operand vectors become **UNPREDICTABLE**.

Operation:

```

ValidateAccessToDSPResources()
temp31..0 ← satAbs32( GPR[rs]31..0 )
GPR[rt]31..0 ← temp31..0

function satAbs32( a31..0 )
  if ( a31..0 = 0x80000000 ) then
    DSPControlouflag:20 ← 1
    temp31..0 ← 0x7FFFFFFF
  else
    if ( a31 = 1 ) then
      temp31..0 ← -a31..0
    else
      temp31..0 ← a31..0
    endif
  endif
  return temp31..0
endfunction satAbs32

```

Exceptions:

Reserved Instruction, DSP Disabled

ABSQ_S.W

Find Absolute Value of Fractional Word

ADDQ[_S].PH**Add Fractional Halfword Vectors**

31	26	25	21	20	16	15	11	10	9	3	2	0
ADDQ.PH												
P32A 001000		rt		rs		rd		0	0000001		101	
ADDQ_S.PH												
P32A 001000		rt		rs		rd		1	0000001		101	
6		5		5		5		1	7		3	

Format: ADDQ[_S].PH
 ADDQ.PH rd, rs, rt
 ADDQ_S.PH rd, rs, rt

DSP
DSP

Purpose: Add Fractional Halfword Vectors

Element-wise addition of two vectors of Q15 fractional values to produce a vector of Q15 fractional results, with optional saturation.

Description: $rd \leftarrow \text{sat16}(rs_{31..16} + rt_{31..16}) \parallel \text{sat16}(rs_{15..0} + rt_{15..0})$

Each of the two fractional halfword elements in register *rt* are added to the corresponding fractional halfword elements in register *rs*.

For the non-saturating version of the instruction, the result of each addition is written into the corresponding element in register *rd*. If the addition results in overflow or underflow, the result modulo 2 is written to the corresponding element in register *rd*.

For the saturating version of the instruction, signed saturating arithmetic is performed, where an overflow is clamped to the largest representable value (0x7FFF hexadecimal) and an underflow to the smallest representable value (0x8000 hexadecimal) before being written to the destination register *rd*.

For each instruction, if either of the individual additions result in underflow, overflow, or saturation, a 1 is written to bit 20 in the *DSPControl* register in the ouflag field.

Restrictions:

No data-dependent exceptions are possible.

The operands must be values in the specified format. If they are not, the results are **UNPREDICTABLE** and the values of the operand vectors become **UNPREDICTABLE**.

Operation:

```

ADDQ.PH:
  ValidateAccessToDSPResources()
  tempB15..0 ← add16( GPR[rs]31..16 , GPR[rt]31..16 )
  tempA15..0 ← add16( GPR[rs]15..0 , GPR[rt]15..0 )
  GPR[rd]31..0 ← tempB15..0 || tempA15..0

ADDQ_S.PH:
  ValidateAccessToDSPResources()
  tempB15..0 ← satAdd16( GPR[rs]31..16 , GPR[rt]31..16 )
  tempA15..0 ← satAdd16( GPR[rs]15..0 , GPR[rt]15..0 )
  GPR[rd]31..0 ← tempB15..0 || tempA15..0

function add16( a15..0, b15..0 )
  temp16..0 ← ( a15 || a15..0 ) + ( b15 || b15..0 )
  if ( temp16 ≠ temp15 ) then
    DSPControlouflag:20 ← 1
  endif

```

ADDQ[_S].PH**Add Fractional Halfword Vectors**

```

    return temp15..0
endfunction add16

function satAdd16( a15..0, b15..0 )
    temp16..0 ← ( a15 || a15..0 ) + ( b15 || b15..0 )
    if ( temp16 ≠ temp15 ) then
        if ( temp16 = 0 ) then
            temp15..0 ← 0x7FFF
        else
            temp15..0 ← 0x8000
        endif
        DSPControlouflag:20 ← 1
    endif
    return temp15..0
endfunction satAdd16

```

Exceptions:

Reserved Instruction, DSP Disabled

ADDQ_S.W**Add Fractional Words**

31	26	25	21	20	16	15	11	10	9	3	2	0
P32A 001000	rt		rs		rd		x	1100000			101	
6	5		5		5		1	7			3	

Format: ADDQ_S.W rd, rs, rt**DSP****Purpose:** Add Fractional Words

Addition of two Q31 fractional values to produce a Q31 fractional result, with saturation.

Description: $rd \leftarrow \text{sat32}(rs_{31..0} + rt_{31..0})$

The Q31 fractional word in register *rt* is added to the corresponding fractional word in register *rs*. The result is then written to the destination register *rd*.

Signed saturating arithmetic is used, where an overflow is clamped to the largest representable value (0x7FFFFFFF hexadecimal) and an underflow to the smallest representable value (0x80000000 hexadecimal) before being sign-extended and written to the destination register *rd*.

If the addition results in underflow, overflow, or saturation, a 1 is written to bit 20 in the *DSPControl* register within the *ouflag* field.

Restrictions:

No data-dependent exceptions are possible.

The operands must be values in the specified format. If they are not, the results are **UNPREDICTABLE** and the values of the operand vectors become **UNPREDICTABLE**.

Operation:

```

ValidateAccessToDSPResources()
temp31..0 ← satAdd32( GPR[rs]31..0 , GPR[rt]31..0 )
GPR[rd]31..0 ← temp31..0

function satAdd32( a31..0, b31..0 )
    temp32..0 ← ( a31 || a31..0 ) + ( b31 || b31..0 )
    if ( temp32 ≠ temp31 ) then
        if ( temp32 = 0 ) then
            temp31..0 ← 0x7FFFFFFF
        else
            temp31..0 ← 0x80000000
        endif
        DSPControlouflag:20 ← 1
    endif
    return temp31..0
endfunction satAdd32

```

Exceptions:

Reserved Instruction, DSP Disabled

ADDQ_S.W

Add Fractional Words

ADDQH[_R].PH**Add Fractional Halfword Vectors And Shift Right to Halve Results**

31	26	25	21	20	16	15	11	10	9	3	2	0
ADDQH.PH												
P32A 001000		rt		rs		rd		0	0001001		101	
ADDQH_R.PH												
P32A 001000		rt		rs		rd		1	0001001		101	
6		5		5		5		1	7		3	

Format: ADDQH[_R].PH

ADDQH.PH rd, rs, rt

DSP-R2

ADDQH_R.PH rd, rs, rt

DSP-R2**Purpose:** Add Fractional Halfword Vectors And Shift Right to Halve Results

Element-wise fractional addition of halfword vectors, with a right shift by one bit to halve each result, with optional rounding.

Description: $rd \leftarrow \text{round}((rs_{31..16} + rt_{31..16}) \gg 1) \parallel \text{round}((rs_{15..0} + rt_{15..0}) \gg 1)$

Each element from the two halfword values in register *rs* is added to the corresponding halfword element in register *rt* to create an interim 17-bit result.

In the non-rounding instruction variant, each interim result is then shifted right by one bit before being written to the corresponding halfword element of destination register *rd*.

In the rounding version of the instruction, a value of 1 is added at the least-significant bit position of each interim result; the interim result is then right-shifted by one bit and written to the destination register.

This instruction does not modify the *DSPControl* register.

Restrictions:

No data-dependent exceptions are possible.

The operands must be a value in the specified format. If they are not, the results are **UNPREDICTABLE** and the values of the operand vectors become **UNPREDICTABLE**.

Operation:

```

ADDQH.PH
    ValidateAccessToDSP2Resources()
    tempB15..0 ← rightShift1AddQ16( GPR[rs]31..16 , GPR[rt]31..16 )
    tempA15..0 ← rightShift1AddQ16( GPR[rs]15..0 , GPR[rt]15..0 )
    GPR[rd]31..0 ← tempB15..0 || tempA15..0

ADDQH_R.PH
    ValidateAccessToDSP2Resources()
    tempB15..0 ← roundRightShift1AddQ16( GPR[rs]31..16 , GPR[rt]31..16 )
    tempA15..0 ← roundRightShift1AddQ16( GPR[rs]15..0 , GPR[rt]15..0 )
    GPR[rd]31..0 ← tempB15..0 || tempA15..0

function rightShift1AddQ16( a15..0 , b15..0 )
    temp16..0 ← (( a15 || a15..0 ) + ( b15 || b15..0 ))
    return temp16..1
endfunction rightShift1AddQ16

function roundRightShift1AddQ16( a15..0 , b15..0 )
    temp16..0 ← (( a15 || a15..0 ) + ( b15 || b15..0 ))
    temp16..0 ← temp16..0 + 1

```

ADDQH[_R].PH

Add Fractional Halfword Vectors And Shift Right to Halve Results

```
    return temp16..1  
endfunction roundRightShift1AddQ16
```

Exceptions:

Reserved Instruction, DSP Disabled

ADDQH[_R].W**Add Fractional Words And Shift Right to Halve Results**

31	26	25	21	20	16	15	11	10	9	3	2	0
ADDQH.W												
P32A 001000	rt		rs		rd		0	0010001			101	
ADDQH_R.W												
P32A 001000	rt		rs		rd		1	0010001			101	
6	5		5		5		1	7			3	

Format: ADDQH[_R].W

ADDQH.W rd, rs, rt

DSP-R2

ADDQH_R.W rd, rs, rt

DSP-R2**Purpose:** Add Fractional Words And Shift Right to Halve Results

Fractional addition of word vectors, with a right shift by one bit to halve the result, with optional rounding.

Description: $rd \leftarrow \text{round}((rs_{31..0} + rt_{31..0}) \gg 1)$ The word in register *rs* is added to the word in register *rt* to create an interim 33-bit result.In the non-rounding instruction variant, the interim result is then shifted right by one bit before being written to the destination register *rd*.

In the rounding version of the instruction, a value of 1 is added at the least-significant bit position of the interim result; the interim result is then right-shifted by one bit and written to the destination register.

This instruction does not modify the *DSPControl* register.**Restrictions:**

No data-dependent exceptions are possible.

The operands must be a value in the specified format. If they are not, the results are **UNPREDICTABLE** and the values of the operand vectors become **UNPREDICTABLE**.**Operation:**

ADDQH.W

ValidateAccessToDSP2Resources()

tempA_{31..0} ← rightShift1AddQ32(GPR[rs]_{31..0} , GPR[rt]_{31..0})GPR[rd]_{31..0} ← tempA_{31..0}

ADDQH_R.W

ValidateAccessToDSP2Resources()

tempA_{31..0} ← roundRightShift1AddQ32(GPR[rs]_{31..0} , GPR[rt]_{31..0})GPR[rd]_{31..0} ← tempA_{31..0}function rightShift1AddQ32(a_{31..0} , b_{31..0})temp_{32..0} ← ((a₃₁ || a_{31..0}) + (b₃₁ || b_{31..0}))return temp_{32..1}

endfunction rightShift1AddQ32

function roundRightShift1AddQ32(a_{31..0} , b_{31..0})temp_{32..0} ← ((a₃₁ || a_{31..0}) + (b₃₁ || b_{31..0}))temp_{32..0} ← temp_{32..0} + 1return temp_{32..1}

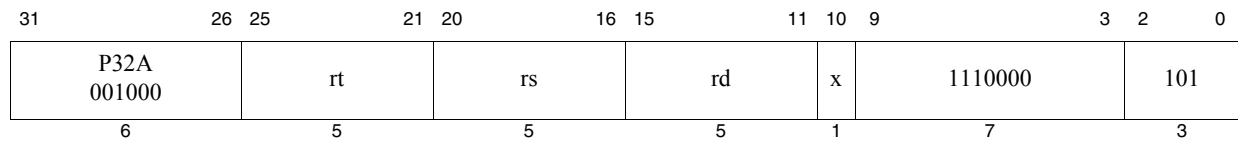
endfunction roundRightShift1AddQ32

ADDQH[_R].W

Add Fractional Words And Shift Right to Halve Results

Exceptions:

Reserved Instruction, DSP Disabled

ADDSC**Add Signed Word and Set Carry Bit****Format:** ADDSC rd, rs, rt**DSP****Purpose:** Add Signed Word and Set Carry Bit

Add two signed 32-bit values and set the carry bit in the *DSPControl* register if the addition generates a carry-out bit.

Description: $\text{DSPControl}[c], rd \leftarrow rs + rt$

The 32-bit signed value in register *rt* is added to the 32-bit signed value in register *rs*. The result is then written into register *rd*. The carry bit result out of the addition operation is written to bit 13 (the *c* field) of the *DSPControl* register.

This instruction does not modify the *ouflag* field in the *DSPControl* register.

Restrictions:

No data-dependent exceptions are possible.

Operation:

```

ValidateAccessToDSPResources()
temp32..0 ← ( 0 || GPR[rs]31..0 ) + ( 0 || GPR[rt]31..0 )
DSPControlc:13 ← temp32
GPR[rd]31..0 ← temp31..0

```

Exceptions:

Reserved Instruction, DSP Disabled

Programming Notes:

Note that this is really two's complement (modulo) arithmetic on the two integer values, where the overflow is preserved in architectural state. The ADDWC instruction can be used to do an add using this carry bit. These instructions are provided in the MIPS32 ISA to support 64-bit addition and subtraction using two pairs of 32-bit GPRs to hold each 64-bit value. In the MIPS64 ISA, 64-bit addition and subtraction can be performed directly, without requiring the use of these instructions.

ADDSC

Add Signed Word and Set Carry Bit

ADDU[_S].PH**Unsigned Add Integer Halfwords**

31	26	25	21	20	16	15	11	10	9	3	2	0
ADDU.PH												
P32A 001000		rt		rs		rd		0	0100001		101	
ADDU_S.PH												
P32A 001000		rt		rs		rd		1	0100001		101	
6		5		5		5		1	7		3	

Format: ADDU[_S].PH
 ADDU.PH rd, rs, rt
 ADDU_S.PH rd, rs, rt

DSP-R2**DSP-R2**

Purpose: Unsigned Add Integer Halfwords

Add two pairs of unsigned integer halfwords, with optional saturation.

Description: $rd \leftarrow \text{sat16}(rs_{31..16} + rt_{31..16}) \parallel \text{sat16}(rs_{15..0} + rt_{15..0})$

The two unsigned integer halfword elements in register *rt* are added to the corresponding unsigned integer halfword elements in register *rs*.

For the non-saturating version of the instruction, the result modulo 65,536 is written into the corresponding element in register *rd*.

For the saturating version of the instruction, the addition is performed using unsigned saturating arithmetic. Results that overflow are clamped to the largest representable value (65,535 decimal, 0xFFFF hexadecimal) before being written to the destination register *rd*.

For either instruction, if any of the individual additions result in overflow or saturation, a 1 is written to bit 20 in the *DSPControl* register within the ouflag field.

Restrictions:

No data-dependent exceptions are possible.

The operands must be a value in the specified format. If they are not, the results are **UNPREDICTABLE** and the values of the operand vectors become **UNPREDICTABLE**.

Operation:

```

ADDU.PH
  ValidateAccessToDSP2Resources()
  tempB15..0 ← addU16( GPR[rs]31..16 , GPR[rt]31..16 )
  tempA15..0 ← addU16( GPR[rs]15..0 , GPR[rt]15..0 )
  GPR[rd]31..0 ← tempB15..0 || tempA15..0

ADDU_S.PH
  ValidateAccessToDSP2Resources()
  tempB15..0 ← satAddU16( GPR[rs]31..16 , GPR[rt]31..16 )
  tempA15..0 ← satAddU16( GPR[rs]15..0 , GPR[rt]15..0 )
  GPR[rd]31..0 ← tempB15..0 || tempA15..0

```

Exceptions:

Reserved Instruction, DSP Disabled

ADDU[_S].PH

Unsigned Add Integer Halfwords

ADDU[_S].QB**Unsigned Add Quad Byte Vectors**

31	26	25	21	20	16	15	11	10	9	3	2	0
ADDU.QB												
P32A 001000		rt		rs		rd		0	0011001		101	
ADDU_S.QB												
P32A 001000		rt		rs		rd		1	0011001		101	
6		5		5		5		1	7		3	

Format: ADDU[_S].QB
 ADDU.QB rd, rs, rt
 ADDU_S.QB rd, rs, rt

DSP
DSP

Purpose: Unsigned Add Quad Byte Vectors

Element-wise addition of two vectors of unsigned byte values to produce a vector of unsigned byte results, with optional saturation.

Description: $rd \leftarrow \text{sat8}(rs_{31..24} + rt_{31..24}) \parallel \text{sat8}(rs_{23..16} + rt_{23..16}) \parallel \text{sat8}(rs_{15..8} + rt_{15..8}) \parallel \text{sat8}(rs_{7..0} + rt_{7..0})$

The four byte elements in register *rt* are added to the corresponding byte elements in register *rs*.

For the non-saturating version of the instruction, the result modulo 256 is written into the corresponding element in register *rd*.

For the saturating version of the instruction, the addition is performed using unsigned saturating arithmetic. Results that overflow are clamped to the largest representable value (255 decimal, 0xFF hexadecimal) before being written to the destination register *rd*.

For either instruction, if any of the individual additions result in overflow or saturation, a 1 is written to bit 20 in the *DSPControl* register within the *ouflag* field.

Restrictions:

No data-dependent exceptions are possible.

The operands must be values in the specified format. If they are not, the results are **UNPREDICTABLE** and the values of the operand vectors become **UNPREDICTABLE**.

Operation:

```

ADDU.QB:
    ValidateAccessToDSPResources()
    tempD7..0 ← addU8( GPR[rs]31..24 , GPR[rt]31..24 )
    tempC7..0 ← addU8( GPR[rs]23..16 , GPR[rt]23..16 )
    tempB7..0 ← addU8( GPR[rs]15..8 , GPR[rt]15..8 )
    tempA7..0 ← addU8( GPR[rs]7..0 , GPR[rt]7..0 )
    GPR[rd]31..0 ← tempD7..0 || tempC7..0 || tempB7..0 || tempA7..0

ADDU_S.QB:
    ValidateAccessToDSPResources()
    tempD7..0 ← satAddU8( GPR[rs]31..24 , GPR[rt]31..24 )
    tempC7..0 ← satAddU8( GPR[rs]23..16 , GPR[rt]23..16 )
    tempB7..0 ← satAddU8( GPR[rs]15..8 , GPR[rt]15..8 )
    tempA7..0 ← satAddU8( GPR[rs]7..0 , GPR[rt]7..0 )
    GPR[rd]31..0 ← tempD7..0 || tempC7..0 || tempB7..0 || tempA7..0

function addU8( a7..0, b7..0 )
    temp8..0 ← ( 0 || a7..0 ) + ( 0 || b7..0 )

```

ADDU[_S].QB**Unsigned Add Quad Byte Vectors**

```

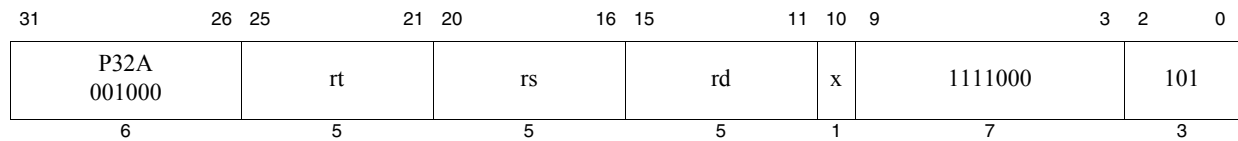
    if ( temp8 = 1 ) then
        DSPControlouflag:20 ← 1
    endif
    return temp7..0
endfunction addU8

function satAddU8( a7..0, b7..0 )
    temp8..0 ← ( 0 || a7..0 ) + ( 0 || b7..0 )
    if ( temp8 = 1 ) then
        temp7..0 ← 0xFF
        DSPControlouflag:20 ← 1
    endif
    return temp7..0
endfunction satAddU8

```

Exceptions:

Reserved Instruction, DSP Disabled

ADDWC**Add Word with Carry Bit****Format:** ADDWC rd, rs, rt**DSP****Purpose:** Add Word with Carry BitAdd two signed 32-bit values with the carry bit in the *DSPControl* register.**Description:** $rd \leftarrow rs + rt + DSPControl_{c:13}$

The 32-bit value in register *rt* is added to the 32-bit value in register *rs* and the carry bit in the *DSPControl* register. The result is then written to destination register *rd*.

If the addition results in either overflow or underflow, this instruction writes a 1 to bit 20 in the *ouflag* field of the *DSPControl* register.

Restrictions:

No data-dependent exceptions are possible.

Operation:

```

ValidateAccessToDSPResources()
temp32..0 ← ( GPR[rs]31 || GPR[rs]31..0 ) + ( GPR[rt]31 || GPR[rt]31..0 ) + ( 032 ||
DSPControlc:13 )
if ( temp32 ≠ temp31 ) then
    DSPControlouflag:20 ← 1
endif
GPR[rd]31..0 ← temp31..0

```

Exceptions:

Reserved Instruction, DSP Disabled

ADDWC

Add Word with Carry Bit

ADDUH[_R].QB**Unsigned Add Vector Quad-Bytes And Right Shift to Halve Results**

31	26	25	21	20	16	15	11	10	9	3	2	0
ADDUH.QB												
P32A 001000		rt		rs		rd		0	0101001		101	
ADDUH_R.QB												
P32A 001000		rt		rs		rd		1	0101001		101	
6		5		5		5		1	7		3	

Format: ADDUH[_R].QB

ADDUH.QB rd, rs, rt

DSP-R2

ADDUH_R.QB rd, rs, rt

DSP-R2**Purpose:** Unsigned Add Vector Quad-Bytes And Right Shift to Halve Results

Element-wise unsigned addition of unsigned byte vectors, with right shift by one bit to halve each result, with optional rounding.

Description $rd \leftarrow \text{round}((rs_{31..24} + rt_{31..24}) \gg 1) \parallel \text{round}((rs_{23..16} + rt_{23..16}) \gg 1) \parallel \text{round}((rs_{15..8} + rt_{15..8}) \gg 1) \parallel \text{round}((rs_{7..0} + rt_{7..0}) \gg 1)$

Each element from the four unsigned byte values in register *rs* is added to the corresponding unsigned byte element in register *rt* to create an unsigned interim result.

In the non-rounding instruction variant, each interim result is then shifted right by one bit before being written to the corresponding unsigned byte element of destination register *rd*.

In the rounding version of the instruction, a value of 1 is added at the least-significant bit position of each interim result before being right-shifted by one bit and written to the destination register.

This instruction does not modify the *DSPControl* register.

Restrictions:

No data-dependent exceptions are possible.

The operands must be a value in the specified format. If they are not, the results are **UNPREDICTABLE** and the values of the operand vectors become **UNPREDICTABLE**.

Operation:

```

ADDUH.QB
    ValidateAccessToDSPResources()
    tempD7..0 ← rightShift1AddU8( GPR[rs]31..24 , GPR[rt]31..24 )
    tempC7..0 ← rightShift1AddU8( GPR[rs]23..16 , GPR[rt]23..16 )
    tempB7..0 ← rightShift1AddU8( GPR[rs]15..8 , GPR[rt]15..8 )
    tempA7..0 ← rightShift1AddU8( GPR[rs]7..0 , GPR[rt]7..0 )
    GPR[rd]31..0 ← tempD7..0 || tempC7..0 || tempB7..0 || tempA7..0

ADDUH_R.QB
    ValidateAccessToDSPResources()
    tempD7..0 ← roundRightShift1AddU8( GPR[rs]31..24 , GPR[rt]31..24 )
    tempC7..0 ← roundRightShift1AddU8( GPR[rs]23..16 , GPR[rt]23..16 )
    tempB7..0 ← roundRightShift1AddU8( GPR[rs]15..8 , GPR[rt]15..8 )
    tempA7..0 ← roundRightShift1AddU8( GPR[rs]7..0 , GPR[rt]7..0 )
    GPR[rd]31..0 ← tempD7..0 || tempC7..0 || tempB7..0 || tempA7..0

function rightShift1AddU8( a7..0 , b7..0 )
    temp8..0 ← (( 0 || a7..0 ) + ( 0 || b7..0 ))

```

ADDUH[_R].QB**Unsigned Add Vector Quad-Bytes And Right Shift to Halve Results**

```

    return temp8..1
endfunction rightShift1AddU8

function roundRightShift1AddU8( a7..0 , b7..0 )
    temp8..0 ← (( 0 || a7..0 ) + ( 0 || b7..0 ))
    temp8..0 ← temp8..0 + 1
    return temp8..1
endfunction roundRightShift1AddU8

```

Exceptions:

Reserved Instruction, DSP Disabled

BALIGN**Byte Align Contents from Two Registers**

Format: BALIGN *rt*, *rs*, *bp*
 EXTW *rt*, *rs*, *rt*, $8 \cdot (4 - bp)$

DSP-R2
 Replaced with EXTW in nanoMIPS

Purpose: Byte Align Contents from Two Registers

Create a word result by combining a specified number of bytes from each of two source registers.

Description: $rt \leftarrow (rt \ll 8 \cdot bp) \mid (rs \gg 8 \cdot (4 - bp))$

The 32-bit word in register *rt* is left-shifted as a 32-bit value by *bp* byte positions, and the right-most word in register *rs* is right-shifted as a 32-bit value by $(4 - bp)$ byte positions. The shifted values are then *or*-ed together to create a 32-bit result that is written to destination register *rt*.

The argument *bp* is provided by the instruction, and is interpreted as an unsigned two-bit integer taking values between zero and three.

Restrictions:

No data-dependent exceptions are possible.

Operation:

```

ValidateAccessToDSP2Resources()
if (bp1..0 = 0) or (bp1..0 = 2) then
    GPR[rt]31..0 ← UNPREDICTABLE
else
    temp31..0 ← ( GPR[rt]31..0 << (8*bp1..0) ) || ( GPR[rs]31..0 >> (8*(4-bp1..0)) )
    GPR[rt]31..0 = temp31..0
endif

```

Implementation Notes:

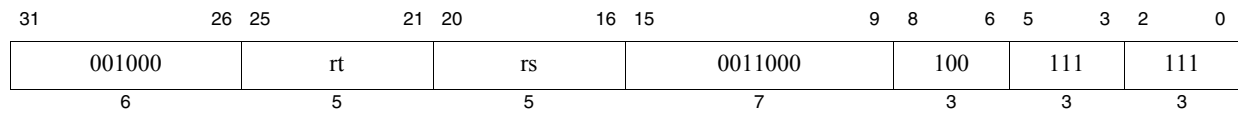
When *bp* is equal to zero, no left-shift is performed. When *bp* is equal to two, the result is equivalent to a PACKRL operation when the destination register is identical to the first source register. The assembler is expected to map these two variants of the BALIGN instructions to the appropriate equivalents. The only valid values of *bp* that the hardware must implement are when *bp* is equal to 1 and 3. If this instruction is passed through to the hardware with *bp* value equal to 0 or 2, the result is **UNPREDICTABLE**.

Exceptions:

Reserved Instruction, DSP Disabled

BALIGN

Byte Align Contents from Two Registers

BITREV**Bit-Reverse Halfword****Format:** BITREV *rt*, *rs***DSP****Purpose:** Bit-Reverse Halfword

To reverse the order of the bits of the least-significant halfword in the specified register.

Description: $rt \leftarrow rs_{0..15}$ The right-most halfword value in register *rs* is bit-reversed into the right-most halfword position in the destination register *rt*. The 16 most-significant bits of the destination register are zero-filled.**Restrictions:**

No data-dependent exceptions are possible.

The operands must be values in the specified format. If they are not, the results are **UNPREDICTABLE** and the values of the operand vectors become **UNPREDICTABLE**.**Operation:**

```

ValidateAccessToDSPResources()
temp15..0 ← GPR[rs]0..15
GPR[rt]31..0 ← 016 || temp15..0

```

Exceptions:

Reserved Instruction, DSP Disabled

BITREV

Bit-Reverse Halfword

BPOSGE32C**Branch on Greater Than or Equal To Value 32 in DSPControl Pos Field**

31	26	25	21	20	16	15	14	13	1	0
P32A 100010	x		00100		01		s[13:1]			s [14]
6	5		5		2		13			1

Format: BPOSGE32C offset**DSP-R3****Purpose:** Branch on Greater Than or Equal To Value 32 in *DSPControl* Pos FieldPerform a PC-relative branch if the value of the pos field in the *DSPControl* register is greater than or equal to 32.**Description:** if (DSPControl_{pos:5..0} >= 32) then goto PC+offset

First, the *offset* argument is left-shifted by one bit to form a 17-bit signed integer value. This value is added to the address of the instruction immediately following the branch to form a target branch address. Then, if the value of the pos field of the *DSPControl* register is greater than or equal to 32, the branch is taken and execution begins from the target address.

Restrictions:

Any instruction may be placed at PC + 4, where PC is that of the branch. An exception on such an instruction does not affect CP0 CAUSE_{BD}, and CP0 EPC is that of instruction in slot after branch.

Availability:

This instruction is introduced by and required as of Revision 3 of the DSP Module.

Operation:

```

I:      ValidateAccessToDSPResources()
          se_offsetGPRLEN..0 ← ( offset15 )GPRLEN-17 || offset15..0 || 01
          branch_condition ← ( DSPControlpos:5..0 >= 32 ? 1 : 0 )
I+1:   if ( branch_condition = 1 ) then
            PCGPRLEN..0 ← PCGPRLEN..0 + se_offsetGPRLEN..0
          endif

```

Exceptions:

Reserved Instruction, DSP Disabled

Programming Notes:

With the 17-bit signed instruction offset, the conditional branch range is ±64 Kbytes. Use jump (J) or jump register (JR) instructions to branch to addresses outside of this range.

BPOSGE32C

Branch on Greater Than or Equal To Value 32 in DSPControl Pos Field

CMP.cond.PH**Compare Vectors of Signed Integer Halfword Values**

31	26	25	21	20	16	15	10	9	3	2	0
CMP.EQ.PH											
P32A 0010000	rt		rs		x		0000000			101	
CMP.LE.PH											
P32A 0010000	rt		rs		x		0010000			101	
CMP.LT.PH											
P32A 0010000	rt		rs		x		0001000			101	
6	5		5		6		7			3	

Format: CMP.cond.PH

CMP.EQ.PH rs, rt

CMP.LT.PH rs, rt

CMP.LE.PH rs, rt

DSP**DSP****DSP****Purpose:** Compare Vectors of Signed Integer Halfword Values

Perform an element-wise comparison of two vectors of two signed integer halfwords, recording the results of the comparison in condition code bits.

Description: $\text{DSPControl}_{\text{ccond}:25..24} \leftarrow (\text{rs}_{31..16} \text{ cond } \text{rt}_{31..16}) \mid \mid (\text{rs}_{15..0} \text{ cond } \text{rt}_{15..0})$

The two signed integer halfword elements in register *rs* are compared with the corresponding signed integer halfword element in register *rt*. The two 1-bit boolean comparison results are written to bits 24 and 25 of the *DSPControl* register's 4-bit condition code field. The values of the two remaining condition code bits (bits 26 through 27 of the *DSPControl* register) are **UNPREDICTABLE**.

Restrictions:

No data-dependent exceptions are possible.

The operands must be values in the specified format. If they are not, the results are **UNPREDICTABLE** and the values of the operand vectors become **UNPREDICTABLE**.

Operation:

```

CMP.EQ.PH
    ValidateAccessToDSPResources()
    ccB ← GPR[rs]31..16 EQ GPR[rt]31..16
    ccA ← GPR[rs]15..0 EQ GPR[rt]15..0
    DSPControlccond:25..24 ← ccB || ccA
    DSPControlccond:27..26 ← UNPREDICTABLE

CMP.LT.PH
    ValidateAccessToDSPResources()
    ccB ← GPR[rs]31..16 LT GPR[rt]31..16
    ccA ← GPR[rs]15..0 LT GPR[rt]15..0
    DSPControlccond:25..24 ← ccB || ccA
    DSPControlccond:27..26 ← UNPREDICTABLE

CMP.LE.PH
    ValidateAccessToDSPResources()
    ccB ← GPR[rs]31..16 LE GPR[rt]31..16
    ccA ← GPR[rs]15..0 LE GPR[rt]15..0
    DSPControlccond:25..24 ← ccB || ccA
    DSPControlccond:27..26 ← UNPREDICTABLE

```

CMP.cond.PH

Compare Vectors of Signed Integer Halfword Values

Exceptions:

Reserved Instruction, DSP Disabled

CMPGDU.cond.QB Compare Unsigned Vector of Four Bytes and Write Result to GPR and DSPControl

31	26	25	21	20	16	15	11	10	9	3	2	0
CMPGDU.EQ.QB												
P32A 001000	rt		rs		rd		x	0110000			101	
CMPGDU.LE.QB												
P32A 001000	rt		rs		rd		x	1000000			101	
CMPGDU.LT.QB												
P32A 001000	rt		rs		rd		x	0111000			101	
6	5		5		5		1	7			3	

Format: CMPGDU.cond.QB

CMPGDU.EQ.QB rd, rs, rt

DSP-R2

CMPGDU.LT.QB rd, rs, rt

DSP-R2

CMPGDU.LE.QB rd, rs, rt

DSP-R2**Purpose:** Compare Unsigned Vector of Four Bytes and Write Result to GPR and DSPControl

Compare two vectors of four unsigned bytes each, recording the comparison results in condition code bits that are written to both the specified destination GPR and the condition code bits in the DSPControl register.

Description: $\text{DSPControl}[\text{ccond}]_{27..24} \leftarrow (\text{rs}_{31..24} \text{ cond } \text{rt}_{31..24}) \mid (\text{rs}_{23..16} \text{ cond } \text{rt}_{23..16}) \mid (\text{rs}_{15..8} \text{ cond } \text{rt}_{15..8}) \mid (\text{rs}_{7..0} \text{ cond } \text{rt}_{7..0});$
 $\text{rd} \leftarrow 0^{(\text{GPRLEN}-4)} \mid \text{DSPControl}[\text{ccond}]_{27..24}$

Each of the unsigned byte elements in register *rs* are compared with the corresponding unsigned byte elements in register *rt*. The four 1-bit boolean comparison results are written to the four least-significant bits of destination register *rd* and to bits 24 through 27 of the *DSPControl* register's 4-bit condition code field. The remaining bits in destination register *rd* are set to zero.

Restrictions:

No data-dependent exceptions are possible.

The operands must be a value in the specified format. If they are not, the results are **UNPREDICTABLE** and the values of the operand vectors become **UNPREDICTABLE**.

Operation:

CMPGDU.EQ.QB

```

ValidateAccessToDSP2Resources()
ccD ← GPR[rs]31..24 EQ GPR[rt]31..24
ccC ← GPR[rs]23..16 EQ GPR[rt]23..16
ccB ← GPR[rs]15..8 EQ GPR[rt]15..8
ccA ← GPR[rs]7..0 EQ GPR[rt]7..0
DSPControlcc:27..24 ← ccD || ccC || ccB || ccA
GPR[rd]31..0 ← 0(GPRLEN-4) || ccD || ccC || ccB || ccA

```

CMPGDU.LT.QB

```

ValidateAccessToDSP2Resources()
ccD ← GPR[rs]31..24 LT GPR[rt]31..24
ccC ← GPR[rs]23..16 LT GPR[rt]23..16
ccB ← GPR[rs]15..8 LT GPR[rt]15..8
ccA ← GPR[rs]7..0 LT GPR[rt]7..0
DSPControlcc:27..24 ← ccD || ccC || ccB || ccA
GPR[rd]31..0 ← 0(GPRLEN-4) || ccD || ccC || ccB || ccA

```

CMPGDU.cond.QB Compare Unsigned Vector of Four Bytes and Write Result to GPR and DSPControl

```

CMPGDU.LE.QB
    ValidateAccessToDSP2Resources()
    ccD ← GPR[rs]31..24 LE GPR[rt]31..24
    ccC ← GPR[rs]23..16 LE GPR[rt]23..16
    ccB ← GPR[rs]15..8 LE GPR[rt]15..8
    ccA ← GPR[rs]7..0 LE GPR[rt]7..0
    DSPControlcc:27..24 ← ccD || ccC || ccB || ccA
    GPR[rd]31..0 ← 0(GPRLEN-4) || ccD || ccC || ccB || ccA

```

Exceptions:

Reserved Instruction, DSP Disabled

CMPGU.cond.QB**Compare Vectors of Unsigned Byte Values and Write Results to a GPR**

31	26	25	21	20	16	15	11	10	9	3	2	0
CMPGU.EQ.QB												
P32A 001000	rt		rs		rd		x	0011000			101	
CMPGU.LE.QB												
P32A 001000	rt		rs		rd		x	0101000			101	
CMPGU.LT.QB												
P32A 001000	rt		rs		rd		x	0100000			101	
6	5		5		5		1	7			3	

Format: CMPGU.cond.QB

CMPGU.EQ.QB rd, rs, rt

CMPGU.LT.QB rd, rs, rt

CMPGU.LE.QB rd, rs, rt

DSP**DSP****DSP****Purpose:** Compare Vectors of Unsigned Byte Values and Write Results to a GPR

Perform an element-wise comparison of two vectors of unsigned bytes, recording the results of the comparison in condition code bits that are written to the specified GPR.

Description: $rd \leftarrow (rs_{31..24} \text{ cond } rt_{31..24}) \mid (rs_{23..16} \text{ cond } rt_{23..16}) \mid (rs_{15..8} \text{ cond } rt_{15..8}) \mid (rs_{7..0} \text{ cond } rt_{7..0})$

Each of the unsigned byte elements in register *rs* are compared with the corresponding unsigned byte elements in register *rt*. The four 1-bit boolean comparison results are written to the four least-significant bits of destination register *rd*. The remaining bits in *rd* are set to zero.

Restrictions:

No data-dependent exceptions are possible.

The operands must be values in the specified format. If they are not, the results are **UNPREDICTABLE** and the values of the operand vectors become **UNPREDICTABLE**.

Operation:

CMPGU.EQ.QB

```

ValidateAccessToDSPResources()
ccD ← GPR[rs]31..24 EQ GPR[rt]31..24
ccC ← GPR[rs]23..16 EQ GPR[rt]23..16
ccB ← GPR[rs]15..8 EQ GPR[rt]15..8
ccA ← GPR[rs]7..0 EQ GPR[rt]7..0
GPR[rd]31..0 ← 0(GPRLEN-4) || ccD || ccC || ccB || ccA

```

CMPGU.LT.QB

```

ValidateAccessToDSPResources()
ccD ← GPR[rs]31..24 LT GPR[rt]31..24
ccC ← GPR[rs]23..16 LT GPR[rt]23..16
ccB ← GPR[rs]15..8 LT GPR[rt]15..8
ccA ← GPR[rs]7..0 LT GPR[rt]7..0
GPR[rd]31..0 ← 0(GPRLEN-4) || ccD || ccC || ccB || ccA

```

CMPGU.LE.QB

```

ValidateAccessToDSPResources()
ccD ← GPR[rs]31..24 LE GPR[rt]31..24
ccC ← GPR[rs]23..16 LE GPR[rt]23..16

```

CMPGU.cond.QB**Compare Vectors of Unsigned Byte Values and Write Results to a GPR**

$$\begin{aligned}
 ccB &\leftarrow GPR[rs]_{15..8} \text{ LE } GPR[rt]_{15..8} \\
 ccA &\leftarrow GPR[rs]_{7..0} \text{ LE } GPR[rt]_{7..0} \\
 GPR[rd]_{31..0} &\leftarrow 0^{(GPRLEN-4)} \parallel ccD \parallel ccC \parallel ccB \parallel ccA
 \end{aligned}$$
Exceptions:

Reserved Instruction, DSP Disabled

CMPU.cond.QB**Compare Vectors of Unsigned Byte Values**

31	26	25	21	20	16	15	11	10	9	3	2	0
CMPU.EQ.QB												
P32A 001000	rt		rs		x		x	1001000			101	
CMPU.LE.QB												
P32A 001000	rt		rs		x		x	1011000			101	
CMPU.LT.QB												
P32A 001000	rt		rs		x		x	1010000			101	
6	5		5		5		1	7			3	

Format: CMPU.cond.QB

CMPU.EQ.QB rs, rt

CMPU.LT.QB rs, rt

CMPU.LE.QB rs, rt

DSP**DSP****DSP****Purpose:** Compare Vectors of Unsigned Byte Values

Perform an element-wise comparison of two vectors of unsigned bytes, recording the results of the comparison in condition code bits.

Description: $\text{DSPControl}_{\text{ccond}:27..24} \leftarrow (\text{rs}_{31..24} \text{ cond } \text{rt}_{31..24}) \mid (\text{rs}_{23..16} \text{ cond } \text{rt}_{23..16}) \mid (\text{rs}_{15..8} \text{ cond } \text{rt}_{15..8}) \mid (\text{rs}_{7..0} \text{ cond } \text{rt}_{7..0})$

Each of the unsigned byte elements in register *rs* are compared with the corresponding unsigned byte elements in register *rt*. The four 1-bit boolean comparison results are written to bits 24 through 27 of the *DSPControl* register's 4-bit condition code field.

Restrictions:

No data-dependent exceptions are possible.

The operands must be values in the specified format. If they are not, the results are **UNPREDICTABLE** and the values of the operand vectors become **UNPREDICTABLE**.

Operation:

CMPU.EQ.QB

```

ValidateAccessToDSPResources()
ccD ← GPR[rs]31..24 EQ GPR[rt]31..24
ccC ← GPR[rs]23..16 EQ GPR[rt]23..16
ccB ← GPR[rs]15..8 EQ GPR[rt]15..8
ccA ← GPR[rs]7..0 EQ GPR[rt]7..0
DSPControlccond:27..24 ← ccD || ccC || ccB || ccA

```

CMPU.LT.QB

```

ValidateAccessToDSPResources()
ccD ← GPR[rs]31..24 LT GPR[rt]31..24
ccC ← GPR[rs]23..16 LT GPR[rt]23..16
ccB ← GPR[rs]15..8 LT GPR[rt]15..8
ccA ← GPR[rs]7..0 LT GPR[rt]7..0
DSPControlccond:27..24 ← ccD || ccC || ccB || ccA

```

CMPU.LE.QB

```

ValidateAccessToDSPResources()
ccD ← GPR[rs]31..24 LE GPR[rt]31..24
ccC ← GPR[rs]23..16 LE GPR[rt]23..16

```


CMPU.cond.QB**Compare Vectors of Unsigned Byte Values**

```

ccB ← GPR[rs]15..8 LE GPR[rt]15..8
ccA ← GPR[rs]7..0 LE GPR[rt]7..0
DSPControlccond:27..24 ← ccD || ccC || ccB || ccA

```

Exceptions:

Reserved Instruction, DSP Disabled

DPA.W.PH**Dot Product with Accumulate on Vector Integer Halfword Elements**

31	26	25	21	20	16	15	14	13	12	11	9	8	6	5	3	2	0
P32A 001000		rt			rs			ac	00	000	010	111	111				
6		5			5			2	2	3	3	3	3				

Format: DPA.W.PH *ac*, *rs*, *rt***DSP-R2****Purpose:** Dot Product with Accumulate on Vector Integer Halfword Elements

Generate the dot-product of two integer halfword vector elements using full-size intermediate products and then accumulate into the specified accumulator register.

Description: $ac \leftarrow ac + ((rs_{31..16} * rt_{31..16}) + (rs_{15..0} * rt_{15..0}))$

Each of the two halfword integer values from register *rt* is multiplied with the corresponding halfword element from register *rs* to create two integer word results. These two products are summed to generate a dot-product result, which is then accumulated into the specified 64-bit *HI/LO* accumulator, creating a 64-bit integer result.

The value of *ac* selects an accumulator numbered from 0 to 3. When *ac*=0, this refers to the original *HI/LO* register pair of the MIPS32 architecture.

This instruction does not set any bits of the *ouflag* field in the *DSPControl* register.

Restrictions:

No data-dependent exceptions are possible.

The operands must be a value in the specified format. If they are not, the result is **UNPREDICTABLE** and the values of the operand vectors become **UNPREDICTABLE**.

Operation:

```

ValidateAccessToDSP2Resources()
tempB31..0 ← (GPR[rs]31..16 * GPR[rt]31..16)
tempA31..0 ← (GPR[rs]15..0 * GPR[rt]15..0)
dotp32..0 ← (tempB31 || tempB31..0) + (tempA31 || tempA31..0)
acc63..0 ← (HI[ac]31..0 || LO[ac]31..0) + ((dotp32)31 || dotp32..0)
(HI[ac]31..0 || LO[ac]31..0) ← acc63..32 || acc31..0

```

Exceptions:

Reserved Instruction, DSP Disabled

DPA.W.PH

Dot Product with Accumulate on Vector Integer Halfword Elements

DPAQ_S.W.PH**Dot Product with Accumulation on Fractional Halfword Elements**

31	26	25	21	20	16	15	14	13	12	11	9	8	6	5	3	2	0
P32A 001000	rt		rs		ac	00	001	010	111	111							
6	5		5		2	2	3	3	3	3							

Format: DPAQ_S.W.PH ac, rs, rt**DSP****Purpose:** Dot Product with Accumulation on Fractional Halfword Elements

Element-wise multiplication of two vectors of fractional halfword elements and accumulation of the accumulated 32-bit intermediate products into the specified 64-bit accumulator register, with saturation.

Description: $ac \leftarrow ac + (\text{sat32}(rs_{31..16} * rt_{31..16}) + \text{sat32}(rs_{15..0} * rt_{15..0}))$

Each of the two Q15 fractional word values from registers *rt* and *rs* are multiplied together, and the results left-shifted by one bit position to generate two Q31 fractional format intermediate products. If both multiplicands for either of the multiplications are equal to -1.0 (0x8000 hexadecimal), the resulting intermediate product is saturated to the maximum positive Q31 fractional value (0x7FFFFFFF hexadecimal).

The two intermediate products are then sign-extended and summed to generate a 64-bit, Q32.31 fractional format dot-product result that is accumulated in to the specified 64-bit *HI/LO* accumulator to produce a final Q32.31 fractional result.

The value of *ac* can range from 0 to 3; a value of 0 refers to the original *HI/LO* register pair of the MIPS32 architecture.

If saturation occurs as a result of a halfword multiplication, a 1 is written to one of bits 16 through 19 of the *DSPControl* register, within the ouflag field. The value of *ac* determines which of these bits is set: bit 16 corresponds to *ac0*, bit 17 to *ac1*, bit 18 to *ac2*, and bit 19 to *ac3*.

Restrictions:

No data-dependent exceptions are possible.

The operands must be values in the specified format. If they are not, the result is **UNPREDICTABLE** and the values of the operand vectors become **UNPREDICTABLE**.

Operation:

```

ValidateAccessToDSPResources()
tempB31..0 ← multiplyQ15Q15( ac, GPR[rs]31..16, GPR[rt]31..16 )
tempA31..0 ← multiplyQ15Q15( ac, GPR[rs]15..0, GPR[rt]15..0 )
dotp63..0 ← ( (tempB31)32 || tempB31..0 ) + ( (tempA31)32 || tempA31..0 )
tempC63..0 ← ( HI[ac]31..0 || LO[ac]31..0 ) + dotp63..0
( HI[ac]31..0 || LO[ac]31..0 ) ← tempC63..32 || tempC31..0

function multiplyQ15Q15( acc1..0, a15..0, b15..0 )
  if ( a15..0 = 0x8000 ) and ( b15..0 = 0x8000 ) then
    temp31..0 ← 0x7FFFFFFF
    DSPControlouflag:16+acc ← 1
  else
    temp31..0 ← ( a15..0 * b15..0 ) << 1
  endif
  return temp31..0
endfunction multiplyQ15Q15

```

Exceptions:

Reserved Instruction, DSP Disabled

DPAQ_S.W.PH

Dot Product with Accumulation on Fractional Halfword Elements

DPAQ_SA.L.W**Dot Product with Accumulate on Fractional Word Element**

31	26	25	21	20	16	15	14	13	12	11	9	8	6	5	3	2	0
P32A 001000	rt				rs				ac	01	001	010	111	111			
6	5				5				2	2	3	3	3	3			

Format: DPAQ_SA.L.W ac, rs, rt**DSP****Purpose:** Dot Product with Accumulate on Fractional Word Element

Multiplication of two fractional word elements, accumulating the product to the specified 64-bit accumulator register, with saturation.

Description: $ac \leftarrow \text{sat}_{64}(ac + \text{sat}_{32}(rs_{31..0} * rt_{31..0}))$

The two right-most Q31 fractional word values from registers *rt* and *rs* are multiplied together and the result left-shifted by one bit position to generate a 64-bit, Q63 fractional format intermediate product. If both multiplicands are equal to -1.0 (0x80000000 hexadecimal), the intermediate product is saturated to the maximum positive Q63 fractional value (0x7FFFFFFFFFFFFFFF hexadecimal).

The intermediate product is then added to the specified 64-bit *HI/LO* accumulator, creating a Q63 fractional result. If the accumulation results in overflow or underflow, the accumulator is saturated to either the maximum positive or minimum negative Q63 fractional value (0x8000000000000000 hexadecimal), respectively.

The value of *ac* can range from 0 to 3; a value of 0 refers to the original *HI/LO* register pair of the MIPS32 architecture.

If saturation occurs, a 1 is written to one of bits 16 through 19 of the *DSPControl* register, within the *ouflag* field. The value of *ac* determines which of these bits is set: bit 16 corresponds to *ac0*, bit 17 to *ac1*, bit 18 to *ac2*, and bit 19 to *ac3*.

Restrictions:

No data-dependent exceptions are possible.

The operands must be values in the specified format. If they are not, the result is **UNPREDICTABLE** and the values of the operand vectors become **UNPREDICTABLE**.

Operation:

```

ValidateAccessToDSPResources()
dotp63..0 ← multiplyQ31Q31( ac, GPR[rs]31..0, GPR[rt]31..0 )
temp64..0 ← HI[ac]31 || HI[ac]31..0 || LO[ac]31..0
temp64..0 ← temp64..0 + dotp63..0
if ( temp64 ≠ temp63 ) then
    if ( temp64 = 1 ) then
        temp63..0 ← 0x8000000000000000
    else
        temp63..0 ← 0x7FFFFFFFFFFFFFFF
    endif
    DSPControlouflag:16+ac ← 1
endif
( HI[ac]31..0 || LO[ac]31..0 ) ← temp63..32 || temp31..0

function multiplyQ31Q31( acc1..0, a31..0, b31..0 )
    if ( ( a31..0 = 0x80000000 ) and ( b31..0 = 0x80000000 ) ) then
        temp63..0 ← 0x7FFFFFFFFFFFFFFF
        DSPControlouflag:16+acc ← 1
    else
        temp63..0 ← ( a31..0 * b31..0 ) << 1
    endif
endfunction

```

DPAQ_SA.L.W

Dot Product with Accumulate on Fractional Word Element

```
    return temp63..0  
endfunction multiplyQ31Q31
```

Exceptions:

Reserved Instruction, DSP Disabled

DPAQX_S.W.PH**Cross Dot Product with Accumulation on Fractional Halfword Elements**

31	26	25	21	20	16	15	14	13	12	11	9	8	6	5	3	2	0
P32A 001000	rt		rs		ac	10	001		010		111		111				
6	5		5		2	2	3		3		3		3		3		

Format: DPAQX_S.W.PH ac, rs, rt**DSP-R2****Purpose:** Cross Dot Product with Accumulation on Fractional Halfword Elements

Element-wise cross multiplication of two vectors of fractional halfword elements and accumulation of the 32-bit intermediate products into the specified 64-bit accumulator register, with saturation.

Description: $ac \leftarrow ac + (sat32(rs_{31..16} * rt_{15..0}) + sat32(rs_{15..0} * rt_{31..16}))$

The left Q15 fractional word value from registers *rt* is multiplied with the right halfword element from register *rs* and the result left-shifted by one bit position to generate a Q31 fractional format intermediate product. Similarly, the right Q15 fractional word value from registers *rt* is multiplied with the left halfword element from register *rs* and the result left-shifted by one bit position to generate a Q31 fractional format intermediate product. If both multiplicands for either of the multiplications are equal to -1.0 (0x8000 hexadecimal), the resulting intermediate product is saturated to the maximum positive Q31 fractional value (0x7FFFFFFF hexadecimal).

The two intermediate products are then sign-extended and summed to generate a 64-bit, Q32.31 fractional format dot-product result that is accumulated in to the specified 64-bit *H/L0* accumulator to produce a final Q32.31 fractional result.

The value of *ac* can range from 0 to 3; a value of 0 refers to the original *H/L0* register pair of the MIPS32 architecture.

If saturation occurs as a result of a halfword multiplication, a 1 is written to one of bits 16 through 19 of the *DSPControl* register, within the ouflag field. The value of *ac* determines which of these bits is set: bit 16 corresponds to *ac0*, bit 17 to *ac1*, bit 18 to *ac2*, and bit 19 to *ac3*.

Restrictions:

No data-dependent exceptions are possible.

The operands must be values in the specified format. If they are not, the result is **UNPREDICTABLE** and the values of the operand vectors become **UNPREDICTABLE**.

Operation:

```

ValidateAccessToDSPResources()
tempB31..0 ← multiplyQ15Q15( ac, GPR[rs]31..16, GPR[rt]15..0 )
tempA31..0 ← multiplyQ15Q15( ac, GPR[rs]15..0, GPR[rt]31..16 )
dotp63..0 ← ( (tempB31)32 || tempB31..0 ) + ( (tempA31)32 || tempA31..0 )
tempC63..0 ← ( HI[ac]31..0 || LO[ac]31..0 ) + dotp63..0
( HI[ac]31..0 || LO[ac]31..0 ) ← tempC63..32 || tempC31..0

function multiplyQ15Q15( acc1..0, a15..0, b15..0 )
  if ( a15..0 = 0x8000 ) and ( b15..0 = 0x8000 ) then
    temp31..0 ← 0x7FFFFFFF
    DSPControlouflag:16+acc ← 1
  else
    temp31..0 ← ( a15..0 * b15..0 ) << 1
  endif
  return temp31..0
endfunction multiplyQ15Q15

```

Exceptions:

Reserved Instruction, DSP Disabled

DPAQX_S.W.PH

Cross Dot Product with Accumulation on Fractional Halfword Elements

DPAQX_SA.W.PH**Cross Dot Product with Accumulation on Fractional Halfword Elements**

31	26	25	21	20	16	15	14	13	12	11	9	8	6	5	3	2	0
P32A 001000	rt		rs		ac	11	001		010		111		111				
6	5		5		2	2	3		3		3		3		3		

Format: DPAQX_SA.W.PH ac, rs, rt**DSP-R2****Purpose:** Cross Dot Product with Accumulation on Fractional Halfword Elements

Element-wise cross multiplication of two vectors of fractional halfword elements and accumulation of the 32-bit intermediate products into the specified 64-bit accumulator register, with saturation of the accumulator.

Description: $ac \leftarrow \text{sat32}(ac + (\text{sat32}(rs_{31..16} * rt_{15..0}) + \text{sat32}(rs_{15..0} * rt_{31..16})))$

The left Q15 fractional word value from registers *rt* is multiplied with the right halfword element from register *rs* and the result left-shifted by one bit position to generate a Q31 fractional format intermediate product. Similarly, the right Q15 fractional word value from registers *rt* is multiplied with the left halfword element from register *rs* and the result left-shifted by one bit position to generate a Q31 fractional format intermediate product. If both multiplicands for either of the multiplications are equal to -1.0 (0x8000 hexadecimal), the resulting intermediate product is saturated to the maximum positive Q31 fractional value (0x7FFFFFFF hexadecimal).

The two intermediate products are then sign-extended and summed to generate a 64-bit, Q32.31 fractional format dot-product result that is accumulated into the specified 64-bit *HI/LO* accumulator to produce a Q32.31 fractional result. If this result is larger than or equal to +1.0, or smaller than -1.0, it is saturated to the Q31 range.

The value of *ac* can range from 0 to 3; a value of 0 refers to the original *HI/LO* register pair of the MIPS32 architecture.

If saturation occurs as a result of halfword multiplication or accumulation, a 1 is written to one of bits 16 through 19 of the *DSPControl* register, within the ouflag field. The value of *ac* determines which of these bits is set: bit 16 corresponds to *ac0*, bit 17 to *ac1*, bit 18 to *ac2*, and bit 19 to *ac3*.

Restrictions:

No data-dependent exceptions are possible.

The operands must be values in the specified format. If they are not, the result is **UNPREDICTABLE** and the values of the operand vectors become **UNPREDICTABLE**.

Operation:

```

ValidateAccessToDSPResources()
tempB31..0 ← multiplyQ15Q15( ac, GPR[rs]31..16, GPR[rt]15..0 )
tempA31..0 ← multiplyQ15Q15( ac, GPR[rs]15..0, GPR[rt]31..16 )
dotp63..0 ← ( (tempB31)32 || tempB31..0 ) + ( (tempA31)32 || tempA31..0 )
tempC63..0 ← ( HI[ac]31..0 || LO[ac]31..0 ) + dotp63..0
if ( tempC63 = 0 ) and ( tempC62..31 ≠ 0 ) then
    tempC63..0 = 032 || 0x7FFFFFFF
    DSPControlouflag:16+acc ← 1
endif
if ( tempC63 = 1 ) and ( tempC62..31 ≠ 132 ) then
    tempC63..0 = 132 || 0x80000000
    DSPControlouflag:16+acc ← 1
endif
( HI[ac]31..0 || LO[ac]31..0 ) ← tempC63..32 || tempC31..0

function multiplyQ15Q15( acc15..0, a15..0, b15..0 )
    if ( a15..0 = 0x8000 ) and ( b15..0 = 0x8000 ) then
        temp31..0 ← 0x7FFFFFFF
        DSPControlouflag:16+acc ← 1
    
```

DPAQX_SA.W.PH**Cross Dot Product with Accumulation on Fractional Halfword Elements**

```
    else
        temp31..0 ← ( a15..0 * b15..0 ) << 1
    endif
    return temp31..0
endfunction multiplyQ15Q15
```

Exceptions:

Reserved Instruction, DSP Disabled

DPAU.H.QBL**Dot Product with Accumulate on Vector Unsigned Byte Elements**

31	26	25	21	20	16	15	14	13	12	11	9	8	6	5	3	2	0
P32A 001000						rt		rs		ac	10	000	010	111	111		
6						5		5		2	2	3	3	3	3		

Format: DPAU.H.QBL *ac*, *rs*, *rt***DSP****Purpose:** Dot Product with Accumulate on Vector Unsigned Byte Elements

Element-wise multiplication of the two left-most elements of the four elements of each of two vectors of unsigned bytes, accumulating the sum of the products into the specified 64-bit accumulator register.

Description: $ac \leftarrow ac + \text{zero_extend}((rs_{31..24} * rt_{31..24}) + (rs_{23..16} * rt_{23..16}))$

The two left-most elements of the four unsigned byte elements of each of registers *rt* and *rs* are multiplied together using unsigned arithmetic to generate two 16-bit unsigned intermediate products. The intermediate products are then zero-extended to 64 bits and accumulated into the specified 64-bit *HI/LO* accumulator.

The value of *ac* can range from 0 to 3; a value of 0 refers to the original *HI/LO* register pair of the MIPS32 architecture.

This instruction does not set any bits in the ouflag field in the *DSPControl* register.

Restrictions:

No data-dependent exceptions are possible.

The operands must be values in the specified format. If they are not, the result is **UNPREDICTABLE** and the values of the operand vectors become **UNPREDICTABLE**.

Operation:

```

ValidateAccessToDSPResources()
tempB15..0 ← multiplyU8U8( GPR[rs]31..24, GPR[rt]31..24 )
tempA15..0 ← multiplyU8U8( GPR[rs]23..16, GPR[rt]23..16 )
dotp63..0 ← ( 048 || tempB15..0 ) + ( 048 || tempA15..0 )
tempC63..0 ← ( HI[ac]31..0 || LO[ac]31..0 ) + dotp63..0
( HI[ac]31..0 || LO[ac]31..0 ) ← tempC63..32 || tempC31..0

function multiplyU8U8( a7..0, b7..0 )
    temp17..0 ← ( 0 || a7..0 ) * ( 0 || b7..0 )
    return temp15..0
endfunction multiplyU8U8

```

Exceptions:

Reserved Instruction, DSP Disabled

DPAU.H.QBL

Dot Product with Accumulate on Vector Unsigned Byte Elements

DPAU.H.QBR**Dot Product with Accumulate on Vector Unsigned Byte Elements**

31	26	25	21	20	16	15	14	13	12	11	9	8	6	5	3	2	0
P32A 001000	rt				rs				ac	11	000	010	111	111			
6	5				5				2	2	3	3	3	3	3		

Format: DPAU.H.QBR *ac*, *rs*, *rt***DSP****Purpose:** Dot Product with Accumulate on Vector Unsigned Byte Elements

Element-wise multiplication of the two right-most elements of the four elements of each of two vectors of unsigned bytes, accumulating the sum of the products into the specified 64-bit accumulator register.

Description: $ac \leftarrow ac + \text{zero_extend}((rs_{15..8} * rt_{15..8}) + (rs_{7..0} * rt_{7..0}))$

The two right-most elements of the four unsigned byte elements of each of registers *rt* and *rs* are multiplied together using unsigned arithmetic to generate two 16-bit unsigned intermediate products. The intermediate products are then zero-extended to 64 bits and accumulated into the specified 64-bit *HI/LO* accumulator.

The value of *ac* can range from 0 to 3; a value of 0 refers to the original *HI/LO* register pair of the MIPS32 architecture.

This instruction does not set any bits in the ouflag field in the *DSPControl* register.

Restrictions:

No data-dependent exceptions are possible.

The operands must be values in the specified format. If they are not, the result is **UNPREDICTABLE** and the values of the operand vectors become **UNPREDICTABLE**.

Operation:

```

ValidateAccessToDSPResources()
tempB15..0 ← multiplyU8U8( GPR[rs]15..8, GPR[rs]15..8 )
tempA15..0 ← multiplyU8U8( GPR[rs]7..0, GPR[rs]7..0 )
dotp63..0 ← ( 048 || tempB15..0 ) + ( 048 || tempA15..0 )
tempC63..0 ← ( HI[ac]31..0 || LO[ac]31..0 ) + dotp63..0
( HI[ac]31..0 || LO[ac]31..0 ) ← tempC63..32 || tempC31..0

```

Exceptions:

Reserved Instruction, DSP Disabled

DPAU.H.QBR

Dot Product with Accumulate on Vector Unsigned Byte Elements

DPAX.W.PH**Cross Dot Product with Accumulate on Vector Integer Halfword Elements**

31	26	25	21	20	16	15	14	13	12	11	9	8	6	5	3	2	0
P32A 001000		rt			rs			ac	01	000		010		111		111	
6		5			5			2	2	3		3		3		3	

Format: DPAX.W.PH *ac*, *rs*, *rt***DSP-R2****Purpose:** Cross Dot Product with Accumulate on Vector Integer Halfword Elements

Generate the cross dot-product of two integer halfword vector elements using full-size intermediate products and then accumulate into the specified accumulator register.

Description: $ac \leftarrow ac + ((rs_{31..16} * rt_{15..0}) + (rs_{15..0} * rt_{31..16}))$

The left halfword integer value from register *rt* is multiplied with the right halfword element from register *rs* to create an integer word result. Similarly, the right halfword integer value from register *rt* is multiplied with the left halfword element from register *rs* to create the second integer word result. These two products are summed to generate the dot-product result, which is then accumulated into the specified 64-bit *HI/LO* accumulator, creating a 64-bit integer result.

The value of *ac* selects an accumulator numbered from 0 to 3. When *ac*=0, this refers to the original *HI/LO* register pair of the MIPS32 architecture.

This instruction will not set any bits of the ouflag field in the *DSPControl* register.

Restrictions:

No data-dependent exceptions are possible.

The operands must be a value in the specified format. If they are not, the result is **UNPREDICTABLE** and the values of the operand vectors become **UNPREDICTABLE**.

Operation:

```

ValidateAccessToDSPResources()
tempB31..0 ← (GPR[rs]31..16 * GPR[rt]15..0)
tempA31..0 ← (GPR[rs]15..0 * GPR[rt]31..16)
dotp32..0 ← ( (tempB31) || tempB31..0 ) + ( (tempA31) || tempA31..0 )
acc63..0 ← ( HI[ac]31..0 || LO[ac]31..0 ) + ( (dotp32)31 || dotp32..0 )
( HI[ac]31..0 || LO[ac]31..0 ) ← acc63..32 acc31..0

```

Exceptions:

Reserved Instruction, DSP Disabled

DPAX.W.PH

Cross Dot Product with Accumulate on Vector Integer Halfword Elements

DPS.W.PH**Dot Product with Subtract on Vector Integer Half-Word Elements**

31	26	25	21	20	16	15	14	13	12	11	9	8	6	5	3	2	0
P32A 001000		rt			rs			ac	00	010	010	111	111				
6		5			5			2	2	3	3	3	3				

Format: DPS.W.PH *ac*, *rs*, *rt***DSP-R2****Purpose:** Dot Product with Subtract on Vector Integer Half-Word Elements

Generate the dot-product of two integer halfword vector elements using full-size intermediate products and then subtract from the specified accumulator register.

Description: $ac \leftarrow ac - ((rs_{31..16} * rt_{31..16}) + (rs_{15..0} * rt_{15..0}))$

Each of the two halfword integer values from register *rt* is multiplied with the corresponding halfword element from register *rs* to create two integer word results. These two products are summed to generate the dot-product result, which is then subtracted from the specified 64-bit *HI/LO* accumulator, creating a 64-bit integer result.

The value of *ac* selects an accumulator numbered from 0 to 3. When *ac*=0, this refers to the original *HI/LO* register pair of the MIPS32 architecture.

This instruction will not set any bits of the ouflag field in the *DSPControl* register.

Restrictions:

No data-dependent exceptions are possible.

The operands must be a value in the specified format. If they are not, the result is **UNPREDICTABLE** and the values of the operand vectors become **UNPREDICTABLE**.

Operation:

```

ValidateAccessToDSP2Resources()
tempB31..0 ← (GPR[rs]31..16 * GPR[rt]31..16)
tempA31..0 ← (GPR[rs]15..0 * GPR[rt]15..0)
dotp32..0 ← ( (tempB31) || tempB31..0 ) + ( (tempA31) || tempA31..0 )
acc63..0 ← ( HI[ac]31..0 || LO[ac]31..0 ) - ( (dotp32)31 || dotp32..0 )
( HI[ac]31..0 || LO[ac]31..0 ) ← acc63..32 || acc31..0

```

Exceptions:

Reserved Instruction, DSP Disabled

DPS.W.PH

Dot Product with Subtract on Vector Integer Half-Word Elements

DPSQ_S.W.PH**Dot Product with Subtraction on Fractional Halfword Elements**

31	26	25	21	20	16	15	14	13	12	11	9	8	6	5	3	2	0
P32A 001000		rt		rs		ac	00	011		010		111		111			
6		5		5		2	2	3		3		3		3			

Format: DPSQ_S.W.PH ac, rs, rt**DSP****Purpose:** Dot Product with Subtraction on Fractional Halfword Elements

Element-wise multiplication of two vectors of fractional halfword elements and subtraction of the accumulated 32-bit intermediate products from the specified 64-bit accumulator register, with saturation.

Description: $ac \leftarrow ac - (\text{sat32}(rs_{31..16} * rt_{31..16}) + \text{sat32}(rs_{15..0} * rt_{15..0}))$

Each of the two Q15 fractional word values from registers *rt* and *rs* are multiplied together, and the results left-shifted by one bit position to generate two Q31 fractional format intermediate products. If both multiplicands for either of the multiplications are equal to -1.0 (0x8000 hexadecimal), the resulting intermediate product is saturated to the maximum positive Q31 fractional value (0x7FFFFFFF hexadecimal).

The two intermediate products are then sign-extended and summed to generate a 64-bit, Q32.31 fractional format dot-product result that is subtracted from the specified 64-bit *HI/LO* accumulator to produce a final Q32.31 fractional result.

The value of *ac* can range from 0 to 3; a value of 0 refers to the original *HI/LO* register pair of the MIPS32 architecture.

If saturation occurs as a result of a halfword multiplication, a 1 is written to one of bits 16 through 19 of the *DSPControl* register, within the ouflag field. The value of *ac* determines which of these bits is set: bit 16 corresponds to *ac0*, bit 17 to *ac1*, bit 18 to *ac2*, and bit 19 to *ac3*.

Restrictions:

No data-dependent exceptions are possible.

The operands must be values in the specified format. If they are not, the result is **UNPREDICTABLE** and the values of the operand vectors become **UNPREDICTABLE**.

Operation:

```

ValidateAccessToDSPResources()
tempB31..0 ← multiplyQ15Q15( ac, GPR[rs]31..16, GPR[rt]31..16 )
tempA31..0 ← multiplyQ15Q15( ac, GPR[rs]15..0, GPR[rt]15..0 )
dotp63..0 ← ( (tempB31)32 || tempB31..0 ) + ( (tempA31)32 || tempA31..0 )
tempC63..0 ← ( HI[ac]31..0 || LO[ac]31..0 ) - dotp63..0
( HI[ac]31..0 || LO[ac]31..0 ) ← tempC63..32 || tempC31..0

```

Exceptions:

Reserved Instruction, DSP Disabled

DPSQ_S.W.PH

Dot Product with Subtraction on Fractional Halfword Elements

DPSQ_SA.L.W**Dot Product with Subtraction on Fractional Word Element**

31	26	25	21	20	16	15	14	13	12	11	9	8	6	5	3	2	0	
P32A 001000			rt			rs			ac	01	011		010		111		111	
6			5			5			2	2	3		3		3		3	

Format: DPSQ_SA.L.W *ac*, *rs*, *rt***DSP****Purpose:** Dot Product with Subtraction on Fractional Word Element

Multiplication of two fractional word elements, subtracting the accumulated product from the specified 64-bit accumulator register, with saturation.

Description: $ac \leftarrow \text{sat64}(ac - \text{sat32}(rs_{31..0} * rt_{31..0}))$

The two right-most Q31 fractional word values from registers *rt* and *rs* are multiplied together and the result left-shifted by one bit position to generate a 64-bit Q63 fractional format intermediate product. If both multiplicands are equal to -1.0 (0x80000000 hexadecimal), the intermediate product is saturated to the maximum positive Q63 fractional value (0x7FFFFFFFFFFFFFFF hexadecimal).

The intermediate product is then subtracted from the specified 64-bit *HI/LO* accumulator, creating a Q63 fractional result. If the accumulation results in overflow or underflow, the accumulator is saturated to either the maximum positive or minimum negative Q63 fractional value (0x8000000000000000 hexadecimal), respectively.

The value of *ac* can range from 0 to 3; a value of 0 refers to the original *HI/LO* register pair of the MIPS32 architecture.

If saturation occurs, a 1 is written to one of bits 16 through 19 of the *DSPControl* register, within the *ouflag* field. The value of *ac* determines which of these bits is set: bit 16 corresponds to *ac0*, bit 17 to *ac1*, bit 18 to *ac2*, and bit 19 to *ac3*.

Restrictions:

No data-dependent exceptions are possible.

The operands must be values in the specified format. If they are not, the result is **UNPREDICTABLE** and the values of the operand vectors become **UNPREDICTABLE**.

Operation:

```

ValidateAccessToDSPResources()
dotp63..0 ← multiplyQ31Q31( ac, GPR[rs]31..0, GPR[rt]31..0 )
temp64..0 ← HI[ac]31 || HI[ac]31..0 || LO[ac]31..0
temp64..0 ← temp - dotp63..0
if ( temp64 ≠ temp63 ) then
    if ( temp64 = 1 ) then
        temp63..0 ← 0x8000000000000000
    else
        temp63..0 ← 0x7FFFFFFFFFFFFFFF
    endif
    DSPControlouflag:16+ac ← 1
endif
( HI[ac]31..0 || LO[ac]31..0 ) ← temp63..32 || temp31..0

```

Exceptions:

Reserved Instruction, DSP Disabled

DPSQ_SA.L.W

Dot Product with Subtraction on Fractional Word Element

DPSQX_S.W.PH**Cross Dot Product with Subtraction on Fractional Halfword Elements**

31	26	25	21	20	16	15	14	13	12	11	9	8	6	5	3	2	0
P32A 001000	rt				rs				ac	10	011	010	111	111			
6	5				5				2	2	3	3	3	3			

Format: DPSQX_S.W.PH *ac*, *rs*, *rt***DSP-R2****Purpose:** Cross Dot Product with Subtraction on Fractional Halfword Elements

Element-wise cross multiplication of two vectors of fractional halfword elements and subtraction of the accumulated 32-bit intermediate products from the specified 64-bit accumulator register, with saturation.

Description: $ac \leftarrow ac - (\text{sat32}(rs_{31..16} * rt_{15..0}) + \text{sat32}(rs_{15..0} * rt_{31..16}))$

The left Q15 fractional word value from registers *rt* is multiplied with the right halfword element from register *rs* and the result left-shifted by one bit position to generate a Q31 fractional format intermediate product. Similarly, the right Q15 fractional word value from registers *rt* is multiplied with the left halfword element from register *rs* and the result left-shifted by one bit position to generate a Q31 fractional format intermediate product. If both multiplicands for either of the multiplications are equal to -1.0 (0x8000 hexadecimal), the resulting intermediate product is saturated to the maximum positive Q31 fractional value (0x7FFFFFFF hexadecimal).

The two intermediate products are then sign-extended and summed to generate a 64-bit, Q32.31 fractional format dot-product result that is subtracted from the specified 64-bit *HI/LO* accumulator to produce a final Q32.31 fractional result.

The value of *ac* can range from 0 to 3; a value of 0 refers to the original *HI/LO* register pair of the MIPS32 architecture.

If saturation occurs as a result of a halfword multiplication, a 1 is written to one of bits 16 through 19 of the *DSPControl* register, within the *ouflag* field. The value of *ac* determines which of these bits is set: bit 16 corresponds to *ac0*, bit 17 to *ac1*, bit 18 to *ac2*, and bit 19 to *ac3*.

Restrictions:

No data-dependent exceptions are possible.

The operands must be values in the specified format. If they are not, the result is **UNPREDICTABLE** and the values of the operand vectors become **UNPREDICTABLE**.

Operation:

```

ValidateAccessToDSPResources()
tempB31..0 ← multiplyQ15Q15( ac, GPR[rs]31..16, GPR[rt]15..0 )
tempA31..0 ← multiplyQ15Q15( ac, GPR[rs]15..0, GPR[rt]31..16 )
dotp63..0 ← ( (tempB31)32 || tempB31..0 ) + ( (tempA31)32 || tempA31..0 )
tempC63..0 ← ( HI[ac]31..0 || LO[ac]31..0 ) - dotp63..0
( HI[ac]31..0 || LO[ac]31..0 ) ← tempC63..32 || tempC31..0

function multiplyQ15Q15( acc1..0, a15..0, b15..0 )
  if ( a15..0 = 0x8000 ) and ( b15..0 = 0x8000 ) then
    temp31..0 ← 0x7FFFFFFF
    DSPControlouflag:16+acc ← 1
  else
    temp31..0 ← ( a15..0 * b15..0 ) << 1
  endif
  return temp31..0
endfunction multiplyQ15Q15

```

Exceptions:

Reserved Instruction, DSP Disabled

DPSQX_S.W.PH

Cross Dot Product with Subtraction on Fractional Halfword Elements

DPSQX_SA.W.PH**Cross Dot Product with Subtraction on Fractional Halfword Elements**

31	26	25	21	20	16	15	14	13	12	11	9	8	6	5	3	2	0
P32A 001000	rt				rs				ac	11	011	010	111	111			
6	5				5				2	2	3	3	3	3	3		

Format: DPSQX_SA.W.PH ac, rs, rt**DSP-R2****Purpose:** Cross Dot Product with Subtraction on Fractional Halfword Elements

Element-wise cross multiplication of two vectors of fractional halfword elements and subtraction of the accumulated 32-bit intermediate products from the specified 64-bit accumulator register, with saturation of the accumulator.

Description: $ac \leftarrow \text{sat32}(ac - (\text{sat32}(rs_{31..16} * rt_{15..0}) + \text{sat32}(rs_{15..0} * rt_{31..16})))$

The left Q15 fractional word value from registers *rt* is multiplied with the right halfword element from register *rs* and the result left-shifted by one bit position to generate a Q31 fractional format intermediate product. Similarly, the right Q15 fractional word value from registers *rt* is multiplied with the left halfword element from register *rs* and the result left-shifted by one bit position to generate a Q31 fractional format intermediate product. If both multiplicands for either of the multiplications are equal to -1.0 (0x8000 hexadecimal), the resulting intermediate product is saturated to the maximum positive Q31 fractional value (0x7FFFFFFF hexadecimal).

The two intermediate products are then sign-extended and summed to generate a 64-bit, Q32.31 fractional format dot-product result that is subtracted from the specified 64-bit *HI/LO* accumulator to produce a Q32.31 fractional result. If this result is larger than or equal to +1.0, or smaller than -1.0, it is saturated to the Q31 range.

The value of *ac* can range from 0 to 3; a value of 0 refers to the original *HI/LO* register pair of the MIPS32 architecture.

If saturation occurs as a result of halfword multiplication or accumulation, a 1 is written to one of bits 16 through 19 of the *DSPControl* register, within the *ouflag* field. The value of *ac* determines which of these bits is set: bit 16 corresponds to *ac0*, bit 17 to *ac1*, bit 18 to *ac2*, and bit 19 to *ac3*.

Restrictions:

No data-dependent exceptions are possible.

The operands must be values in the specified format. If they are not, the result is **UNPREDICTABLE** and the values of the operand vectors become **UNPREDICTABLE**.

Operation:

```

ValidateAccessToDSPResources()
tempB31..0 ← multiplyQ15Q15( ac, GPR[rs]31..16, GPR[rt]15..0 )
tempA31..0 ← multiplyQ15Q15( ac, GPR[rs]15..0, GPR[rt]31..16 )
dotp63..0 ← ( (tempB31)32 || tempB31..0 ) + ( (tempA31)32 || tempA31..0 )
tempC63..0 ← ( HI[ac]31..0 || LO[ac]31..0 ) - dotp63..0
if ( tempC63 = 0 ) and ( tempC62..31 ≠ 0 ) then
    tempC63..0 = 032 || 0x7FFFFFFF
    DSPControlouflag:16+acc ← 1
endif
if ( tempC63 = 1 ) and ( tempC62..31 ≠ 132 ) then
    tempC63..0 = 132 || 0x80000000
    DSPControlouflag:16+acc ← 1
endif
( HI[ac]31..0 || LO[ac]31..0 ) ← tempC63..32 || tempC31..0

function multiplyQ15Q15( acc15..0, a15..0, b15..0 )
    if ( a15..0 = 0x8000 ) and ( b15..0 = 0x8000 ) then
        temp31..0 ← 0x7FFFFFFF
        DSPControlouflag:16+acc ← 1
    
```

DPSQX_SA.W.PH**Cross Dot Product with Subtraction on Fractional Halfword Elements**

```
    else
        temp31..0 ← ( a15..0 * b15..0 ) << 1
    endif
    return temp31..0
endfunction multiplyQ15Q15
```

Exceptions:

Reserved Instruction, DSP Disabled

DPSU.H.QBL**Dot Product with Subtraction on Vector Unsigned Byte Elements**

31	26	25	21	20	16	15	14	13	12	11	9	8	6	5	3	2	0
P32A 001000	rt				rs				ac	10	010	010	111			111	
6	5				5				2	2	3	3	3			3	

Format: DPSU.H.QBL *ac*, *rs*, *rt***DSP****Purpose:** Dot Product with Subtraction on Vector Unsigned Byte Elements

Element-wise multiplication of two left-most elements from the four elements of each of two vectors of unsigned bytes, subtracting the sum of the products from the specified 64-bit accumulator register.

Description: $ac \leftarrow ac - \text{zero_extend}((rs_{31..24} * rt_{31..24}) + (rs_{23..16} * rt_{23..16}))$

The two left-most elements of the four unsigned byte elements of each of registers *rt* and *rs* are multiplied together using unsigned arithmetic to generate two 16-bit unsigned intermediate products. The intermediate products are then zero-extended to 64 bits and subtracted from the specified 64-bit *HI/LO* accumulator. The result of the subtraction is written back to the specified 64-bit *HI/LO* accumulator.

The value of *ac* can range from 0 to 3; a value of 0 refers to the original *HI/LO* register pair of the MIPS32 architecture.

This instruction does not set any bits in the *ouflag* field in the *DSPControl* register.

Restrictions:

No data-dependent exceptions are possible.

The operands must be values in the specified format. If they are not, the result is **UNPREDICTABLE** and the values of the operand vectors become **UNPREDICTABLE**.

Operation:

```

ValidateAccessToDSPResources()
tempB15..0 ← multiplyU8U8( GPR[rs]31..24, GPR[rt]31..24 )
tempA15..0 ← multiplyU8U8( GPR[rs]23..16, GPR[rt]23..16 )
dotp63..0 ← ( 048 || tempB15..0 ) + ( 048 || tempA15..0 )
tempC63..0 ← ( HI[ac]31..0 || LO[ac]31..0 ) - dotp63..0
( HI[ac]31..0 || LO[ac]31..0 ) ← tempC63..32 || tempC31..0

```

Exceptions:

Reserved Instruction, DSP Disabled

DPSU.H.QBL

Dot Product with Subtraction on Vector Unsigned Byte Elements

DPSU.H.QBR**Dot Product with Subtraction on Vector Unsigned Byte Elements**

31	26	25	21	20	16	15	14	13	12	11	9	8	6	5	3	2	0
P32A 001000	rt				rs				ac	11	010	010	111	111			
6	5				5				2	2	3	3	3	3	3		

Format: DPSU.H.QBR *ac*, *rs*, *rt***DSP****Purpose:** Dot Product with Subtraction on Vector Unsigned Byte Elements

Element-wise multiplication of the two right-most elements of the four elements of each of two vectors of unsigned bytes, subtracting the sum of the products from the specified 64-bit accumulator register.

Description: $ac \leftarrow ac - \text{zero_extend}((rs_{15..8} * rt_{15..8}) + (rs_{7..0} * rt_{7..0}))$

The two right-most elements of the four unsigned byte elements of each of registers *rt* and *rs* are multiplied together using unsigned arithmetic to generate two 16-bit unsigned intermediate products. The intermediate products are then zero-extended to 64 bits and subtracted from the specified 64-bit *HI/LO* accumulator.

The value of *ac* can range from 0 to 3; a value of 0 refers to the original *HI/LO* register pair of the MIPS32 architecture.

This instruction does not set any bits in the *ouflag* field in the *DSPControl* register.

Restrictions:

No data-dependent exceptions are possible.

The operands must be values in the specified format. If they are not, the result is **UNPREDICTABLE** and the values of the operand vectors become **UNPREDICTABLE**.

Operation:

```

ValidateAccessToDSPResources()
tempB15..0 ← multiplyU8U8( GPR[rs]15..8, GPR[rt]15..8 )
tempA15..0 ← multiplyU8U8( GPR[rs]7..0, GPR[rt]7..0 )
dotp63..0 ← ( 048 || tempB15..0 ) + ( 048 || tempA15..0 )
tempC63..0 ← ( HI[ac]31..0 || LO[ac]31..0 ) - dotp63..0
( HI[ac]31..0 || LO[ac]31..0 ) ← tempC63..32 || tempC31..0

```

Exceptions:

Reserved Instruction, DSP Disabled

DPSU.H.QBR

Dot Product with Subtraction on Vector Unsigned Byte Elements

DPSX.W.PH**Cross Dot Product with Subtract on Vector Integer Halfword Elements**

31	26	25	21	20	16	15	14	13	12	11	9	8	6	5	3	2	0
P32A 001000			rt			rs			ac	01	010	010	111	111			
6			5			5			2	2	3	3	3	3	3		

Format: DPSX.W.PH *ac*, *rs*, *rt***DSP-R2****Purpose:** Cross Dot Product with Subtract on Vector Integer Halfword Elements

Generate the cross dot-product of two integer halfword vector elements using full-size intermediate products and then subtract from the specified accumulator register.

Description: $ac \leftarrow ac - ((rs_{31..16} * rt_{15..0}) + (rs_{15..0} * rt_{31..16}))$

The left halfword integer value from register *rt* is multiplied with the right halfword element from register *rs* to create an integer word result. Similarly, the right halfword integer value from register *rt* is multiplied with the left halfword element from register *rs* to create the second integer word result. These two products are summed to generate the dot-product result, which is then subtracted from the specified 64-bit *HI/LO* accumulator, creating a 64-bit integer result.

The value of *ac* selects an accumulator numbered from 0 to 3. When *ac*=0, this refers to the original *HI/LO* register pair of the MIPS32 architecture.

This instruction will not set any bits of the *ouflag* field in the *DSPControl* register.

Restrictions:

No data-dependent exceptions are possible.

The operands must be a value in the specified format. If they are not, the result is **UNPREDICTABLE** and the values of the operand vectors become **UNPREDICTABLE**.

Operation:

```

ValidateAccessToDSPResources()
tempB31..0 ← (GPR[rs]31..16 * GPR[rt]15..0)
tempA31..0 ← (GPR[rs]15..0 * GPR[rt]31..16)
dotp32..0 ← ( (tempB31) || tempB31..0 ) + ( (tempA31) || tempA31..0 )
acc63..0 ← ( HI[ac]31..0 || LO[ac]31..0 ) - ( (dotp32)31 || dotp32..0 )
( HI[ac]31..0 || LO[ac]31..0 ) ← acc63..32 acc31..0

```

Exceptions:

Reserved Instruction, DSP Disabled

DPSX.W.PH

Cross Dot Product with Subtract on Vector Integer Halfword Elements

EXTP**Extract Fixed Bitfield From Arbitrary Position in Accumulator to GPR**

31	26	25	21	20	16	15	14	13	12	11	9	8	6	5	3	2	0
P32A 001000	rt		size		ac	10	011		001	111		111					
6	5		5		2	2	3		3	3		3		3			

Format: EXTP *rt*, *ac*, *size***DSP****Purpose:** Extract Fixed Bitfield From Arbitrary Position in Accumulator to GPR

Extract *size*+1 contiguous bits from a 64-bit accumulator from a position specified in the *DSPControl* register, writing the bits to a GPR with zero-extension.

Description: $rt \leftarrow \text{zero_extend}(ac_{pos..pos-size})$

A set of *size*+1 contiguous bits are extracted from an arbitrary position in accumulator *ac*, zero-extended to 32 bits, and then written to register *rt*.

The bit position, *start_pos*, of the first bit of the contiguous set to extract is specified by the *pos* field in bits 0 through 5 of the *DSPControl* register. The last bit in the set is *start_pos* - *size*, where *size* is specified in the instruction.

The value of *ac* can range from 0 to 3. When *ac*=0, this refers to the original *HI/LO* register pair of the MIPS32 architecture. After the execution of this instruction, accumulator *ac* remains unmodified.

If $start_pos - (size + 1) \geq -1$, the extraction is valid, otherwise the extraction is invalid and is said to have failed. The value of the destination register is **UNPREDICTABLE** when the extraction is invalid. Upon an invalid extraction this instruction writes a 1 to bit 14, the Extract Failed Indicator (EFI) bit of the *DSPControl* register, and 0 otherwise.

The values of bits 0 to 5 in the *pos* field of the *DSPControl* register are unchanged by this instruction.

Restrictions:

No data-dependent exceptions are possible.

The operands must be values in the specified format. If they are not, the results are **UNPREDICTABLE** and the values of the operand vectors become **UNPREDICTABLE**.

Operation:

```

ValidateAccessToDSPResources()
start_pos5..0 ← DSPControlpos:5..0
if ( start_pos - (size+1) >= -1 ) then
    tempsize..0 ← ( HI[ac]31..0 || LO[ac]31..0 )start_pos..start_pos-size
    temp31..0 ← 0(32-(size+1)) || tempsize..0
    GPR[rt]31..0 ← temp31..0
    DSPControlEFI:14 ← 0
else
    DSPControlEFI:14 ← 1
    GPR[rt] ← UNPREDICTABLE
endif

```

Exceptions:

Reserved Instruction, DSP Disabled

EXTP

Extract Fixed Bitfield From Arbitrary Position in Accumulator to GPR

EXTPDP Extract Fixed Bitfield From Arbitrary Position in Accumulator to GPR and Decrement Pos

31	26	25	21	20	16	15	14	13	12	11	9	8	6	5	3	2	0
P32A 001000						rt		size			ac	11	011	001	111	111	
6						5		5			2	2	3	3	3	3	

Format: EXTPDP rt, ac, size**DSP****Purpose:** Extract Fixed Bitfield From Arbitrary Position in Accumulator to GPR and Decrement Pos

Extract $size+1$ contiguous bits from a 64-bit accumulator from a position specified in the *DSPControl* register, writing the bits to a GPR with zero-extension and modifying the extraction position.

Description: $rt \leftarrow \text{zero_extend}(ac_{pos..pos-size})$; $DSPControl_{pos:5..0} -= (size+1)$

A set of $size+1$ contiguous bits are extracted from an arbitrary position in accumulator *ac*, zero-extended to 32 bits, then written to register *rt*.

The bit position, *start_pos*, of the first bit of the contiguous set to extract is specified by the *pos* field in bits 0 through 5 of the *DSPControl* register. The position of the last bit in the extracted set is $start_pos - size$, where the *size* argument is specified in the instruction.

The value of *ac* can range from 0 to 3. When $ac=0$, this refers to the original *HI/LO* register pair of the MIPS32 architecture. After the execution of this instruction, accumulator *ac* remains unmodified.

If $start_pos - (size + 1) \geq -1$, the extraction is valid and the value of the *pos* field in the *DSPControl* register is decremented by $size+1$. Otherwise, the extraction is invalid and is said to have failed. The value of the destination register is **UNPREDICTABLE** when the extraction is invalid, and the value of the *pos* field in the *DSPControl* register (bits 0 through 5) is not modified.

Upon an invalid extraction this instruction writes a 1 to bit 14, the Extract Failed Indicator (EFI) bit of the *DSPControl* register, and 0 otherwise.

Restrictions:

No data-dependent exceptions are possible.

The operands must be values in the specified format. If they are not, the results are **UNPREDICTABLE** and the values of the operand vectors become **UNPREDICTABLE**.

Operation:

```

ValidateAccessToDSPResources()
start_pos5..0 ← DSPControlpos:5..0
if ( start_pos - (size+1) >= -1 ) then
    tempsize..0 ← ( HI[ac]31..0 || LO[ac]31..0 )start_pos..start_pos-size
    GPR[rt] ← 0(GPRLEN-(size+1)) || tempsize..0
    DSPControlpos:5..0 ← DSPControlpos:5..0 - (size + 1)
    DSPControlEFI:14 ← 0
else
    DSPControlEFI:14 ← 1
    GPR[rt] ← UNPREDICTABLE
endif

```

Exceptions:

Reserved Instruction, DSP Disabled

EXTPDP	Extract Fixed Bitfield From Arbitrary Position in Accumulator to GPR and Decrement Pos
---------------	---

EXTDPDV Extract Variable Bitfield From Arbitrary Position in Accumulator to GPR and Decrement Pos

31	26	25	21	20	16	15	14	13	12	11	9	8	6	5	3	2	0
P32A 001000	rt				rs				ac	11	100	010	111	111			
6	5				5				2	2	3	3	3	3	3		

Format: EXTPDPV *rt*, *ac*, *rs***DSP****Purpose:** Extract Variable Bitfield From Arbitrary Position in Accumulator to GPR and Decrement Pos

Extract a fixed number of contiguous bits from a 64-bit accumulator from a position specified in the *DSPControl* register, writing the bits to a GPR with zero-extension and modifying the extraction position.

Description: $rt \leftarrow \text{zero_extend}(ac_{pos..pos-GPR[rs][4:0]})$; $DSPControl_{pos:5..0} -= (GPR[rs]_{4..0} + 1)$

A fixed number of contiguous bits are extracted from an arbitrary position in accumulator *ac*, zero-extended to 32 bits, then written to destination register *rt*. The number of bits extracted is *size*+1, where *size* is specified by the five least-significant bits in register *rs*, interpreted as a five-bit unsigned integer. The remaining bits in register *rs* are ignored.

The bit position, *start_pos*, of the first bit of the contiguous set to extract is specified by the *pos* field in bits 0 through 5 of the *DSPControl* register. The position of the last bit in the extracted set is *start_pos* - *size*.

The value of *ac* can range from 0 to 3. When *ac*=0, this refers to the original *H/LO* register pair of the MIPS32 architecture. After the execution of this instruction, accumulator *ac* remains unmodified.

If $start_pos - (size + 1) \geq -1$, the extraction is valid and the value of the *pos* field in the *DSPControl* register is decremented by *size*+1. Otherwise, the extraction is invalid and is said to have failed. The value of the destination register is **UNPREDICTABLE** when the extraction is invalid, and the value of the *pos* field in the *DSPControl* register (bits 0 through 5) is not modified.

Upon an invalid extraction this instruction writes a 1 to bit 14, the Extract Failed Indicator (EFI) bit of the *DSPControl* register, and 0 otherwise.

Restrictions:

No data-dependent exceptions are possible.

The operands must be values in the specified format. If they are not, the results are **UNPREDICTABLE** and the values of the operand vectors become **UNPREDICTABLE**.

Operation:

```

ValidateAccessToDSPResources()
start_pos5..0 ← DSPControlpos:5..0
size4..0 ← GPR[rs]4..0
if ( start_pos - (size+1) >= -1 ) then
    tempsize..0 ← ( HI[ac]31..0 || LO[ac]31..0 )start_pos..start_pos-size
    GPR[rt] ← 0(GPRLEN-(size+1)) || tempsize..0
    DSPControlpos:5..0 ← DSPControlpos:5..0 - (size + 1)
    DSPControlEFI:14 ← 0
else
    DSPControlEFI:14 ← 1
    GPR[rt] ← UNPREDICTABLE
endif

```

Exceptions:

Reserved Instruction, DSP Disabled

EXTDPV Extract Variable Bitfield From Arbitrary Position in Accumulator to GPR and Decrement Pos

EXTPV**Extract Variable Bitfield From Arbitrary Position in Accumulator to GPR**

31	26	25	21	20	16	15	14	13	12	11	9	8	6	5	3	2	0
P32A 001000		rt			rs			ac	10	100		010	111		111		
6		5			5			2	2	3		3	3		3		

Format: EXTPV *rt*, *ac*, *rs***DSP****Purpose:** Extract Variable Bitfield From Arbitrary Position in Accumulator to GPR

Extract a variable number of contiguous bits from a 64-bit accumulator from a position specified in the *DSPControl* register, writing the bits to a GPR with zero-extension.

Description: $rt \leftarrow \text{zero_extend}(ac_{pos..pos-rs[4:0]})$

A variable number of contiguous bits are extracted from an arbitrary position in accumulator *ac*, zero-extended to 32 bits, then written to register *rt*. The number of bits extracted is *size*+1, where *size* is specified by the five least-significant bits in register *rs*, interpreted as a five-bit unsigned integer. The remaining bits in register *rs* are ignored.

The position of the first bit of the contiguous set to extract, *start_pos*, is specified by the *pos* field in bits 0 through 5 of the *DSPControl* register. The position of the last bit in the contiguous set is *start_pos* - *size*.

The value of *ac* can range from 0 to 3. When *ac*=0, this refers to the original *HI/LO* register pair of the MIPS32 architecture. After the execution of this instruction, accumulator *ac* remains unmodified.

An extraction is valid if $start_pos - (size + 1) \geq -1$; otherwise, the extraction is invalid and is said to have failed. The value of the destination register is **UNPREDICTABLE** when the extraction is invalid. Upon an invalid extraction this instruction writes a 1 to bit 14, the Extract Failed Indicator (EFI) bit of the *DSPControl* register, and 0 otherwise.

The values of bits 0 to 5 in the *pos* field of the *DSPControl* register are unchanged by this instruction.

Restrictions:

No data-dependent exceptions are possible.

The operands must be values in the specified format. If they are not, the results are **UNPREDICTABLE** and the values of the operand vectors become **UNPREDICTABLE**.

Operation:

```

ValidateAccessToDSPResources()
start_pos5..0 ← DSPControlpos:5..0
size4..0 ← GPR[rs]4..0
if ( start_pos - (size+1) >= -1 ) then
    tempsize..0 ← ( HI[ac]31..0 || LO[ac]31..0 )start_pos..start_pos-size
    GPR[rt] ← 0(GPRLEN-(size+1)) || tempsize..0
    DSPControlEFI:14 ← 0
else
    DSPControlEFI:14 ← 1
    GPR[rt] ← UNPREDICTABLE
endif

```

Exceptions:

Reserved Instruction, DSP Disabled

EXTPV

Extract Variable Bitfield From Arbitrary Position in Accumulator to GPR

EXTR[_RS].W**Extract Word Value With Right Shift From Accumulator to GPR**

31	26	25	21	20	16	15	14	13	12	11	9	8	6	5	3	2	0
EXTR.W																	
P32A 001000	rt		shift		ac	00	111		001		111		111				
EXTR_R.W																	
P32A 001000	rt		shift		ac	01	111		001		111		111				
EXTR_RS.W																	
P32A 001000	rt		shift		ac	10	111		001		111		111				
6	5	5	2	2	3	3	3	3	3	3	3	3	3	3	3	3	3

Format: EXTR[_RS].W
EXTR.W rt, ac, shift
EXTR_R.W rt, ac, shift
EXTR_RS.W rt, ac, shift

DSP
DSP
DSP

Purpose: Extract Word Value With Right Shift From Accumulator to GPR

Extract a word value from a 64-bit accumulator to a GPR with right shift, and with optional rounding or rounding and saturation.

Description: $rt \leftarrow \text{sat32}(\text{round}(\text{ac} \gg \text{shift}))$

The value in accumulator *ac* is shifted right by *shift* bits with sign extension (arithmetic shift right). The 32 least-significant bits of the shifted value are then written to the destination register *rs*.

The rounding variant of the instruction adds a 1 at the most-significant discarded bit position. The 32 least-significant bits of the rounded result are then written to the destination register.

The rounding and saturating variant of the instruction adds a 1 at the most-significant discarded bit position. If the rounding operation results in an overflow, the shifted value is clamped to the maximum positive Q31 fractional value (0x7FFFFFFF hexadecimal). The rounded and saturated result is then written to the destination register.

The value of *ac* can range from 0 to 3. When *ac*=0, this refers to the original *H/L*O register pair of the MIPS32 architecture. After the execution of this instruction, *ac* remains unmodified.

For all variants of the instruction, including EXTR.W, bit 23 of the *DSPControl* register is set to 1 if either of the rounded or non-rounded calculation results in overflow or saturation.

Restrictions:

No data-dependent exceptions are possible.

The operands must be values in the specified format. If they are not, the results are **UNPREDICTABLE** and the values of the operand vectors become **UNPREDICTABLE**.

Operation:

```
EXTR.W
    ValidateAccessToDSPResources()
    temp64..0 ← _shiftShortAccRightArithmetic( ac, shift )
    if ( ( temp64..32 ≠ 0 ) and ( temp64..32 ≠ 0x1FFFFFFFF ) ) then
        DSPControlouflag:23 ← 1
    endif
    GPR[rt]31..0 ← temp32..1
    temp64..0 ← temp + 1
    if ( ( temp64..32 ≠ 0 ) and ( temp64..32 ≠ 0x1FFFFFFFF ) ) then
        DSPControlouflag:23 ← 1
```

EXTR[_RS].W**Extract Word Value With Right Shift From Accumulator to GPR**

```

endif

EXTR_R.W
    ValidateAccessToDSPResources()
    temp64..0 ← _shiftShortAccRightArithmetic( ac, shift )
    if ( ( temp64..32 ≠ 0 ) and ( temp64..32 ≠ 0x1FFFFFFF ) ) then
        DSPControlouflag:23 ← 1
    endif
    temp64..0 ← temp + 1
    if ( ( temp64..32 ≠ 0 ) and ( temp64..32 ≠ 0x1FFFFFFF ) ) then
        DSPControlouflag:23 ← 1
    endif
    GPR[rt]31..0 ← temp32..1

EXTR_RS.W
    ValidateAccessToDSPResources()
    temp64..0 ← _shiftShortAccRightArithmetic( ac, shift )
    if ( ( temp64..32 ≠ 0 ) and ( temp64..32 ≠ 0x1FFFFFFF ) ) then
        DSPControlouflag:23 ← 1
    endif
    temp64..0 ← temp + 1
    if ( ( temp64..32 ≠ 0 ) and ( temp64..32 ≠ 0x1FFFFFFF ) ) then
        if ( temp64 = 0 ) then
            temp32..1 ← 0x7FFFFFFF
        else
            temp32..1 ← 0x80000000
        endif
        DSPControlouflag:23 ← 1
    endif
    GPR[rt]31..0 ← temp32..1

function _shiftShortAccRightArithmetic( ac1..0, shift4..0 )
    if ( shift4..0 = 0 ) then
        temp64..0 ← ( HI[ac]31..0 || LO[ac]31..0 || 0 )
    else
        temp64..0 ← ( (HI[ac]31)shift || HI[ac]31..0 || LO[ac]31..shift-1 )
    endif
    return temp64..0
endfunction _shiftShortAccRightArithmetic

```

Exceptions:

Reserved Instruction, DSP Disabled

EXTR_S.H**Extract Halfword Value From Accumulator to GPR With Right Shift and Saturate**

31	26	25	21	20	16	15	14	13	12	11	9	8	6	5	3	2	0
P32A 001000						rt		shift			ac	11	111	001	111	111	
6						5		5			2	2	3	3	3	3	

Format: EXTR_S.H *rt*, *ac*, *shift***DSP****Purpose:** Extract Halfword Value From Accumulator to GPR With Right Shift and Saturate

Extract a halfword value from a 64-bit accumulator to a GPR with right shift and saturation.

Description: $rt \leftarrow \text{sat16}(ac \gg \text{shift})$

The value in the 64-bit accumulator *ac* is shifted right by *shift* bits with sign extension (arithmetic shift right). The 64-bit value is then saturated to 16-bits, sign extended to 32 bits, and written to the destination register *rt*. The shift argument is provided in the instruction.

The value of *ac* can range from 0 to 3. When *ac*=0, this refers to the original *H/LO* register pair of the MIPS32 architecture. After the execution of this instruction, *ac* remains unmodified.

This instruction sets bit 23 of the *DSPControl* register in the *ouflag* field if the operation results in saturation.

Restrictions:

No data-dependent exceptions are possible.

The operands must be values in the specified format. If they are not, the results are **UNPREDICTABLE** and the values of the operand vectors become **UNPREDICTABLE**.

Operation:

```

ValidateAccessToDSPResources()
temp63..0 ← shiftShortAccRightArithmetic( ac, shift )
if ( temp63..0 > 0x00000000000007FFF ) then
    temp31..0 ← 0x00007FFF
    DSPControlouflag:23 ← 1
else if ( temp63..0 < 0xFFFFFFFFFFFF8000 ) then
    temp31..0 ← 0xFFFF8000
    DSPControlouflag:23 ← 1
endif
GPR[rt]31..0 ← temp31..0

function shiftShortAccRightArithmetic( ac1..0, shift4..0 )
    sign ← HI[ac]31
    if ( shift = 0 ) then
        temp63..0 ← HI[ac]31..0 || LO[ac]31..0
    else
        temp63..0 ← signshift || (( HI[ac]31..0 || LO[ac]31..0 ) >> shift )
    endif
    if ( sign ≠ temp31 ) then
        DSPControlouflag:23 ← 1
    endif
    return temp63..0
endfunction shiftShortAccRightArithmetic

```

Exceptions:

Reserved Instruction, DSP Disabled

EXTR_S.H

Extract Halfword Value From Accumulator to GPR With Right Shift and Saturate

EXTRV[_RS].W**Extract Word Value With Variable Right Shift From Accumulator to GPR**

31	26	25	21	20	16	15	14	13	12	11	9	8	6	5	3	2	0
EXTRV.W																	
P32A 001000		rt		rs		ac	00	111			010		111		111		
EXTRV_R.W																	
P32A 001000		rt		rs		ac	01	111			010		111		111		
EXTRV_RS.W																	
P32A 001000		rt		rs		ac	10	111			010		111		111		
6		5		5		2	2	3			3		3		3		

Format: EXTRV[_RS].W

EXTRV.W rt, ac, rs

EXTRV_R.W rt, ac, rs

EXTRV_RS.W rt, ac, rs

DSP**DSP****DSP****Purpose:** Extract Word Value With Variable Right Shift From Accumulator to GPR

Extract a word value from a 64-bit accumulator to a GPR with variable right shift, and with optional rounding or rounding and saturation.

Description: $rt \leftarrow \text{sat32}(\text{round}(ac \gg rs_{5..0}))$

The value in accumulator *ac* is shifted right by *shift* bits with sign extension (arithmetic shift right). The lower 32 bits of the shifted value are then written to the destination register *rt*. The number of bits to shift is given by the five least-significant bits of register *rs*; the remaining bits of *rs* are ignored.

The rounding variant of the instruction adds a 1 at the most-significant discarded bit position. The 32 least-significant bits of the rounded result are then written to the destination register.

The rounding and saturating variant of the instruction adds a 1 at the most-significant discarded bit position. If the rounding operation results in an overflow, the shifted value is clamped to the maximum positive Q31 fractional value (0x7FFFFFFF hexadecimal). The rounded and saturated result is then written to the destination register.

The value of *ac* can range from 0 to 3. When *ac*=0, this refers to the original *H/L*O register pair of the MIPS32 architecture. After the execution of this instruction, *ac* remains unmodified.

For all variants of the instruction, including EXTRV.W, bit 23 of the *DSPControl* register is set to 1 if either of the rounded or non-rounded calculation results in overflow or saturation.

Restrictions:

No data-dependent exceptions are possible.

The operands must be values in the specified format. If they are not, the results are **UNPREDICTABLE** and the values of the operand vectors become **UNPREDICTABLE**.

Operation:

```

EXTRV.W
    ValidateAccessToDSPResources()
    temp64..0 ← _shiftShortAccRightArithmetic( ac, GPR[rt]4..0 )
    if ( ( temp64..32 ≠ 0 ) and ( temp64..32 ≠ 0x1FFFFFFFF ) ) then
        DSPControlouflag:23 ← 1
    endif
    GPR[rt]31..0 ← temp32..1
    temp64..0 ← temp + 1
    if ( ( temp64..32 ≠ 0 ) and ( temp64..32 ≠ 0x1FFFFFFFF ) ) then

```

EXTRV[_RS].W**Extract Word Value With Variable Right Shift From Accumulator to GPR**

```

        DSPControl_ouflag:23 ← 1
    endif

EXTRV_R.W
    ValidateAccessToDSPResources()
    temp64..0 ← _shiftShortAccRightArithmetic( ac, GPR[rt]4..0 )
    if ( ( temp64..32 ≠ 0 ) and ( temp64..32 ≠ 0x1FFFFFFF ) ) then
        DSPControl_ouflag:23 ← 1
    endif
    temp64..0 ← temp + 1
    if ( ( temp64..32 ≠ 0 ) and ( temp64..32 ≠ 0x1FFFFFFF ) ) then
        DSPControl_ouflag:23 ← 1
    endif
    GPR[rt]31..0 ← temp32..1

EXTRV_RS.W
    ValidateAccessToDSPResources()
    temp64..0 ← _shiftShortAccRightArithmetic( ac, GPR[rt]4..0 )
    if ( ( temp64..32 ≠ 0 ) and ( temp64..32 ≠ 0x1FFFFFFF ) ) then
        DSPControl_ouflag:23 ← 1
    endif
    temp64..0 ← temp + 1
    if ( ( temp64..32 ≠ 0 ) and ( temp64..32 ≠ 0x1FFFFFFF ) ) then
        if ( temp64 = 0 ) then
            temp32..1 ← 0x7FFFFFFF
        else
            temp32..1 ← 0x80000000
        endif
        DSPControl_ouflag:23 ← 1
    endif
    GPR[rt]31..0 ← temp32..1

```

Exceptions:

Reserved Instruction, DSP Disabled

EXTRV_S.H Extract Halfword Value Variable From Accumulator to GPR With Right Shift and Saturate

31	26	25	21	20	16	15	14	13	12	11	9	8	6	5	3	2	0
P32A 001000						rt		rs		ac	11	111	010	111	111		
6						5		5		2	2	3	3	3	3		

Format: EXTRV_S.H *rt*, *ac*, *rs***DSP****Purpose:** Extract Halfword Value Variable From Accumulator to GPR With Right Shift and Saturate

Extract a halfword value from a 64-bit accumulator to a GPR with right shift and saturation.

Description: $rt \leftarrow \text{sat16}(ac \gg rs_{4..0})$

The value in the 64-bit accumulator *ac* is shifted right by *shift* bits with sign extension (arithmetic shift right). The 64-bit value is then saturated to 16-bits and sign-extended to 32 bits before being written to the destination register *rt*. The five least-significant bits of register *rs* provide the shift argument, interpreted as a five-bit unsigned integer; the remaining bits in *rs* are ignored.

The value of *ac* can range from 0 to 3. When *ac*=0, this refers to the original *H/L*O register pair of the MIPS32 architecture. After the execution of this instruction, *ac* remains unmodified.

This instruction sets bit 23 of the *DSPControl* register in the *ouflag* field if the operation results in saturation.

Restrictions:

No data-dependent exceptions are possible.

The operands must be values in the specified format. If they are not, the results are **UNPREDICTABLE** and the values of the operand vectors become **UNPREDICTABLE**.

Operation:

```

ValidateAccessToDSPResources()
shift4..0 ← GPR[rs]4..0
temp31..0 ← shiftShortAccRightArithmetic( ac, shift )
if ( temp63..0 > 0x00000000000007FFF ) then
    temp31..0 ← 0x00007FFF
    DSPControl23 ← 1
else if ( temp63..0 < 0xFFFFFFFFFFFF8000 ) then
    temp31..0 ← 0xFFFFF8000
    DSPControl23 ← 1
endif
GPR[rt]31..0 ← temp31..0

```

Exceptions:

Reserved Instruction, DSP Disabled

EXTRV_S.H **Extract Halfword Value Variable From Accumulator to GPR With Right Shift and Saturate**

INSV**Insert Bit Field Variable**

31	26	25	21	20	16	15	9	8	6	5	3	2	0
P32A 001000			rt		rs		0100000			100		111	
6			5		5		7			3		3	

Format: INSV *rt*, *rs***DSP****Purpose:** Insert Bit Field VariableTo merge a right-justified bit field from register *rs* into a specified field in register *rt*.**Description:** $rt \leftarrow \text{InsertFieldVar}(rt, rs, \text{Scount}, \text{Pos})$

The *DSPControl* register provides the *size* value from the *Scount* field, and the *pos* value from the *pos* field. The right-most *size* bits from register *rs* are merged into the value from register *rt* starting at bit position *pos*. The result is put back in register *rt*. These *pos* and *size* values are converted by the instruction into the fields *msb* (the most significant bit of the field), and *lsb* (least significant bit of the field), as follows:

```

pos ← DSPControl5..0
size ← DSPControl12..7
msb ← pos+size-1
lsb ← pos

```

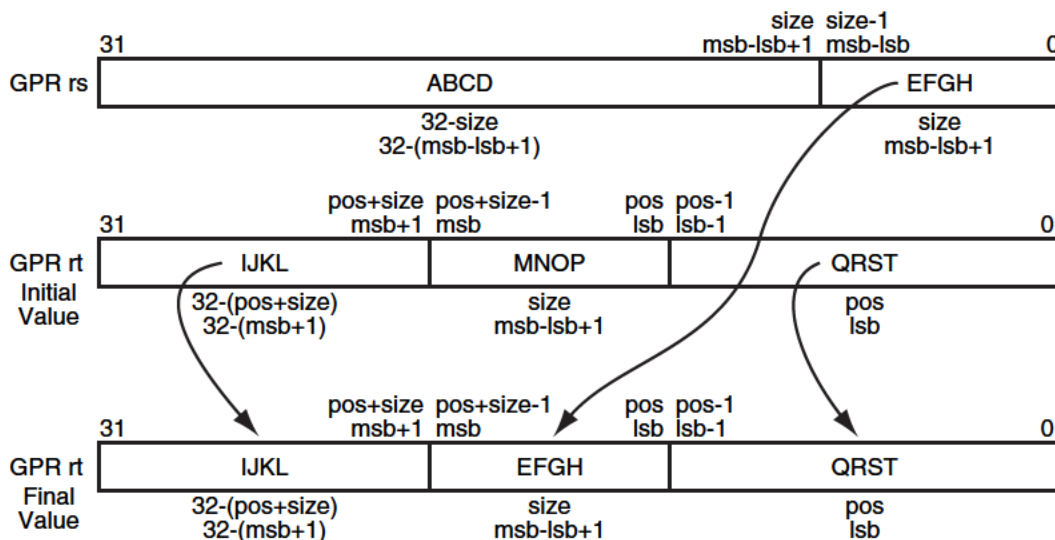
The values of *pos* and *size* must satisfy all of the following relations, or the instruction results in UNPREDICTABLE results:

```

0 ≤ pos < 32
0 < size ≤ 32
0 < pos+size ≤ 32

```

Figure 6.3 shows the symbolic operation of the instruction.

Figure 6.3 Operation of the INSV Instruction**Restrictions:**The operation is **UNPREDICTABLE** if *lsb* > *msb*.**Operation:**

ValidateAccessToDSPResources()

INSV**Insert Bit Field Variable**

```

if (lsb > msb) then
    UNPREDICTABLE
endif
GPR[rt]31..0 ← GPR[rt]31..msb+1 || GPR[rs]msb-lsb..0 || GPR[rt]lsb-1..0

```

Exceptions:

Reserved Instruction, DSP Disabled

Implementation Notes

The destination of this instruction is register *rt* because that register is used as both a source and destination of the instruction. Since most implementations have potential critical paths around source register decode, and typically decode registers *rs* and *rt* as source registers, the instruction is defined with the destination as register *rt* instead of register *rd* to minimize the impact on source register decode.

One implementation method is to shift the register *rs* value left by *lsb* bits and merge that value into the register *rt* value based on a merge mask. The merge mask has a 1 in every bit position from which the corresponding output bit comes from register *rs* and a 0 in every bit position from which the corresponding output bit comes from register *rt*. The mask can be calculated by subtracting two constants generated from the fields of the instruction, as follows:

$$\begin{aligned}
 k1 &\leftarrow 0^{32-lsb-1} || 1 || 0^{lsb} \\
 k2 &\leftarrow (0^{33-(msb+2)} || 1 || 0^{msb+1})_{31..0} \\
 merge_mask &\leftarrow k2 - k1
 \end{aligned}$$

Some implementations may choose to use the ALU to calculate the *merge_mask* in parallel with shifting the register *rs* value to the left, then using the *merge_mask* to bit-select from the register *rt* value or the shifted register *rs* value.

LBUX**Load Unsigned Byte Indexed**

Format: LBUX rd, index(base)
LBUX rd, rs(rt)

DSP
Replaced with LBUX in nanoMIPS

Purpose: Load Unsigned Byte Indexed

To load a byte from memory as an unsigned value, using indexed addressing.

Description: $rd \leftarrow \text{memory}[\text{base} + \text{index}]$

The contents of GPR *index* is added to the contents of GPR *base* to form an effective address. The contents of the 8-bit byte at the memory location specified by the aligned effective address are fetched, zero-extended to the GPR register length and placed in GPR *rd*.

Restrictions:

None.

Operation:

```

ValidateAccessToDSPResources()
vAddr31..0 ← GPR[index]31..0 + GPR[base]31..0
( pAddr, CCA ) ← AddressTranslation( vAddr, DATA, LOAD )
pAddr ← pAddrPSIZE-1..2 || ( pAddr1..0 xor ReverseEndian2 )
memwordGPRLEN..0 ← LoadMemory ( CCA, BYTE, pAddr, vAddr, DATA )
GPR[rd]31..0 ← zero_extend( memword7..0 )

```

Exceptions:

Reserved Instruction, DSP Disabled, TLB Refill, TLB Invalid, Bus Error, Address Error, Watch

LBUX

Load Unsigned Byte Indexed

LHX**Load Halfword Indexed**

Format: LHX rd, index(base)
LHX rd, rs(rt)

DSP
Replaced with LHX in nanoMIPS

Purpose: Load Halfword Indexed

To load a halfword value from memory as a signed value, using indexed addressing.

Description: $rd \leftarrow \text{memory}[\text{base} + \text{index}]$

The contents of GPR *index* is added to the contents of GPR *base* to form an effective address. The contents of the 16-bit halfword at the memory location specified by the aligned effective address are fetched, sign-extended to the length of the destination GPR, and placed in GPR *rd*.

Restrictions:

The effective address must be naturally-aligned. If the least-significant bit of the effective address is non-zero, an Address Error exception occurs.

Operation:

```

ValidateAccessToDSPResources()
vAddr31..0 ← GPR[index]31..0 + GPR[base]31..0
if ( vAddr0 ≠ 0 ) then
    SignalException( AddressError )
endif
( pAddr, CCA ) ← AddressTranslation( vAddr, DATA, LOAD )
halfwordGPRLEN..0 ← LoadMemory( CCA, HALFWORD, pAddr, vAddr, DATA )
GPR[rd]31..0 ← sign_extend( halfword15..0 )

```

Exceptions:

Reserved Instruction, DSP Disabled, TLB Refill, TLB Invalid, Bus Error, Address Error, Watch

LHX

Load Halfword Indexed

LWX**Load Word Indexed**

Format: LWX rd, index(base)
LWX rd, rs(rt)

DSP
Replaced with LWX in nanoMIPS

Purpose: Load Word Indexed

To load a word value from memory as a signed value, using indexed addressing.

Description: $rd \leftarrow \text{memory}[\text{base} + \text{index}]$

The contents of GPR *index* is added to the contents of GPR *base* to form an effective address. The contents of the 32-bit word at the memory location specified by the aligned effective address are fetched and placed in GPR *rd*.

Restrictions:

The effective address must be naturally-aligned. If either of the two least-significant bits of the address are non-zero, an Address Error exception occurs.

Operation:

```

ValidateAccessToDSPResources()
vAddr31..0 ← GPR[index]31..0 + GPR[base]31..0
if ( vAddr1..0 ≠ 02 ) then
    SignalException( AddressError )
endif
( pAddr, CCA ) ← AddressTranslation( vAddr, DATA, LOAD )
memwordGPREN..0 ← LoadMemory( CCA, WORD, pAddr, vAddr, DATA )
GPR[rd]31..0 ← memword31..0

```

Exceptions:

Reserved Instruction, DSP Disabled, TLB Refill, TLB Invalid, Bus Error, Address Error, Watch

LWX

Load Word Indexed

MADD**Multiply Word and Add to Accumulator**

31	26	25	21	20	16	15	14	13	12	11	9	8	6	5	3	2	0
P32A 001000		rt		rs		ac		00		101		010		111		111	
6		5		5		2		2		3		3		3		3	

Format: MADD ac, rs, rt**DSP****Purpose:** Multiply Word and Add to Accumulator

To multiply two 32-bit integer words and add the 64-bit result to the specified accumulator.

Description: $(HI[ac] || LO[ac]) \leftarrow (HI[ac] || LO[ac]) + (rs_{31..0} * rt_{31..0})$ The 32-bit signed integer word in register *rs* is multiplied by the corresponding 32-bit signed integer word in register *rt* to produce a 64-bit result. The 64-bit product is added to the specified 64-bit accumulator.These special registers *HI* and *LO* are specified by the value of *ac*. When *ac*=0, this refers to the original *HI/LO* register pair of the MIPS32 architecture.

In Release 6 of the MIPS Architecture, accumulators are eliminated from MIPS32.

No arithmetic exception occurs under any circumstances.

Restrictions:

None

This instruction does not provide the capability of writing directly to a target GPR.

Operation:

```

if (( ac ≠ 0 ) or (ConfigAR ≥ 2)) then
    ValidateAccessToDSP2Resources()
endif
temp63..0 ← ((GPR[rs]31)32 || GPR[rs]31..0) * ((GPR[rt]31)32 || GPR[rt]31..0)
acc63..0 ← ( HI[ac]31..0 || LO[ac]31..0 ) + temp63..0
( HI[ac]31..0 || LO[ac]31..0 ) ← acc63..32 || acc31..0

```

Exceptions:

Reserved Instruction, DSP Disabled

Implementation Notes:Processors which implement a multiplier array which is not square (for example, 32 x 16), and which therefore has an operation latency which is data dependent, should assume that the shorter operand is in register *rt*.**Programming Notes:**In some processors the integer multiply operation may proceed asynchronously and allow other CPU instructions to execute before it is complete. An attempt to read *LO* or *HI* before the results are written interlocks until the results are ready. Asynchronous execution does not affect the program result, but offers an opportunity for performance improvement by scheduling the multiply so that other instructions can execute in parallel.

Programs that require overflow detection must check for it explicitly.

Where the size of the operands are known, software should place the shorter operand in register *rt*. This may reduce the latency of the instruction on those processors which implement data-dependent instruction latencies.

MADD

Multiply Word and Add to Accumulator

MADDU**Multiply Unsigned Word and Add to Accumulator**

31	26	25	21	20	16	15	14	13	12	11	9	8	6	5	3	2	0
P32A 001000		rt		rs		ac	01	101		010		111		111			
6		5		5		2	2	3		3		3		3			

Format: MADDU ac, rs, rt**DSP****Purpose:** Multiply Unsigned Word and Add to Accumulator

To multiply two 32-bit unsigned integer words and add the 64-bit result to the specified accumulator.

Description: $(HI[ac] || LO[ac]) \leftarrow (HI[ac] || LO[ac]) + (rs_{31..0} * rt_{31..0})$ The 32-bit unsigned integer word in register *rs* is multiplied by the corresponding 32-bit unsigned integer word in register *rt* to produce a 64-bit result. The 64-bit product is added to the specified 64-bit accumulator.These special registers *HI* and *LO* are specified by the value of *ac*. When *ac*=0, this refers to the original *HI/LO* register pair of the MIPS32 architecture.

In Release 6 of the MIPS Architecture, accumulators are eliminated from MIPS32.

No arithmetic exception occurs under any circumstances.

Restrictions:

None

This instruction does not provide the capability of writing directly to a target GPR.

Operation:

```

if (( ac ≠ 0 ) or (ConfigAR ≥ 2)) then
    ValidateAccessToDSP2Resources()
endif
temp64..0 ← (032 || GPR[rs]31..0) * (032 || GPR[rt]31..0)
acc63..0 ← (HI[ac]31..0 || LO[ac]31..0) + temp63..0
(HI[ac]31..0 || LO[ac]31..0) ← acc63..32 || acc31..0

```

Exceptions:

Reserved Instruction, DSP Disabled

Implementation Notes:Processors which implement a multiplier array which is not square (for example, 32 x 16), and which therefore has an operation latency which is data dependent, should assume that the shorter operand is in register *rt*.**Programming Notes:**In some processors the integer multiply operation may proceed asynchronously and allow other CPU instructions to execute before it is complete. An attempt to read *LO* or *HI* before the results are written interlocks until the results are ready. Asynchronous execution does not affect the program result, but offers an opportunity for performance improvement by scheduling the multiply so that other instructions can execute in parallel.

Programs that require overflow detection must check for it explicitly.

Where the size of the operands are known, software should place the shorter operand in register *rt*. This may reduce the latency of the instruction on those processors which implement data-dependent instruction latencies.

MADDU

Multiply Unsigned Word and Add to Accumulator

MAQ_S[A].W.PHL

Multiply with Accumulate Single Vector Fractional Halfword Element

31	26	25	21	20	16	15	14	13	12	11	9	8	6	5	3	2	0
MAQ_S.W.PHL																	
P32A 001000	rt			rs			ac	0	1	101	001	111	111				
MAQ_SA.W.PHL																	
P32A 001000	rt			rs			ac	1	1	101	001	111	111				
6	5			5			2	1	1	3	3	3	3				

Format: MAQ_S[A].W.PHL
MAQ_S.W.PHL ac, rs, rt
MAQ_SA.W.PHL ac, rs, rt

DSP
DSP

Purpose: Multiply with Accumulate Single Vector Fractional Halfword Element

To multiply one pair of elements from two vectors of fractional halfword values using full-sized intermediate products and accumulate the result into the specified 64-bit accumulator, with optional saturating accumulation.

Description: $ac \leftarrow \text{sat32}(ac + \text{sat32}(rs_{31..16} * rt_{31..16}))$

The left-most Q15 fractional halfword values from the paired halfword vectors in each of registers *rt* and *rs* are multiplied together, and the product left-shifted by one bit position to generate a Q31 fractional format intermediate result. If both multiplicands are equal to -1.0 in Q15 fractional format (0x8000 hexadecimal), the intermediate result is saturated to the maximum positive Q31 fractional value (0x7FFFFFFF hexadecimal). The intermediate result is then sign-extended and accumulated into accumulator *ac* to generate a 64-bit Q32.31 fractional format result.

In the saturating accumulation variant of this instruction, if the accumulation of the intermediate product with the accumulator results in a value that cannot be represented as a Q31 fractional format value, the accumulator is saturated to either the maximum positive Q31 fractional format value (0x7FFFFFFF hexadecimal) or the minimum negative Q31 fractional format value (0x80000000), sign-extended to 64 bits.

The value of *ac* can range from 0 to 3; a value of 0 refers to the original *HI/LO* register pair of the MIPS32 architecture.

If overflow or saturation occurs, a 1 is written to one of bits 16 through 19 of the *DSPControl* register, within the *ouflag* field. The value of *ac* determines which of these bits is set: bit 16 corresponds to *ac0*, bit 17 to *ac1*, bit 18 to *ac2*, and bit 19 to *ac3*.

Restrictions:

No data-dependent exceptions are possible.

The operands must be values in the specified format. If they are not, the result is **UNPREDICTABLE** and the values of the operand vectors become **UNPREDICTABLE**.

Operation:

```

MAQ_S.W.PHL
    ValidateAccessToDSPResources()
    tempA31..0 ← multiplyQ15Q15( ac, GPR[rs]31..16, GPR[rt]31..16 )
    tempB63..0 ← ( HI[ac]31..0 || LO[ac]31..0 ) + ( (tempA31)32 || tempA31..0 )
    ( HI[ac]31..0 || LO[ac]31..0 ) ← tempB63..32 || tempB31..0

MAQ_SA.W.PHL
    ValidateAccessToDSPResources()
    tempA31..0 ← multiplyQ15Q15( ac, GPR[rs]31..16, GPR[rt]31..16 )
    tempA31..0 ← sat32AccumulateQ31( ac, temp )
    tempB63..0 ← (tempA31)32 || tempA31..0
    ( HI[ac]31..0 || LO[ac]31..0 ) ← tempB63..32 || tempB31..0

```

MAQ_SA.W.PHL**Multiply with Accumulate Single Vector Fractional Halfword Element**

```

function sat32AccumulateQ31( acc1..0, a31..0 )
    signA ← a31
    temp63..0 ← HI[acc]31..0 || LO[acc]31..0
    temp63..0 ← temp + ( (signA)32 || a31..0 )
    if ( temp32 ≠ temp31 ) then
        if ( temp32 = 0 ) then
            temp31..0 ← 0x80000000
        else
            temp31..0 ← 0x7FFFFFFF
        endif
        DSPControlouflag:16+acc ← 1
    endif
    return temp31..0
endfunction sat32AccumulateQ31

```

Exceptions:

Reserved Instruction, DSP Disabled

Programming Notes:

The MAQ_SA version of the instruction is useful for compliance with some ITU speech processing codecs that require a 32-bit saturation after every multiply-accumulate operation.

MAQ_S[A].W.PHR**Multiply with Accumulate Single Vector Fractional Halfword Element**

31	26	25	21	20	16	15	14	13	12	11	9	8	6	5	3	2	0
MAQ_S.W.PHR																	
P32A 001000	rt		rs		ac	0	0	101		001		111		111			
MAQ_SA.W.PHR																	
P32A 001000	rt		rs		ac	1	0	101		001		111		111			
6	5		5		2	1	1	3		3		3		3			

Format: MAQ_S[A].W.PHR
MAQ_S.W.PHR ac, rs, rt
MAQ_SA.W.PHR ac, rs, rt

DSP
DSP

Purpose: Multiply with Accumulate Single Vector Fractional Halfword Element

To multiply one pair of elements from two vectors of fractional halfword values using full-sized intermediate products and accumulate the result into the specified 64-bit accumulator, with optional saturating accumulation.

Description: $ac \leftarrow \text{sat32}(ac + \text{sat32}(rs_{15..0} * rt_{15..0}))$

The right-most Q15 fractional halfword values from each of the registers *rt* and *rs* are multiplied together and the product left-shifted by one bit position to generate a Q31 fractional format intermediate result. If both multiplicands are equal to -1.0 in Q15 fractional format (0x8000 hexadecimal), the intermediate result is saturated to the maximum positive Q31 fractional value (0x7FFFFFFF hexadecimal). The intermediate result is then sign-extended and accumulated into accumulator *ac* to generate a 64-bit Q32.31 fractional format result.

In the saturating accumulation variant of this instruction, if the accumulation of the intermediate product with the accumulator results in a value that cannot be represented as a Q31 fractional format value, the accumulator is saturated to either the maximum positive Q31 fractional format value (0x7FFFFFFF hexadecimal) or the minimum negative Q31 fractional format value (0x80000000), sign-extended to 64 bits.

The value of *ac* can range from 0 to 3; a value of 0 refers to the original *HI/LO* register pair of the MIPS32 architecture.

If overflow or saturation occurs, a 1 is written to one of bits 16 through 19 of the *DSPControl* register, within the *ouflag* field. The value of *ac* determines which of these bits is set: bit 16 corresponds to *ac0*, bit 17 to *ac1*, bit 18 to *ac2*, and bit 19 to *ac3*.

Restrictions:

No data-dependent exceptions are possible.

The operands must be values in the specified format. If they are not, the result is **UNPREDICTABLE** and the values of the operand vectors become **UNPREDICTABLE**.

Operation:

```

MAQ_S.W.PHR
    ValidateAccessToDSPResources()
    tempA31..0 ← multiplyQ15Q15( ac, GPR[rs]15..0, GPR[rt]15..0 )
    tempB63..0 ← ( HI[ac]31..0 || LO[ac]31..0 ) + ( (tempA31)32 || tempA31..0 )
    ( HI[ac]31..0 || LO[ac]31..0 ) ← tempB63..32 || tempB31..0

MAQ_SA.W.PHR
    ValidateAccessToDSPResources()
    tempA31..0 ← multiplyQ15Q15( ac, GPR[rs]15..0, GPR[rt]15..0 )
    tempA31..0 ← sat32AccumulateQ31( ac, temp )
    tempB63..0 ← (tempA31)32 || tempA31..0
    ( HI[ac]31..0 || LO[ac]31..0 ) ← tempB63..32 || tempB31..0

```


MAQ_S[A].W.PHR

Multiply with Accumulate Single Vector Fractional Halfword Element

Exceptions:

Reserved Instruction, DSP Disabled

Programming Notes:

The MAQ_SA version of the instruction is useful for compliance with some ITU speech processing codecs that require a 32-bit saturation after every multiply-accumulate operation.

MFHI**Move from HI register**

31	26	25	21	20	16	15	14	13	12	11	9	8	6	5	3	2	0
P32A 001000						rt		x		ac	00	000	001	111	111		
6						5		5		2	2	3	3	3	3		

Format: MFHI rs, ac**DSP****Purpose:** Move from HI registerTo copy the special purpose *HI* register to a GPR.**Description:** $rs \leftarrow HI[ac]$

The *HI* part of accumulator *ac* is copied to the general-purpose register *rs*. The *HI* part of the accumulator is defined to be bits 32 through 63 of the DSP Module accumulator register.

The value of *ac* can range from 0 to 3. When *ac*=0, this refers to the original *HI/LO* register pair of the MIPS32 architecture.

In Release 6 of the MIPS Architecture, accumulators are eliminated from MIPS32.

Restrictions:

No data-dependent exceptions are possible.

The operands must be values in the specified format. If they are not, the results are **UNPREDICTABLE** and the values of the operand vectors become **UNPREDICTABLE**.

Operation:

```

if (( ac ≠ 0 ) or (ConfigAR ≥ 2)) then
    ValidateAccessToDSPResources()
endif
GPR[rs]31..0 ← HI[ac]31..0

```

Exceptions:

Reserved Instruction, DSP Disabled

MFHI

Move from HI register

MFLO**Move from LO register**

31	26	25	21	20	16	15	14	13	12	11	9	8	6	5	3	2	0
P32A 001000						rt			x			ac	01	000	001	111	111
6						5			5			2	2	3	3	3	3

Format: MFLO rt, ac**DSP****Purpose:** Move from LO registerTo copy the special purpose *LO* register to a GPR.**Description:** $rt \leftarrow LO[ac]$

The *LO* part of accumulator *ac* is copied to the general-purpose register *rt*. The *LO* part of the accumulator is defined to be bits 0 through 31 of the DSP Module accumulator register.

The value of *ac* can range from 0 to 3. When *ac*=0, this refers to the original *H/LO* register pair of the MIPS32 architecture.

In Release 6 of the MIPS Architecture, accumulators are eliminated from MIPS32.

Restrictions:

No data-dependent exceptions are possible.

The operands must be values in the specified format. If they are not, the results are **UNPREDICTABLE** and the values of the operand vectors become **UNPREDICTABLE**.

Operation:

```

if (( ac ≠ 0 ) or (ConfigAR ≥ 2)) then
    ValidateAccessToDSPResources()
endif
GPR[rt]31..0 ← LO[ac]31..0

```

Exceptions:

Reserved Instruction, DSP Disabled

MFLO

Move from LO register

MODSUB**Modular Subtraction on an Index Value**

31	26	25	21	20	16	15	11	10	9	3	2	0
P32A 001000	rt		rs		rd		x	1010010			101	
6	5		5		5		1	7			3	

Format: MODSUB rd, rs, rt**DSP****Purpose:** Modular Subtraction on an Index Value

Do a modular subtraction on a specified index value, using the specified decrement and modular roll-around values.

Description: $rd \leftarrow (GPR[rs] == 0 ? \text{zero_extend}(GPR[rt]_{23..8}) : GPR[rs] - GPR[rt]_{7..0})$

The 32-bit value in register *rs* is compared to the value zero. If it is zero, then the index value has reached the bottom of the buffer and must be rolled back around to the top of the buffer. The index value of the top element of the buffer is obtained from bits 8 through 23 in register *rt*; this value is zero-extended to 32 bits and written to destination register *rd*.

If the value of register *rs* is not zero, then it is simply decremented by the size of the elements in the buffer. The size of the elements, in bytes, is specified by bits 0 through 7 of register *rt*, interpreted as an unsigned integer.

This instruction does not modify the *ouflag* field in the *DSPControl* register.

Restrictions:

No data-dependent exceptions are possible.

The operands must be values in the specified format. If they are not, the results are **UNPREDICTABLE** and the values of the operand vectors become **UNPREDICTABLE**.

Operation:

```

ValidateAccessToDSPResources()
decr7..0 ← GPR[rt]7..0
lastindex15..0 ← GPR[rt]23..8
if ( GPR[rs]31..0 = 0x00000000 ) then
    GPR[rd]31..0 ← 0(GPRLEN-16) || lastindex15..0
else
    GPR[rd]31..0 ← GPR[rs]31..0 - decr7..0
endif

```

Exceptions:

Reserved Instruction, DSP Disabled

MODSUB

Modular Subtraction on an Index Value

MSUB**Multiply Word and Subtract from Accumulator**

31	26	25	21	20	16	15	14	13	12	11	9	8	6	5	3	2	0
P32A 001000		rt		rs		ac		10	101		010		111		111		
6		5		5		2		2	3		3		3		3		

Format: MSUB ac, rs, rt**DSP****Purpose:** Multiply Word and Subtract from Accumulator

To multiply two 32-bit integer words and subtract the 64-bit result from the specified accumulator.

Description: $(HI[ac] || LO[ac]) \leftarrow (HI[ac] || LO[ac]) - (rs_{31..0} * rt_{31..0})$ The 32-bit signed integer word in register *rs* is multiplied by the corresponding 32-bit signed integer word in register *rt* to produce a 64-bit result. The 64-bit product is subtracted from the specified 64-bit accumulator.These special registers *HI* and *LO* are specified by the value of *ac*. When *ac*=0, this refers to the original *HI/LO* register pair of the MIPS32 architecture.

In Release 6 of the MIPS Architecture, accumulators are eliminated from MIPS32.

No arithmetic exception occurs under any circumstances.

Restrictions:

None

This instruction does not provide the capability of writing directly to a target GPR.

Operation:

```

if (( ac ≠ 0 ) or (ConfigAR ≥ 2)) then
    ValidateAccessToDSP2Resources()
endif
temp63..0 ← ((GPR[rs]31)32 || GPR[rs]31..0) * ((GPR[rt]31)32 || GPR[rt]31..0)
acc63..0 ← ( HI[ac]31..0 || LO[ac]31..0 ) - temp63..0
( HI[ac]31..0 || LO[ac]31..0 ) ← acc63..32 || acc31..0

```

Exceptions:

Reserved Instruction, DSP Disabled

Implementation Notes:Processors which implement a multiplier array which is not square (for example, 32 x 16), and which therefore has an operation latency which is data dependent, should assume that the shorter operand is in register *rt*.**Programming Notes:**In some processors the integer multiply operation may proceed asynchronously and allow other CPU instructions to execute before it is complete. An attempt to read *LO* or *HI* before the results are written interlocks until the results are ready. Asynchronous execution does not affect the program result, but offers an opportunity for performance improvement by scheduling the multiply so that other instructions can execute in parallel.

Programs that require overflow detection must check for it explicitly.

Where the size of the operands are known, software should place the shorter operand in register *rt*. This may reduce the latency of the instruction on those processors which implement data-dependent instruction latencies.

MSUB

Multiply Word and Subtract from Accumulator

MSUBU**Multiply Unsigned Word and Add to Accumulator**

31	26	25	21	20	16	15	14	13	12	11	9	8	6	5	3	2	0
P32A 001000	rt		rs		ac	11	101		010		111		111				
6	5		5		2	2	3		3		3		3		3		

Format: MSUBU *ac*, *rs*, *rt***DSP****Purpose:** Multiply Unsigned Word and Add to Accumulator

To multiply two 32-bit unsigned integer words and subtract the 64-bit result from the specified accumulator.

Description: $(HI[ac] || LO[ac]) \leftarrow (HI[ac] || LO[ac]) - (rs_{31..0} * rt_{31..0})$ The 32-bit unsigned integer word in register *rs* is multiplied by the corresponding 32-bit unsigned integer word in register *rt* to produce a 64-bit result. The 64-bit product is subtracted from the specified 64-bit accumulator.These special registers *HI* and *LO* are specified by the value of *ac*. When *ac*=0, this refers to the original *HI/LO* register pair of the MIPS32 architecture.

In Release 6 of the MIPS Architecture, accumulators are eliminated from MIPS32.

No arithmetic exception occurs under any circumstances.

Restrictions:

None

This instruction does not provide the capability of writing directly to a target GPR.

Operation:

```

if (( ac ≠ 0 ) or (ConfigAR ≥ 2)) then
    ValidateAccessToDSP2Resources()
endif
temp64..0 ← (032 || GPR[rs]31..0) * (032 || GPR[rt]31..0)
acc63..0 ← (HI[ac]31..0 || LO[ac]31..0) - temp63..0
(HI[ac]31..0 || LO[ac]31..0) ← acc63..32 || acc31..0

```

Exceptions:

Reserved Instruction, DSP Disabled

Implementation Notes:Processors which implement a multiplier array which is not square (for example, 32 x 16), and which therefore has an operation latency which is data dependent, should assume that the shorter operand is in register *rt*.**Programming Notes:**In some processors the integer multiply operation may proceed asynchronously and allow other CPU instructions to execute before it is complete. An attempt to read *LO* or *HI* before the results are written interlocks until the results are ready. Asynchronous execution does not affect the program result, but offers an opportunity for performance improvement by scheduling the multiply so that other instructions can execute in parallel.

Programs that require overflow detection must check for it explicitly.

Where the size of the operands are known, software should place the shorter operand in register *rt*. This may reduce the latency of the instruction on those processors which implement data-dependent instruction latencies.

MSUBU

Multiply Unsigned Word and Add to Accumulator

MTHI**Move to HI register**

31	26	25	21	20	16	15	14	13	12	11	9	8	6	5	3	2	0
P32A 001000						x		rs		ac	10	000		001	111		111
6						5		5		2	2	3		3	3		3

Format: MTHI *rs*, *ac***DSP****Purpose:** Move to HI registerTo copy a GPR to the special purpose *HI* part of the specified accumulator register.**Description:** $HI[ac] \leftarrow GPR[rs]$

The source register *rs* is copied to the *HI* part of accumulator *ac*. The *HI* part of the accumulator is defined to be bits 32 to 63 of the DSP Module accumulator register.

The value of *ac* can range from 0 to 3. When *ac*=0, this refers to the original *HI/LO* register pair of the MIPS32 architecture.

In Release 6 of the MIPS Architecture, accumulators are eliminated from MIPS32.

Restrictions:

A computed result written to the *HI/LO* pair by DIV, DIVU, DDIV, DDIVU, DMULT, DMULTU, MULT, or MULTU must be read by MFHI or MFLO before a new result can be written into either HI or LO. Note that this restriction only applies to the original *HI/LO* accumulator pair, and does not apply to the new accumulators, *ac1*, *ac2*, and *ac3*.

If an MTHI instruction is executed following one of these arithmetic instructions, but before an MFLO or MFHI instruction, the contents of *LO* are **UNPREDICTABLE**. The following example shows this illegal situation:

```

MULT r2,r4  # start operation that will eventually write to HI,LO
...        # code not containing mfhi or mflo
MTHI r6
...        # code not containing mflo
MFLO r3     # this mflo would get an UNPREDICTABLE value

```

Operation:

```

if (( ac ≠ 0 ) or (ConfigAR ≥ 2)) then
    ValidateAccessToDSPResources()
endif
HI[ac]31..0 ← GPR[rs]31..0

```

Exceptions:

Reserved Instruction, DSP Disabled

MTHI

Move to HI register

MTHLIP**Copy LO to HI and a GPR to LO and Increment Pos by 32**

31	26	25	21	20	16	15	14	13	12	11	9	8	6	5	3	2	0
P32A 001000	x		rs		ac	00	001	001	111	111							
6	5		5		2	2	3	3	3	3							

Format: MTHLIP *rs*, *ac***DSP****Purpose:** Copy LO to HI and a GPR to LO and Increment Pos by 32

Copy the LO part of an accumulator to the HI part, copy a GPR to LO, and increment the *pos* field in the *DSPControl* register by 32.

Description: $ac \leftarrow LO[ac]_{31..0} \parallel GPR[rs]_{31..0} ; DSPControl_{pos:5..0} += 32$

The 32 least-significant bits of the specified accumulator are copied to the most-significant 32 bits of the same accumulator. Then the 32 least-significant bits of register *rs* are copied to the least-significant 32 bits of the accumulator. The instruction then increments the value of bits 0 through 5 of the *DSPControl* register (the *pos* field) by 32.

The result of this instruction is **UNPREDICTABLE** if the value of the *pos* field before the execution of the instruction is greater than 32.

The value of *ac* can range from 0 to 3. When *ac*=0, this refers to the original *HI/LO* register pair of the MIPS32 architecture.

Restrictions:

No data-dependent exceptions are possible.

The operands must be values in the specified format. If they are not, the results are **UNPREDICTABLE** and the values of the operand vectors become **UNPREDICTABLE**.

Operation:

```

ValidateAccessToDSPResources()
tempA31..0 ← GPR[rs]31..0
tempB31..0 ← LO[ac]31..0
( HI[ac]31..0 || LO[ac]31..0 ) ← tempB31..0 || tempA31..0
oldpos5..0 ← DSPControlpos:5..0
if ( oldpos5..0 > 32 ) then
    DSPControlpos:5..0 ← UNPREDICTABLE
else
    DSPControlpos:5..0 ← oldpos5..0 + 32
endif

```

Exceptions:

Reserved Instruction, DSP Disabled

MTHLIP

Copy LO to HI and a GPR to LO and Increment Pos by 32

MTLO**Move to LO register**

31	26	25	21	20	16	15	14	13	12	11	9	8	6	5	3	2	0
P32A 001000						x		rs		ac	11	000	001	111	111		
6						5		5		2	2	3	3	3	3		

Format: MTLO rs, ac**DSP****Purpose:** Move to LO registerTo copy a GPR to the special purpose *LO* part of the specified accumulator register.**Description:** $LO[ac] \leftarrow GPR[rs]$

The source register *rs* is copied to the *LO* part of accumulator *ac*. The *LO* part of the accumulator is defined to be bits 0 to 32 of the DSP Module accumulator register.

The value of *ac* can range from 0 to 3. When *ac*=0, this refers to the original *HI/LO* register pair of the MIPS32 architecture.

In Release 6 of the MIPS Architecture, accumulators are eliminated from MIPS32.

Restrictions:

A computed result written to the *HI/LO* pair by DIV, DIVU, DDIV, DDIVU, DMULT, DMULTU, MULT, or MULTU must be read by MFHI or MFLO before a new result can be written into either HI or LO. Note that this restriction only applies to the original *HI/LO* accumulator pair, and does not apply to the new accumulators, *ac1*, *ac2*, and *ac3*.

If an MTHI instruction is executed following one of these arithmetic instructions, but before an MFLO or MFHI instruction, the contents of *LO* are **UNPREDICTABLE**. The following example shows this illegal situation:

```

MULT r2,r4 # start operation that will eventually write to HI,LO
...       # code not containing mfhi or mflo
MTHI r6
...       # code not containing mflo
MFLO r3   # this mflo would get an UNPREDICTABLE value

```

Operation:

```

if (( ac ≠ 0 ) or (ConfigAR ≥ 2)) then
    ValidateAccessToDSPResources()
endif
LO[ac]31..0 ← GPR[rs]31..0

```

Exceptions:

Reserved Instruction, DSP Disabled

MTLO

Move to LO register

MUL[_S].PH**Multiply Vector Integer HalfWords to Same Size Products**

31	26	25	21	20	16	15	11	10	9	3	2	0
MUL.PH												
P32A 001000	rt		rs		rd		0	0000101			101	
MUL_S.PH												
P32A 001000	rt		rs		rd		1	0000101			101	
6	5		5		5		1	7			3	

Format: MUL[_S].PH

MUL.PH rd, rs, rt

DSP-R2

MUL_S.PH rd, rs, rt

DSP-R2**Purpose:** Multiply Vector Integer HalfWords to Same Size Products

Multiply two vector halfword values.

Description: $rd \leftarrow (rs_{31..16} * rt_{31..16}) \parallel (rs_{15..0} * rt_{15..0})$

Each of the two integer halfword elements in register *rs* is multiplied by the corresponding integer halfword element in register *rt* to create a 32-bit signed integer intermediate result.

In the non-saturation version of the instruction, the 16 least-significant bits of each 32-bit intermediate result are written to the corresponding vector element in destination register *rd*.

In the saturating version of the instruction, intermediate results that cannot be represented in 16 bits are clipped to either the maximum positive 16-bit value (0x7FFF hexadecimal) or the minimum negative 16-bit value (0x8000 hexadecimal), depending on the sign of the intermediate result. The saturated results are then written to the destination register.

To stay compliant with the base architecture, this instruction leaves the base *H/LO* pair (accumulator *ac0*) **UNPREDICTABLE** after the operation completes. The other DSP Module accumulators, *ac1*, *ac2*, and *ac3*, are unchanged.

In the saturating instruction variant, if either multiplication results in an overflow or underflow, the instruction writes a 1 to bit 21 in the *ouflag* field in the *DSPControl* register.

Restrictions:

No data-dependent exceptions are possible.

The operands must be a value in the specified format. If they are not, the results are **UNPREDICTABLE** and the values of the operand vectors become **UNPREDICTABLE**.

Operation:

```

MUL.PH
    ValidateAccessToDSPResources()
    tempB31..0 ← MultiplyI16I16( GPR[rs]31..16, GPR[rt]31..16 )
    tempA31..0 ← MultiplyI16I16( GPR[rs]15..0, GPR[rt]15..0 )
    GPR[rd]31..0 ← tempB15..0 || tempA15..0
    HI31..0 ← UNPREDICTABLE
    LO31..0 ← UNPREDICTABLE

MUL_S.PH
    ValidateAccessToDSPResources()
    tempB31..0 ← sat16MultiplyI16I16( GPR[rs]31..16, GPR[rt]31..16 )
    tempA31..0 ← sat16MultiplyI16I16( GPR[rs]15..0, GPR[rt]15..0 )
    GPR[rd]31..0 ← tempB15..0 || tempA15..0
    HI31..0 ← UNPREDICTABLE

```

MUL[_S].PH**Multiply Vector Integer HalfWords to Same Size Products**

```

LO31..0 ← UNPREDICTABLE

function MultiplyI16I16( a15..0, b15..0 )
    temp31..0 ← a15..0 * b15..0
    if ( temp31..0 > 0x7FFF ) or ( temp31..0 < 0xFFFF8000 ) then
        DSPControlouflag:21 ← 1
    endif
    return temp15..0
endfunction MultiplyI16I16

function satMultiplyI16I16( a15..0, b15..0 )
    temp31..0 ← a15..0 * b15..0
    if ( temp31..0 > 0x7FFF ) then
        temp31..0 ← 0x00007FFF
        DSPControlouflag:21 ← 1
    else
        if ( temp31..0 < 0xFFFF8000 ) then
            temp31..0 ← 0xFFFF8000
            DSPControlouflag:21 ← 1
        endif
    endif
    return temp15..0
endfunction satMultiplyI16I16

```

Exceptions:

Reserved Instruction, DSP Disabled

Programming Notes:

The base MIPS32 architecture states that upon the after a GPR-targeting multiply instruction such as MUL, the contents of *HI* and *LO* are **UNPREDICTABLE**. To stay compliant with the base architecture, this multiply instruction states the same requirement. But this requirement does not apply to the new accumulators *ac1-ac3* and hence a programmer must save the value in *ac0* (which is the same as *HI* and *LO*) across a GPR-targeting multiply instruction, it needed, while the values in *ac1-ac3* do not need to be saved.

MULEQ_S.W.PHL**Multiply Vector Fractional Left Halfwords to Expanded Width Products**

31	26	25	21	20	16	15	11	10	9	3	2	0
P32A 001000	rt		rs		rd		x	0000100			101	
6	5		5		5		1	7			3	

Format: MULEQ_S.W.PHL rd, rs, rt**DSP****Purpose:** Multiply Vector Fractional Left Halfwords to Expanded Width Products

Multiply two Q15 fractional halfword values to produce a Q31 fractional word result, with saturation.

Description: $rd \leftarrow \text{sat32}(rs_{31..16} * rt_{31..16})$

The left-most Q15 fractional halfword value from the paired halfword vector in register *rs* is multiplied by the corresponding Q15 fractional halfword value from register *rt*. The result is left-shifted one bit position to create a Q31 format result and written into the destination register *rd*. If both input values are -1.0 in Q15 format (0x8000 in hexadecimal) the result is clamped to the maximum positive Q31 fractional value (0x7FFFFFFF in hexadecimal) before being written to the destination register.

To stay compliant with the base architecture, this instruction leaves the base *HI/LO* pair (accumulator *ac0*) **UNPREDICTABLE** after the operation completes. The other DSP Module accumulators, *ac1*, *ac2*, and *ac3* are unmodified.

If the result is saturated, this instruction writes a 1 to bit 21 in the *ouflag* field of the *DSPControl* register.

Restrictions:

No data-dependent exceptions are possible.

The operands must be values in the specified format. If they are not, the results are **UNPREDICTABLE** and the values of the operand vectors become **UNPREDICTABLE**.

Operation:

```

ValidateAccessToDSPResources()
temp31..0 ← multiplyQ15Q15ouflag21( GPR[rs]31..16, GPR[rt]31..16 )
GPR[rd]31..0 ← temp31..0
HI[0]31..0 ← UNPREDICTABLE
LO[0]31..0 ← UNPREDICTABLE

function multiplyQ15Q15ouflag21( a15..0, b15..0 )
    if ( a15..0 = 0x8000 ) and ( b15..0 = 0x8000 ) then
        temp31..0 ← 0x7FFFFFFF
        DSPControlouflag:21 ← 1
    else
        temp31..0 ← ( a15..0 * b15..0 ) << 1
    endif
    return temp31..0
endfunction multiplyQ15Q15ouflag21

```

Exceptions:

Reserved Instruction, DSP Disabled

Programming Notes:

The base MIPS32 architecture states that after a GPR-targeting multiply instruction such as MUL, the contents of registers *HI* and *LO* are **UNPREDICTABLE**. To maintain compliance with the base architecture this multiply instruction, MULEQ_S.W.PHL, has the same requirement. Software must save and restore the *ac0* register if the previous value in the *ac0* register is needed following the MULEQ_S.W.PHL instruction.

Note that the requirement on *HI* and *LO* does not apply to the new accumulator registers *ac1*, *ac2*, and *ac3*; as a result

MULEQ_S.W.PHL

Multiply Vector Fractional Left Halfwords to Expanded Width Products

the values in these accumulators need not be saved.

MULEQ_S.W.PHR**Multiply Vector Fractional Right Halfwords to Expanded Width Products**

31	26	25	21	20	16	15	11	10	9	3	2	0
P32A 001000	rt		rs		rd		x	0001100			101	
6	5		5		5		1	7			3	

Format: MULEQ_S.W.PHR rd, rs, rt**DSP****Purpose:** Multiply Vector Fractional Right Halfwords to Expanded Width Products

Multiply two Q15 fractional halfword values to produce a Q31 fractional word result, with saturation.

Description: $rd \leftarrow \text{sat32}(rs_{15..0} * rt_{15..0})$

The right-most Q15 fractional halfword value from register *rs* is multiplied by the corresponding Q15 fractional halfword value from register *rt*. The result is left-shifted one bit position to create a Q31 format result and written into the destination register *rd*. If both input values are -1.0 in Q15 format (0x8000 in hexadecimal) the result is clamped to the maximum positive Q31 fractional value (0x7FFFFFFF in hexadecimal) before being written to the destination register.

To stay compliant with the base architecture, this instruction leaves the base *HI/LO* pair (accumulator *ac0*) **UNPREDICTABLE** after the operation completes. The other DSP Module accumulators, *ac1*, *ac2*, and *ac3* are unmodified.

If the result is saturated, this instruction writes a 1 to bit 21 in the *ouflag* field of the *DSPControl* register.

Restrictions:

No data-dependent exceptions are possible.

The operands must be values in the specified format. If they are not, the results are **UNPREDICTABLE** and the values of the operand vectors become **UNPREDICTABLE**.

Operation:

```

ValidateAccessToDSPResources()
temp31..0 ← multiplyQ15Q15ouflag21( GPR[rs]15..0, GPR[rt]15..0 )
GPR[rd]31..0 ← temp31..0
HI[0]31..0 ← UNPREDICTABLE
LO[0]31..0 ← UNPREDICTABLE

function multiplyQ15Q15ouflag21( a15..0, b15..0 )
    if ( a15..0 = 0x8000 ) and ( b15..0 = 0x8000 ) then
        temp31..0 ← 0x7FFFFFFF
        DSPControlouflag:21 ← 1
    else
        temp31..0 ← ( a15..0 * b15..0 ) << 1
    endif
    return temp31..0
endfunction multiplyQ15Q15ouflag21

```

Exceptions:

Reserved Instruction, DSP Disabled

Programming Notes:

The base MIPS32 architecture states that after a GPR-targeting multiply instruction such as MUL, the contents of registers *HI* and *LO* are **UNPREDICTABLE**. To maintain compliance with the base architecture this multiply instruction, MULEQ_S.W.PHR, has the same requirement. Software must save and restore the *ac0* register if the previous value in the *ac0* register is needed following the MULEQ_S.W.PHR instruction.

MULEQ_S.W.PHR

Multiply Vector Fractional Right Halfwords to Expanded Width Products

Note that the requirement on *HI* and *LO* does not apply to the new accumulator registers *ac1*, *ac2*, and *ac3*; as a result the values in these accumulators need not be saved.

MULEU_S.PH.QBL**Multiply Unsigned Vector Left Bytes by Halfwords to Halfword Products**

31	26	25	21	20	16	15	11	10	9	3	2	0
P32A 001000	rt		rs		rd		x	0010010			101	
6	5		5		5		1	7			3	

Format: MULEU_S.PH.QBL rd, rs, rt**DSP****Purpose:** Multiply Unsigned Vector Left Bytes by Halfwords to Halfword Products

Multiply two left-most unsigned byte vector elements in a byte vector by two unsigned halfword vector elements to produce two unsigned halfword results, with saturation.

Description: $rd \leftarrow \text{sat16}(rs_{31..24} * rt_{31..16}) \mid \mid \text{sat16}(rs_{23..16} * rt_{15..0})$

The two left-most unsigned byte elements in a four-element byte vector in register *rs* are multiplied as unsigned integer values with the four corresponding unsigned halfword elements from register *rt*. The eight most-significant bits of each 24-bit result are discarded, and the remaining 16 least-significant bits are written to the corresponding elements in halfword vector register *rd*. The instruction saturates the result to the maximum positive value (0xFFFF hexadecimal) if any of the discarded bits from each intermediate result are non-zero.

To stay compliant with the base architecture, this instruction leaves the base *HI/LO* pair (accumulator *ac0*) **UNPREDICTABLE** after the operation completes. The other DSP Module accumulators, *ac1*, *ac2*, and *ac3* are unmodified.

If either result is saturated this instruction writes a 1 to bit 21 in the *DSPControl* register in the *ouflag* field.

Restrictions:

No data-dependent exceptions are possible.

The operands must be values in the specified format. If they are not, the results are **UNPREDICTABLE** and the values of the operand vectors become **UNPREDICTABLE**.

Operation:

```

ValidateAccessToDSPResources()
tempB15..0 ← multiplyU8U16( GPR[rs]31..24, GPR[rt]31..16 )
tempA15..0 ← multiplyU8U16( GPR[rs]23..16, GPR[rt]15..0 )
GPR[rd]31..0 ← tempB15..0 || tempA15..0
HI[0]31..0 ← UNPREDICTABLE
LO[0]31..0 ← UNPREDICTABLE

function multiplyU8U16( a7..0, b15..0 )
    temp25..0 ← (0 || a) * (0 || b)
    if ( temp25..16 > 0x00 ) then
        temp25..0 ← 010 || 0xFFFF
        DSPControlouflag:21 ← 1
    endif
    return temp15..0
endfunction multiplyU8U16

```

Exceptions:

Reserved Instruction, DSP Disabled

Programming Notes:

The base MIPS32 architecture states that after a GPR-targeting multiply instruction such as MUL, the contents of registers *HI* and *LO* are **UNPREDICTABLE**. To maintain compliance with the base architecture this multiply instruction, MULEU_S.PH.QBL, has the same requirement. Software must save and restore the *ac0* register if the previous value in the *ac0* register is needed following the MULEU_S.PH.QBL instruction.

MULEU_S.PH.QBL

Multiply Unsigned Vector Left Bytes by Halfwords to Halfword Products

Note that the requirement on *HI* and *LO* does not apply to the new accumulator registers *ac1*, *ac2*, and *ac3*; as a result the values in these accumulators need not be saved.

MULEU_S.PH.QBR**Multiply Unsigned Vector Right Bytes with halfwords to Half Word Products**

31	26	25	21	20	16	15	11	10	9	3	2	0
P32A 001000	rt		rs		rd		x	0011010			101	
6	5		5		5		1	7			3	

Format: MULEU_S.PH.QBR rd, rs, rt**DSP****Purpose:** Multiply Unsigned Vector Right Bytes with halfwords to Half Word Products

Element-wise multiplication of unsigned byte elements with corresponding unsigned halfword elements, with saturation.

Description: $rd \leftarrow \text{sat16}(rs_{15..8} * rt_{31..16}) \mid \mid \text{sat16}(rs_{7..0} * rt_{15..0})$

The two right-most unsigned byte elements in a four-element byte vector in register *rs* are multiplied as unsigned integer values with the corresponding right-most 16-bit unsigned values from register *rt*. Each result is clipped to preserve the 16 least-significant bits and written back into the respective halfword element positions in the destination register *rd*. The instruction saturates the result to the maximum positive value (0xFFFF hexadecimal) if any of the clipped bits are non-zero.

To stay compliant with the base architecture, this instruction leaves the base *HI/LO* pair (accumulator *ac0*) **UNPREDICTABLE** after the operation completes. The other DSP Module accumulators, *ac1*, *ac2*, and *ac3* must be unmodified.

This instruction writes a 1 to bit 21 in the *ouflag* field in the *DSPControl* register if either multiplication results in saturation.

Restrictions:

No data-dependent exceptions are possible.

The operands must be values in the specified format. If they are not, the results are **UNPREDICTABLE** and the values of the operand vectors become **UNPREDICTABLE**.

Operation:

```

ValidateAccessToDSPResources()
tempB15..0 ← multiplyU8U16( GPR[rs]15..8, GPR[rt]31..16 )
tempA15..0 ← multiplyU8U16( GPR[rs]7..0, GPR[rt]15..0 )
GPR[rd] ← tempB15..0 || tempA15..0
HI[0]31..0 ← UNPREDICTABLE
LO[0]31..0 ← UNPREDICTABLE

```

Exceptions:

Reserved Instruction, DSP Disabled

Programming Notes:

The base MIPS32 architecture states that after a GPR-targeting multiply instruction such as MUL, the contents of registers *HI* and *LO* are **UNPREDICTABLE**. To maintain compliance with the base architecture this multiply instruction, MULEU_S.PH.QBR, has the same requirement. Software must save and restore the *ac0* register if the previous value in the *ac0* register is needed following the MULEU_S.PH.QBR instruction.

Note that the requirement on *HI* and *LO* does not apply to the new accumulator registers *ac1*, *ac2*, and *ac3*; as a result the values in these accumulators need not be saved.

MULEU_S.PH.QBR

Multiply Unsigned Vector Right Bytes with halfwords to Half Word Products

MULQ_RS.PH**Multiply Vector Fractional Halfwords to Fractional Halfword Products**

31	26	25	21	20	16	15	11	10	9	3	2	0
P32A 001000	rt		rs		rd		x	0100010			101	
6	5		5		5		1	7			3	

Format: MULQ_RS.PH rd, rs, rt**DSP****Purpose:** Multiply Vector Fractional Halfwords to Fractional Halfword Products

Multiply Q15 fractional halfword vector elements with rounding and saturation to produce two Q15 fractional halfword results.

Description: $rd \leftarrow \text{rndQ15}(rs_{31..16} * rt_{31..16}) \parallel \text{rndQ15}(rs_{15..0} * rt_{15..0})$

The two Q15 fractional halfword elements from register *rs* are separately multiplied by the corresponding Q15 fractional halfword elements from register *rt* to produce 32-bit intermediate results. Each intermediate result is left-shifted by one bit position to produce a Q31 fractional value, then rounded by adding 0x00008000 hexadecimal. The rounded intermediate result is then truncated to a Q15 fractional value and written to the corresponding position in destination register *rd*.

If the two input values to either multiplication are both -1.0 (0x8000 in hexadecimal), the final halfword result is saturated to the maximum positive Q15 value (0x7FFF in hexadecimal) and rounding and truncation are not performed.

To stay compliant with the base architecture, this instruction leaves the base *H/L*O pair (accumulator *ac0*) **UNPREDICTABLE** after the operation completes. The other DSP Module accumulators, *ac1*, *ac2*, and *ac3* must be unmodified.

If either result is saturated this instruction writes a 1 to bit 21 in the *DSPControl* register in the *ouflag* field.

Restrictions:

No data-dependent exceptions are possible.

The operands must be values in the specified format. If they are not, the results are **UNPREDICTABLE** and the values of the operand vectors become **UNPREDICTABLE**.

Operation:

```

ValidateAccessToDSPResources()
tempB15..0 ← rndQ15MultiplyQ15Q15( GPR[rs]31..16, GPR[rt]31..16 )
tempA15..0 ← rndQ15MultiplyQ15Q15( GPR[rs]15..0, GPR[rt]15..0 )
GPR[rd]31..0 ← tempB15..0 || tempA15..0
HI[0]31..0 ← UNPREDICTABLE
LO[0]31..0 ← UNPREDICTABLE

function rndQ15MultiplyQ15Q15( a15..0, b15..0 )
  if ( a15..0 = 0x8000 ) and ( b15..0 = 0x8000 ) then
    temp31..0 ← 0x7FFF0000
    DSPControlouflag:21 ← 1
  else
    temp31..0 ← ( a15..0 * b15..0 ) << 1
    temp31..0 ← temp31..0 + 0x00008000
  endif
  return temp31..16
endfunction rndQ15MultiplyQ15Q15

```

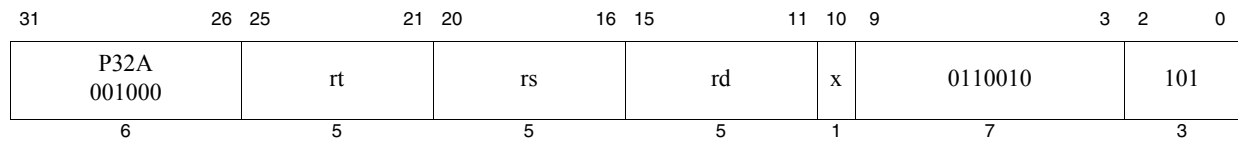
Exceptions:

Reserved Instruction, DSP Disabled

MULQ_RS.PH**Multiply Vector Fractional Halfwords to Fractional Halfword Products****Programming Notes:**

The base MIPS32 architecture states that after a GPR-targeting multiply instruction such as `MUL`, the contents of registers *HI* and *LO* are **UNPREDICTABLE**. To maintain compliance with the base architecture, this multiply instruction, `MULQ_RS.PH`, has the same requirement. Software must save and restore the *ac0* register if the previous value in the *ac0* register is needed following the `MULQ_RS.PH` instruction.

Note that the requirement on *HI* and *LO* does not apply to the new accumulator registers *ac1*, *ac2*, and *ac3*; as a result, the values in these accumulators need not be saved.

MULQ_RS.W**Multiply Fractional Words to Same Size Product with Saturation and Rounding****Format:** MULQ_RS.W rd, rs, rt**DSP-R2****Purpose:** Multiply Fractional Words to Same Size Product with Saturation and Rounding

Multiply fractional Q31 word values, with saturation and rounding.

Description: $rd \leftarrow \text{round}(\text{sat32}(rs_{31..0} * rt_{31..0}))$

The Q31 fractional format words in registers *rs* and *rt* are multiplied together and the product shifted left by one bit position to create a 64-bit fractional format intermediate result. The intermediate result is rounded up by adding a 1 at bit position 31, and then truncated by discarding the 32 least-significant bits to create a 32-bit fractional format result. The result is then written to destination register *rd*.

If both input multiplicands are equal to -1 (0x80000000 hexadecimal), rounding is not performed and the maximum positive Q31 fractional format value (0x7FFFFFFF hexadecimal) is written to the destination register.

To stay compliant with the base architecture, this instruction leaves the base *H/L*O pair (accumulator *ac0*) **UNPREDICTABLE** after the operation completes. The other DSP Module accumulators, *ac1*, *ac2*, and *ac3*, are unchanged.

This instruction, on an overflow or underflow of the operation, writes a 1 to bit 21 in the *DSPControl* register in the *ouflag* field.

Restrictions:

No data-dependent exceptions are possible.

The operands must be a value in the specified format. If they are not, the results are **UNPREDICTABLE** and the values of the operand vectors become **UNPREDICTABLE**.

Operation:

```

ValidateAccessToDSP2Resources()
if ( GPR[rs]_31..0 = 0x80000000 ) and ( GPR[rt]_31..0 = 0x80000000 ) then
    temp_63..0 ← 0x7FFFFFFF00000000
    DSPControl_ouflag:21 ← 1
else
    temp_63..0 ← ( GPR[rs]_31..0 * GPR[rt]_31..0 ) << 1
    temp_63..0 ← temp_63..0 + ( 032 || 0x80000000 )
endif
GPR[rd]_31..0 ← temp_63..32
HI[0]_31..0 ← UNPREDICTABLE
LO[0]_31..0 ← UNPREDICTABLE

```

Exceptions:

Reserved Instruction, DSP Disabled

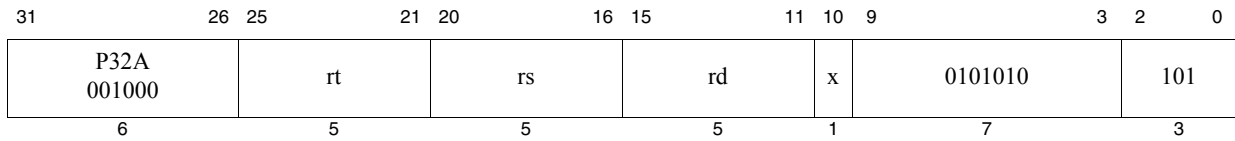
Programming Notes:

The base MIPS32 architecture states that after a GPR-targeting multiply instruction such as MUL, the contents of registers *H* and *L*O are **UNPREDICTABLE**. To maintain compliance with the base architecture, this multiply instruction, MULQ_RS.W, has the same requirement. Software must save and restore the *ac0* register if the previous value in the *ac0* register is needed following the MULQ_RS.W instruction.

Note that the requirement on *H* and *L*O does not apply to the new accumulator registers *ac1*, *ac2*, and *ac3*; as a result, the values in these accumulators need not be saved.

MULQ_RS.W

Multiply Fractional Words to Same Size Product with Saturation and Rounding

MULQ_S.PH**Multiply Vector Fractional Half-Words to Same Size Products****Format:** MULQ_S.PH rd, rs, rt**DSP-R2****Purpose:** Multiply Vector Fractional Half-Words to Same Size Products

Multiply two vector fractional Q15 values to create a Q15 result, with saturation.

Description: $rd \leftarrow \text{sat16}(rs_{31..16} * rt_{31..16}) \parallel \text{sat16}(rs_{15..0} * rt_{15..0})$

The two vector fractional Q15 values in register *rs* are multiplied with the corresponding elements in register *rt* to produce two 32-bit products. Each product is left-shifted by one bit position to create a Q31 fractional word intermediate result. The two 32-bit intermediate results are then each truncated by discarding the 16 least-significant bits of each result, and the resulting Q15 fractional format halfwords are then written to the corresponding positions in destination register *rd*. For each halfword result, if both input multiplicands are equal to -1 (0x8000 hexadecimal), the final halfword result is saturated to the maximum positive Q15 value (0x7FFF hexadecimal).

To stay compliant with the base architecture, this instruction leaves the base *H/L*O pair (accumulator *ac0*) **UNPREDICTABLE** after the operation completes. The other DSP Module accumulators, *ac1*, *ac2*, and *ac3*, must be untouched.

This instruction, on an overflow or underflow of any one of the two vector operation, writes bit 21 in the *ouflag* field in the *DSPControl* register.

Restrictions:

No data-dependent exceptions are possible.

The operands must be a value in the specified format. If they are not, the results are **UNPREDICTABLE** and the values of the operand vectors become **UNPREDICTABLE**.

Operation:

```

ValidateAccessToDSP2Resources()
tempB31..0 ← sat16MultiplyQ15Q15( GPR[rs]31..16, GPR[rt]31..16 )
tempA31..0 ← sat16MultiplyQ15Q15( GPR[rs]15..0, GPR[rt]15..0 )
GPR[rd]31..0 ← tempB15..0 || tempA15..0
HI[0]31..0 ← UNPREDICTABLE
LO[0]31..0 ← UNPREDICTABLE

function sat16MultiplyQ15Q15( a15..0, b15..0 )
  if ( a15..0 = 0x8000 ) and ( b15..0 = 0x8000 ) then
    temp31..0 ← 0x7FFF0000
    DSPControlouflag:21 ← 1
  else
    temp31..0 ← ( a15..0 * b15..0 )
    temp31..0 ← ( temp30..0 || 0 )
  endif
  return temp31..16
endfunction sat16MultiplyQ15Q15

```

Exceptions:

Reserved Instruction, DSP Disabled

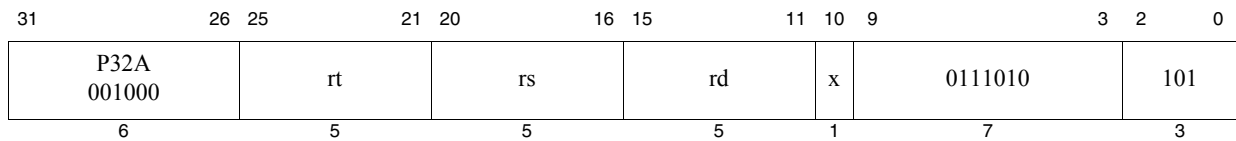
Programming Notes:

The base MIPS32 architecture states that after a GPR-targeting multiply instruction such as MUL, the contents of reg-

MULQ_S.PH**Multiply Vector Fractional Half-Words to Same Size Products**

isters *HI* and *LO* are **UNPREDICTABLE**. To maintain compliance with the base architecture, this multiply instruction, MULQ_S.PH, has the same requirement. Software must save and restore the *ac0* register if the previous value in the *ac0* register is needed following the MULQ_S.PH instruction.

Note that the requirement on *HI* and *LO* does not apply to the new accumulator registers *ac1*, *ac2*, and *ac3*; as a result, the values in these accumulators need not be saved.

MULQ_S.W**Multiply Fractional Words to Same Size Product with Saturation****Format:** MULQ_S.W rd, rs, rt**DSP-R2****Purpose:** Multiply Fractional Words to Same Size Product with Saturation

Multiply two Q31 fractional format word values to create a fractional Q31 result, with saturation.

Description: $rd \leftarrow \text{sat32}(rs_{31..0} * rt_{31..0})$

The Q31 fractional format words in registers *rs* and *rt* are multiplied together to create a 64-bit fractional format intermediate result. The intermediate result is left-shifted by one bit position, and then truncated by discarding the 32 least-significant bits to create a Q31 fractional format result. This result is then written to destination register *rd*.

If both input multiplicands are equal to -1 (0x80000000 hexadecimal), the product is clipped to the maximum positive Q31 fractional format value (0x7FFFFFFF hexadecimal), and written to the destination register.

To stay compliant with the base architecture, this instruction leaves the base *H/L*O pair (accumulator *ac0*) **UNPREDICTABLE** after the operation completes. The other DSP Module accumulators, *ac1*, *ac2*, and *ac3*, are unchanged.

This instruction, on an overflow or underflow of the operation, writes a 1 to bit 21 in the *DSPControl* register in the *ouflag* field.

Restrictions:

No data-dependent exceptions are possible.

The operands must be a value in the specified format. If they are not, the results are **UNPREDICTABLE** and the values of the operand vectors become **UNPREDICTABLE**.

Operation:

```

ValidateAccessToDSP2Resources()
if ( GPR[rs]_31..0 = 0x80000000 ) and ( GPR[rt]_31..0 = 0x80000000 ) then
    temp_63..0 ← 0x7FFFFFFF00000000
    DSPControl_ouflag:21 ← 1
else
    temp_63..0 ← ( GPR[rs]_31..0 * GPR[rt]_31..0 ) << 1
endif
GPR[rd]_31..0 ← temp_63..32
HI[0]_31..0 ← UNPREDICTABLE
LO[0]_31..0 ← UNPREDICTABLE

```

Exceptions:

Reserved Instruction, DSP Disabled

Programming Notes:

The base MIPS32 architecture states that after a GPR-targeting multiply instruction such as MUL, the contents of registers *H* and *L*O are **UNPREDICTABLE**. To maintain compliance with the base architecture, this multiply instruction, MULQ_S.W, has the same requirement. Software must save and restore the *ac0* register if the previous value in the *ac0* register is needed following the MULQ_S.W instruction.

Note that the requirement on *H* and *L*O does not apply to the new accumulator registers *ac1*, *ac2*, and *ac3*; as a result, the values in these accumulators need not be saved.

MULQ_S.W

Multiply Fractional Words to Same Size Product with Saturation

MULSA.W.PH**Multiply and Subtract Vector Integer Halfword Elements and Accumulate**

31	26	25	21	20	16	15	14	13	12	11	9	8	6	5	3	2	0
P32A 001000						rt			rs			ac	10	110	010	111	111
6						5			5			2	2	3	3	3	3

Format: MULSA.W.PH *ac*, *rs*, *rt***DSP-R2****Purpose:** Multiply and Subtract Vector Integer Halfword Elements and Accumulate

To multiply and subtract two integer vector elements using full-size intermediate products, accumulating the result into the specified accumulator.

Description: $ac \leftarrow ac + ((rs_{31..16} * rt_{31..16}) - (rs_{15..0} * rt_{15..0}))$

Each of the two halfword integer elements from register *rt* are multiplied by the corresponding elements in *rs* to create two word results. The right-most result is subtracted from the left-most result to generate the intermediate result, which is then added to the specified 64-bit accumulator.

The value of *ac* selects an accumulator numbered from 0 to 3. When *ac*=0, this refers to the original *HI/LO* register pair of the MIPS32 architecture.

This instruction does not set any bits of the *ouflag* field in the *DSPControl* register.

Restrictions:

No data-dependent exceptions are possible.

The operands must be a value in the specified format. If they are not, the result is **UNPREDICTABLE** and the values of the operand vectors become **UNPREDICTABLE**.

Operation:

```

ValidateAccessToDSP2Resources()
tempB31..0 ← (GPR[rs]31..16 * GPR[rt]31..16)
tempA31..0 ← (GPR[rs]15..0 * GPR[rt]15..0)
dotp32..0 ← ( (tempB31) || tempB31..0 ) - ( (tempA31) || tempA31..0 )
acc63..0 ← ( HI[ac]31..0 || LO[ac]31..0 ) + ( (dotp32)31 || dotp32..0 )
( HI[ac]31..0 || LO[ac]31..0 ) ← acc63..32 || acc31..0

```

Exceptions:

Reserved Instruction, DSP Disabled

MULSA.W.PH

Multiply and Subtract Vector Integer Halfword Elements and Accumulate

MULSAQ_S.W.PH**Multiply And Subtract Vector Fractional Halfwords And Accumulate**

31	26	25	21	20	16	15	14	13	12	11	9	8	6	5	3	2	0
P32A 001000		rt		rs		ac	11	110		010		111		111			
6		5		5		2	2	3		3		3		3			

Format: MULSAQ_S.W.PH *ac*, *rs*, *rt***DSP****Purpose:** Multiply And Subtract Vector Fractional Halfwords And Accumulate

Multiply and subtract two Q15 fractional halfword vector elements using full-size intermediate products, accumulating the result from the specified accumulator, with saturation.

Description: $ac \leftarrow ac + (\text{sat32}(rs_{31..16} * rt_{31..16}) - \text{sat32}(rs_{15..0} * rt_{15..0}))$

The two corresponding Q15 fractional values from registers *rt* and *rs* are multiplied together and left-shifted by 1 bit to generate two Q31 fractional format intermediate products. If the input multiplicands to either of the multiplications are both -1.0 (0x8000 hexadecimal), the intermediate result is saturated to 0x7FFFFFFF hexadecimal.

The two intermediate products (named left and right) are summed with alternating sign to create a sum-of-products, i.e., the sign of the right product is negated before summation. The sum-of-products is then sign-extended to 64 bits and accumulated into the specified 64-bit accumulator, producing a Q32.31 result.

The value of *ac* can range from 0 to 3; a value of 0 refers to the original *HI/LO* register pair of the MIPS32 architecture.

If saturation occurs, a 1 is written to one of bits 16 through 19 of the *DSPControl* register, within the *ouflag* field. The value of *ac* determines which of these bits is set: bit 16 corresponds to *ac0*, bit 17 to *ac1*, bit 18 to *ac2*, and bit 19 to *ac3*.

Restrictions:

No data-dependent exceptions are possible.

The operands must be values in the specified format. If they are not, the results are **UNPREDICTABLE** and the values of the operand vectors become **UNPREDICTABLE**.

Operation:

```

ValidateAccessToDSPResources()
tempB31..0 ← multiplyQ15Q15( ac, rs31..16, rt31..16 )
tempA31..0 ← multiplyQ15Q15( ac, rs15..0, rt15..0 )
dotp63..0 ← ( (tempB31)32 || tempB31..0 ) - ( (tempA31)32 || tempA31..0 )
tempC63..0 ← ( HI[ac]31..0 || LO[ac]31..0 ) + dotp63..0
( HI[ac]31..0 || LO[ac]31..0 ) ← tempC63..32 || tempC31..0

```

Exceptions:

Reserved Instruction, DSP Disabled

MULSAQ_S.W.PH

Multiply And Subtract Vector Fractional Halfwords And Accumulate

MULT**Multiply Word**

31	26	25	21	20	16	15	14	13	12	11	9	8	6	5	3	2	0
P32A 001000						rt			rs			ac	00	110	010	111	111
6						5			5			2	2	3	3	3	3

Format: MULT ac, rs, rt**DSP****Purpose:** Multiply Word

To multiply two 32-bit signed integers, writing the 64-bit result to the specified accumulator.

Description: $ac \leftarrow rs_{31..0} * rt_{31..0}$

The 32-bit signed integer value in register *rt* is multiplied by the corresponding 32-bit signed integer value in register *rs*, to produce a 64-bit result that is written to the specified accumulator register.

The value of *ac* selects an accumulator numbered from 0 to 3. When *ac*=0, this refers to the original *HI/LO* register pair of the MIPS32 architecture.

In Release 6 of the MIPS Architecture, accumulators are eliminated from MIPS32.

No arithmetic exception occurs under any circumstances.

Restrictions:

None

Operation:

```

if (( ac ≠ 0 ) or (ConfigAR ≥ 2)) then
    ValidateAccessToDSP2Resources()
endif
temp63..0 ← ((GPR[rs]31)32 || GPR[rs]31..0) * ((GPR[rt]31)32 || GPR[rt]31..0)
( HI[ac]31..0 || LO[ac]31..0 ) ← temp63..32 || temp31..0

```

Exceptions:

Reserved Instruction, DSP Disabled

Programming Notes:

In some processors the integer multiply operation may proceed asynchronously and allow other CPU instructions to execute before it is complete. An attempt to read *LO* or *HI* before the results are written interlocks until the results are ready. Asynchronous execution does not affect the program result, but offers an opportunity for performance improvement by scheduling the multiply so that other instructions can execute in parallel.

Programs that require overflow detection must check for it explicitly.

Where the size of the operands are known, software should place the shorter operand in register *rt*. This may reduce the latency of the instruction on those processors which implement data-dependent instruction latencies.

Implementation Note:

Processors which implement a multiplier array which is not square (for example, 32 x 16), and which therefore has an operation latency which is data dependent, should assume that the shorter operand is in register *rt*.

MULT

Multiply Word

MULTU**Multiply Unsigned Word**

31	26	25	21	20	16	15	14	13	12	11	9	8	6	5	3	2	0
P32A 001000		rt			rs			ac	01	110		010		111		111	
6		5			5			2	2	3		3		3		3	

Format: MULTU ac, rs, rt**DSP****Purpose:** Multiply Unsigned Word

To multiply 32-bit unsigned integers, writing the 64-bit result to the specified accumulator.

Description: $ac \leftarrow rs_{31..0} * rt_{31..0}$

The 32-bit unsigned integer value in register *rt* is multiplied by the corresponding 32-bit unsigned integer value in register *rs*, to produce a 64-bit unsigned result that is written to the specified accumulator register.

The value of *ac* selects an accumulator numbered from 0 to 3. When *ac*=0, this refers to the original *HI/LO* register pair of the MIPS32 architecture.

In Release 6 of the MIPS Architecture, accumulators are eliminated from MIPS32.

No arithmetic exception occurs under any circumstances.

Restrictions:

None

Operation:

```

if (( ac ≠ 0 ) or (ConfigAR ≥ 2)) then
    ValidateAccessToDSP2Resources()
endif
temp64..0 ← ( 032 || GPR[rs]31..0 ) * ( 032 || GPR[rt]31..0 )
( HI[ac]31..0 || LO[ac]31..0 ) ← temp63..32 || temp31..0

```

Exceptions:

Reserved Instruction, DSP Disabled

Programming Notes:

In some processors the integer multiply operation may proceed asynchronously and allow other CPU instructions to execute before it is complete. An attempt to read *LO* or *HI* before the results are written interlocks until the results are ready. Asynchronous execution does not affect the program result, but offers an opportunity for performance improvement by scheduling the multiply so that other instructions can execute in parallel.

Programs that require overflow detection must check for it explicitly.

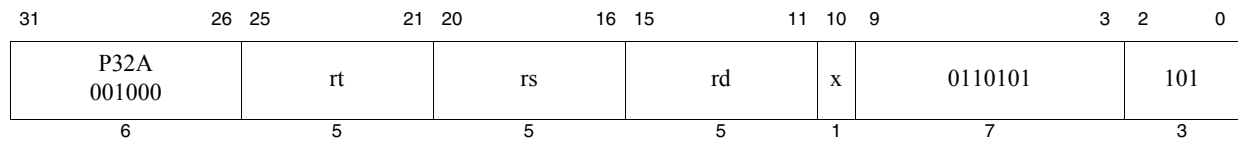
Where the size of the operands are known, software should place the shorter operand in register *rt*. This may reduce the latency of the instruction on those processors which implement data-dependent instruction latencies.

Implementation Note:

Processors which implement a multiplier array which is not square (for example, 32 x 16), and which therefore has an operation latency which is data dependent, should assume that the shorter operand is in register *rt*.

MULTU

Multiply Unsigned Word

PACKRL.PH**Pack a Vector of Halfwords from Vector Halfword Sources****Format:** PACKRL.PH rd, rs, rt**DSP****Purpose:** Pack a Vector of Halfwords from Vector Halfword Sources

Pick two elements for a halfword vector using the right halfword and left halfword respectively from the two source registers.

Description: $rd \leftarrow rs_{15..0} \parallel rt_{31..16}$

The right halfword element from register *rs* and the left halfword from register *rt* are packed into the two halfword positions of the destination register *rd*.

Restrictions:

No data-dependent exceptions are possible.

The operands must be values in the specified format. If they are not, the results are **UNPREDICTABLE** and the values of the operand vectors become **UNPREDICTABLE**.

Operation:

```

ValidateAccessToDSPResources()
tempB15..0 ← GPR[rs]15..0
tempA15..0 ← GPR[rt]31..16
GPR[rd]31..0 ← tempB15..0 || tempA15..0

```

Exceptions:

Reserved Instruction, DSP Disabled

PACKRL.PH

Pack a Vector of Halfwords from Vector Halfword Sources

PICK.PH**Pick a Vector of Halfword Values Based on Condition Code Bits**

31	26	25	21	20	16	15	11	10	9	3	2	0
P32A 001000						rt	rs	rd	x	1000101		101
6						5	5	5	1	7		3

Format: PICK.PH rd, rs, rt**DSP****Purpose:** Pick a Vector of Halfword Values Based on Condition Code Bits

Select two halfword elements from either of two source registers based on condition code bits, writing the selected elements to the destination register.

Description: $rd \leftarrow \text{pick}(cc_{25}, rs_{31..16}, rt_{31..16}) \mid \mid \text{pick}(cc_{24}, rs_{15..0}, rt_{15..0})$

The two right-most condition code bits in the *DSPControl* register are used to select halfword values from the corresponding element of either source register *rs* or source register *rt*. If the value of the corresponding condition code bit is 1, then the halfword value is selected from register *rs*; otherwise, it is selected from *rt*. The selected halfwords are written to the destination register *rd*.

Restrictions:

No data-dependent exceptions are possible.

The operands must be in the specified format. If they are not, the results are **UNPREDICTABLE** and the values of the operand vectors become **UNPREDICTABLE**.

Operation:

```

ValidateAccessToDSPResources()
tempB15..0 ← ( DSPControlccond:25 = 1 ? GPR[rs]31..16 : GPR[rt]31..16 )
tempA15..0 ← ( DSPControlccond:24 = 1 ? GPR[rs]15..0 : GPR[rt]15..0 )
GPR[rd]31..0 ← tempB15..0 || tempA15..0

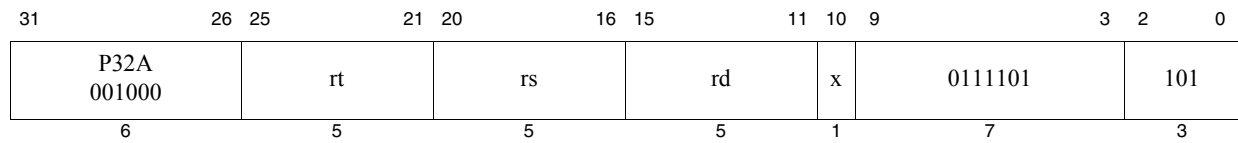
```

Exceptions:

Reserved Instruction, DSP Disabled

PICK.PH

Pick a Vector of Halfword Values Based on Condition Code Bits

PICK.QB**Pick a Vector of Byte Values Based on Condition Code Bits****Format:** PICK.QB rd, rs, rt**DSP****Purpose:** Pick a Vector of Byte Values Based on Condition Code Bits

Select four byte elements from either of two source registers based on condition code bits, writing the selected elements to the destination register.

Description: $rd \leftarrow \text{pick}(cc_{27}, rs_{31..24}, rt_{31..24}) \ || \ \text{pick}(cc_{26}, rs_{23..16}, rt_{23..16}) \ || \ \text{pick}(cc_{25}, rs_{15..8}, rt_{15..8}) \ || \ \text{pick}(cc_{24}, rs_{7..0}, rt_{7..0})$

Four condition code bits in the *DSPControl* register are used to select byte values from the corresponding byte element of either source register *rs* or source register *rt*. If the value of the corresponding condition code bit is 1, then the byte value is selected from register *rs*; otherwise, it is selected from *rt*. The selected bytes are written to the destination register *rd*.

Restrictions:

No data-dependent exceptions are possible.

The operands must be in the specified format. If they are not, the results are **UNPREDICTABLE** and the values of the operand vectors become **UNPREDICTABLE**.

Operation:

```

ValidateAccessToDSPResources()
tempD7..0 ← ( DSPControl_ccond:27 = 1 ? GPR[rs]31..24 : GPR[rt]31..24 )
tempC7..0 ← ( DSPControl_ccond:26 = 1 ? GPR[rs]23..16 : GPR[rt]23..16 )
tempB7..0 ← ( DSPControl_ccond:25 = 1 ? GPR[rs]15..8 : GPR[rt]15..8 )
tempA7..0 ← ( DSPControl_ccond:24 = 1 ? GPR[rs]7..0 : GPR[rt]7..0 )
GPR[rd]31..0 ← tempD7..0 || tempC7..0 || tempB7..0 || tempA7..0

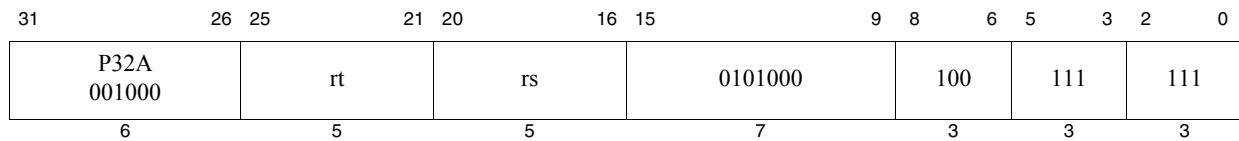
```

Exceptions:

Reserved Instruction, DSP Disabled

PICK.QB

Pick a Vector of Byte Values Based on Condition Code Bits

PRECEQ.W.PHL**Precision Expand Fractional Halfword to Fractional Word Value****Format:** PRECEQ.W.PHL *rt*, *rs***DSP****Purpose:** Precision Expand Fractional Halfword to Fractional Word Value

Expand the precision of a Q15 fractional value taken from the left element of a paired halfword vector to create a Q31 fractional word value.

Description: $rt \leftarrow \text{expand_prec}(rs_{31..16})$

The left Q15 fractional halfword value from the paired halfword vector in register *rs* is expanded to a Q31 fractional value and written to destination register *rt*. The precision expansion is achieved by appending 16 least-significant zero bits to the original halfword value to generate the 32-bit fractional value.

Restrictions:

No data-dependent exceptions are possible.

The operands must be values in the specified format. If they are not, the results are **UNPREDICTABLE** and the values of the operand vectors become **UNPREDICTABLE**.

Operation:

```

ValidateAccessToDSPResources()
temp31..0 ← GPR[rs]31..16 || 016
GPR[rt]31..0 ← temp31..0

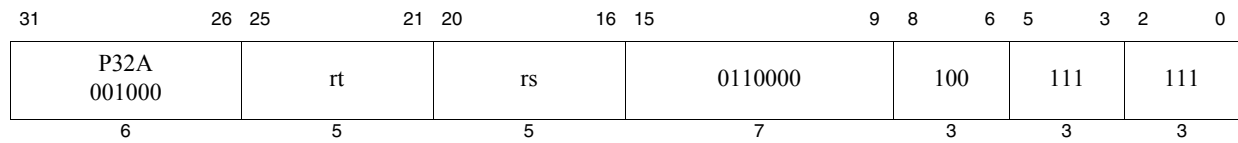
```

Exceptions:

Reserved Instruction, DSP Disabled

PRECEQ.W.PHL

Precision Expand Fractional Halfword to Fractional Word Value

PRECEQ.W.PHR**Precision Expand Fractional Halfword to Fractional Word Value****Format:** PRECEQ.W.PHR *rt*, *rs***DSP****Purpose:** Precision Expand Fractional Halfword to Fractional Word Value

Expand the precision of a Q15 fractional value taken from the right element of a paired halfword vector to create a Q31 fractional word value.

Description: $rt \leftarrow \text{expand_prec}(rs_{15..0})$

The right Q15 fractional halfword value from the paired halfword vector in register *rs* is expanded to a Q31 fractional value and written to destination register *rt*. The precision expansion is achieved by appending 16 least-significant zero bits to the original halfword value to generate the 32-bit fractional value.

Restrictions:

No data-dependent exceptions are possible.

The operands must be values in the specified format. If they are not, the results are **UNPREDICTABLE** and the values of the operand vectors become **UNPREDICTABLE**.

Operation:

```

ValidateAccessToDSPResources()
temp31..0 ← GPR[rs]15..0 || 016
GPR[rt]31..0 ← temp31..0

```

Exceptions:

Reserved Instruction, DSP Disabled

PRECEQ.W.PHR

Precision Expand Fractional Halfword to Fractional Word Value

PRECEQU.PH.QBL**Precision Expand two Unsigned Bytes to Fractional Halfword Values**

31	26	25	21	20	16	15	9	8	6	5	3	2	0	
P32A 001000			rt		rs		0111000			100		111		111
6			5		5		7			3		3		3

Format: PRECEQU.PH.QBL *rt*, *rs***DSP****Purpose:** Precision Expand two Unsigned Bytes to Fractional Halfword Values

Expand the precision of two unsigned byte values taken from the two left-most elements of a quad byte vector to create two Q15 fractional halfword values.

Description: $rt \leftarrow \text{expand_prec}(rs_{31..24}) \parallel \text{expand_prec}(rs_{23..16})$

The two left-most unsigned integer byte values from the four byte elements in register *rs* are expanded to create two Q15 fractional values that are then written to destination register *rt*. The precision expansion is achieved by pre-pending a single zero bit (for positive sign) to the original byte value and appending seven least-significant zeros to generate each 16-bit fractional value.

Restrictions:

No data-dependent exceptions are possible.

The operands must be values in the specified format. If they are not, the results are **UNPREDICTABLE** and the values of the operand vectors become **UNPREDICTABLE**.

Operation:

```

ValidateAccessToDSPResources()
tempB15..0 ← 01 || GPR[rs]31..24 || 07
tempA15..0 ← 01 || GPR[rs]23..16 || 07
GPR[rt]31..0 ← tempB15..0 || tempA15..0

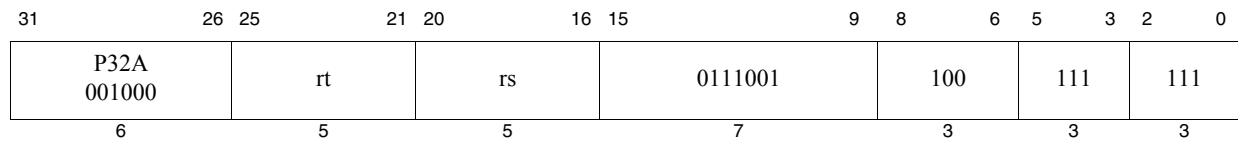
```

Exceptions:

Reserved Instruction, DSP Disabled

PRECEQU.PH.QBL

Precision Expand two Unsigned Bytes to Fractional Halfword Values

PRECEQU.PH.QBLA**Precision Expand two Unsigned Bytes to Fractional Halfword Values****Format:** PRECEQU.PH.QBLA *rt*, *rs***DSP****Purpose:** Precision Expand two Unsigned Bytes to Fractional Halfword Values

Expand the precision of two unsigned byte values taken from the two left-alternate aligned elements of a quad byte vector to create two Q15 fractional halfword values.

Description: $rt \leftarrow \text{expand_prec}(rs_{31..24}) \parallel \text{expand_prec}(rs_{15..8})$

The two left-alternate aligned unsigned integer byte values from the four byte elements in register *rs* are expanded to create two Q15 fractional values that are then written to destination register *rt*. The precision expansion is achieved by pre-pending a single zero bit (for positive sign) to the original byte value and appending seven least-significant zeros to generate each 16-bit fractional value.

Restrictions:

No data-dependent exceptions are possible.

The operands must be values in the specified format. If they are not, the results are **UNPREDICTABLE** and the values of the operand vectors become **UNPREDICTABLE**.

Operation:

```

ValidateAccessToDSPResources()
tempB15..0 ← 01 || GPR[rs]31..24 || 07
tempA15..0 ← 01 || GPR[rs]15..8 || 07
GPR[rt]31..0 ← tempB15..0 || tempA15..0

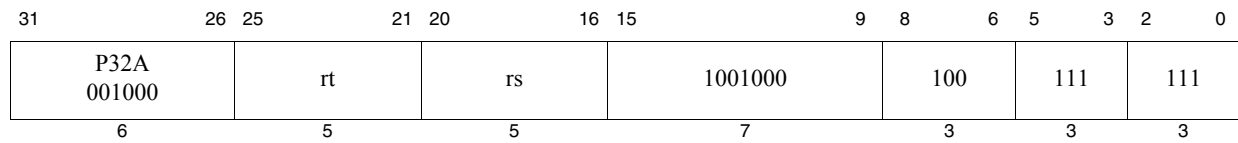
```

Exceptions:

Reserved Instruction, DSP Disabled

PRECEQU.PH.QBLA

Precision Expand two Unsigned Bytes to Fractional Halfword Values

PRECEQU.PH.QBR**Precision Expand two Unsigned Bytes to Fractional Halfword Values****Format:** PRECEQU.PH.QBR *rt*, *rs***DSP****Purpose:** Precision Expand two Unsigned Bytes to Fractional Halfword Values

Expand the precision of two unsigned byte values taken from the two right-most elements of a quad byte vector to create two Q15 fractional halfword values.

Description: $rt \leftarrow \text{expand_prec}(rs_{15..8}) \parallel \text{expand_prec}(rs_{7..0})$

The two right-most unsigned integer byte values from the four byte elements in register *rs* are expanded to create two Q15 fractional values that are then written to destination register *rt*. The precision expansion is achieved by pre-pending a single zero bit (for positive sign) to the original byte value and appending seven least-significant zeros to generate each 16-bit fractional value.

Restrictions:

No data-dependent exceptions are possible.

The operands must be values in the specified format. If they are not, the results are **UNPREDICTABLE** and the values of the operand vectors become **UNPREDICTABLE**.

Operation:

```

ValidateAccessToDSPResources()
tempB15..0 ← 01 || GPR[rs]15..8 || 07
tempA15..0 ← 01 || GPR[rs]7..0 || 07
GPR[rt]31..0 ← tempB15..0 || tempA15..0

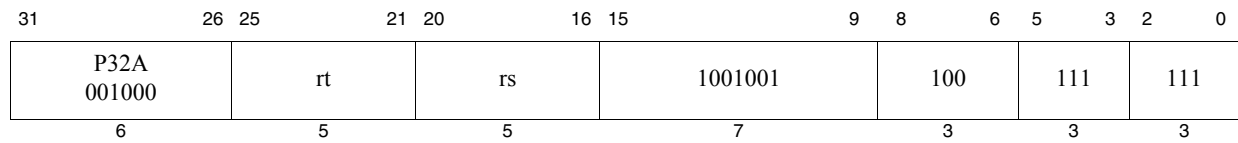
```

Exceptions:

Reserved Instruction, DSP Disabled

PRECEQU.PH.QBR

Precision Expand two Unsigned Bytes to Fractional Halfword Values

PRECEQU.PH.QBRA**Precision Expand two Unsigned Bytes to Fractional Halfword Values****Format:** PRECEQU.PH.QBRA *rt*, *rs***DSP****Purpose:** Precision Expand two Unsigned Bytes to Fractional Halfword Values

Expand the precision of two unsigned byte values taken from the two right-alternate aligned elements of a quad byte vector to create two Q15 fractional halfword values.

Description: $rt \leftarrow \text{expand_prec}(rs_{23..16}) \parallel \text{expand_prec}(rs_{7..0})$

The two right-alternate aligned unsigned integer byte values from the four byte elements in register *rs* are expanded to create two Q15 fractional values that are then written to destination register *rt*. The precision expansion is achieved by pre-pending a single zero bit (for positive sign) to the original byte value and appending seven least-significant zeros to generate each 16-bit fractional value.

Restrictions:

No data-dependent exceptions are possible.

The operands must be values in the specified format. If they are not, the results are **UNPREDICTABLE** and the values of the operand vectors become **UNPREDICTABLE**.

Operation:

```

ValidateAccessToDSPResources()
tempB15..0 ← 01 || GPR[rs]23..16 || 07
tempA15..0 ← 01 || GPR[rs]7..0 || 07
GPR[rt]31..0 ← tempB15..0 || tempA15..0

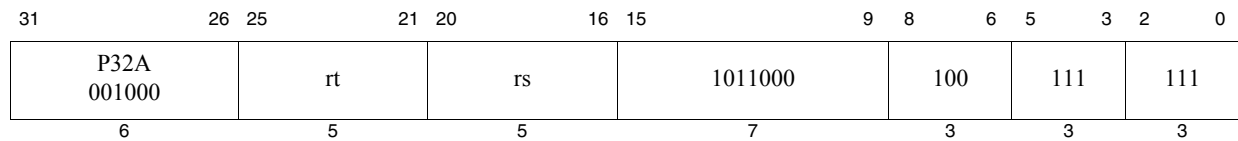
```

Exceptions:

Reserved Instruction, DSP Disabled

PRECEQU.PH.QBRA

Precision Expand two Unsigned Bytes to Fractional Halfword Values

PRECEU.PH.QBL**Precision Expand Two Unsigned Bytes to Unsigned Halfword Values****Format:** PRECEU.PH.QBL *rt*, *rs***DSP****Purpose:** Precision Expand Two Unsigned Bytes to Unsigned Halfword Values

Expand the precision of two unsigned byte values taken from the two left-most elements of a quad byte vector to create two unsigned halfword values.

Description: $rt \leftarrow \text{expand_prec8u16}(rs_{31..24}) \parallel \text{expand_prec8u16}(rs_{23..16})$

The two left-most unsigned integer byte values from the four byte elements in register *rs* are expanded to create two unsigned halfword values that are then written to destination register *rt*. The precision expansion is achieved by pre-pending eight most-significant zeros to each original value to generate each 16 bit unsigned value.

Restrictions:

No data-dependent exceptions are possible.

The operands must be values in the specified format. If they are not, the results are **UNPREDICTABLE** and the values of the operand vectors become **UNPREDICTABLE**.

Operation:

```

ValidateAccessToDSPResources()
tempB15..0 ← 08 || GPR[rs]31..24
tempA15..0 ← 08 || GPR[rs]23..16
GPR[rt]31..0 ← tempB15..0 || tempA15..0

```

Exceptions:

Reserved Instruction, DSP Disabled

PRECEU.PH.QBL

Precision Expand Two Unsigned Bytes to Unsigned Halfword Values

PRECEU.PH.QBLA**Precision Expand Two Unsigned Bytes to Unsigned Halfword Values**

31	26	25	21	20	16	15	9	8	6	5	3	2	0
P32A 001000		rt		rs		1011001		100		111		111	
6		5		5		7		3		3		3	

Format: PRECEU.PH.QBLA *rt*, *rs***DSP****Purpose:** Precision Expand Two Unsigned Bytes to Unsigned Halfword Values

Expand the precision of two unsigned integer byte values taken from the two left-alternate aligned positions of a quad byte vector to create four unsigned halfword values.

Description: $rt \leftarrow \text{expand_prec8ul6}(rs_{31..24}) \parallel \text{expand_prec8ul6}(rs_{15..8})$

The two left-alternate aligned unsigned integer byte values from the four right-most byte elements in register *rs* are each expanded to unsigned halfword values and written to destination register *rt*. The precision expansion is achieved by pre-pending eight most-significant zero bits to the original byte value to generate each 16 bit unsigned halfword value.

Restrictions:

No data-dependent exceptions are possible.

The operands must be values in the specified format. If they are not, the results are **UNPREDICTABLE** and the values of the operand vectors become **UNPREDICTABLE**.

Operation:

```

ValidateAccessToDSPResources()
tempB15..0 ← 08 || GPR[rs]31..24
tempA15..0 ← 08 || GPR[rs]15..8
GPR[rt]31..0 ← tempB15..0 || tempA15..0

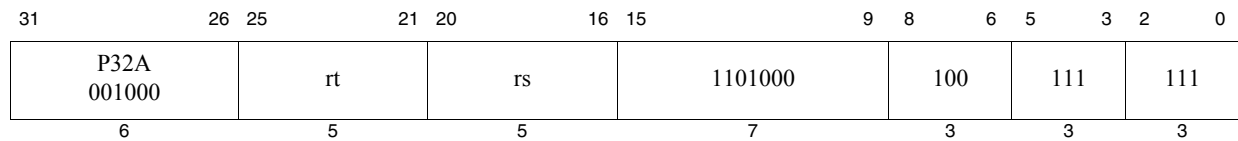
```

Exceptions:

Reserved Instruction, DSP Disabled

PRECEU.PH.QBLA

Precision Expand Two Unsigned Bytes to Unsigned Halfword Values

PRECEU.PH.QBR**Precision Expand two Unsigned Bytes to Unsigned Halfword Values****Format:** PRECEU.PH.QBR *rt*, *rs***DSP****Purpose:** Precision Expand two Unsigned Bytes to Unsigned Halfword Values

Expand the precision of two unsigned integer byte values taken from the two right-most elements of a quad byte vector to create two unsigned halfword values.

Description: $rt \leftarrow \text{expand_prec8u16}(rs_{15..8}) \parallel \text{expand_prec8u16}(rs_{7..0})$

The two right-most unsigned integer byte values from the four byte elements in register *rs* are expanded to create two unsigned halfword values that are then written to destination register *rt*. The precision expansion is achieved by pre-pending eight most-significant zero bits to each original value to generate each 16 bit halfword value.

Restrictions:

No data-dependent exceptions are possible.

The operands must be values in the specified format. If they are not, the results are **UNPREDICTABLE** and the values of the operand vectors become **UNPREDICTABLE**.

Operation:

```

ValidateAccessToDSPResources()
tempB15..0 ← 08 || GPR[rs]15..8
tempA15..0 ← 08 || GPR[rs]7..0
GPR[rt]31..0 ← tempB15..0 || tempA15..0

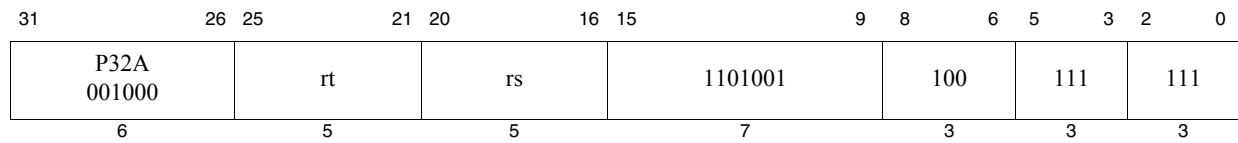
```

Exceptions:

Reserved Instruction, DSP Disabled

PRECEU.PH.QBR

Precision Expand two Unsigned Bytes to Unsigned Halfword Values

PRECEU.PH.QBRA**Precision Expand Two Unsigned Bytes to Unsigned Halfword Values****Format:** PRECEU.PH.QBRA *rt*, *rs***DSP****Purpose:** Precision Expand Two Unsigned Bytes to Unsigned Halfword Values

Expand the precision of two unsigned byte values taken from the two right-alternate aligned positions of a quad byte vector to create two unsigned halfword values.

Description: $rt \leftarrow \text{expand_prec8ul6}(rs_{23..16}) \parallel \text{expand_prec8ul6}(rs_{7..0})$

The two right-alternate aligned unsigned integer byte values from the four byte elements in register *rs* are each expanded to unsigned halfword values and written to destination register *rt*. The precision expansion is achieved by pre-pending eight most-significant zero bits to the original byte value to generate each 16 bit unsigned halfword value.

Restrictions:

No data-dependent exceptions are possible.

The operands must be values in the specified format. If they are not, the results are **UNPREDICTABLE** and the values of the operand vectors become **UNPREDICTABLE**.

Operation:

```

ValidateAccessToDSPResources()
tempB15..0 ← 08 || GPR[rs]23..16
tempA15..0 ← 08 || GPR[rs]7..0
GPR[rt]31..0 ← tempB15..0 || tempA15..0

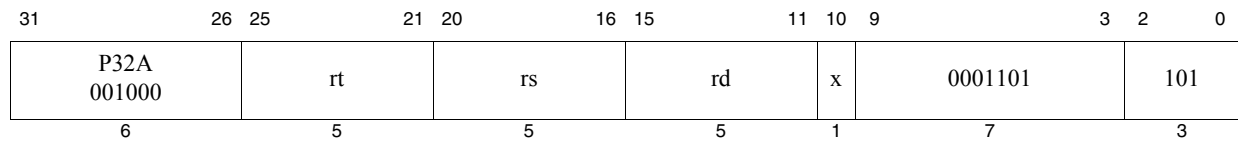
```

Exceptions:

Reserved Instruction, DSP Disabled

PRECEU.PH.QBRA

Precision Expand Two Unsigned Bytes to Unsigned Halfword Values

PRECR.QB.PH**Precision Reduce Four Integer Halfwords to Four Bytes****Format:** PRECR.QB.PH rd, rs, rt**DSP-R2****Purpose:** Precision Reduce Four Integer Halfwords to Four Bytes

Reduce the precision of four integer halfwords to four byte values.

Description: $rd \leftarrow rs_{23..16} \parallel rs_{7..0} \parallel rt_{23..16} \parallel rt_{7..0}$

The 8 least-significant bits from each of the two integer halfword values in registers *rs* and *rt* are taken to produce four byte-sized results that are written to the four byte elements in destination register *rd*. The two bytes values obtained from *rs* are written to the two left-most destination byte elements, and the two bytes obtained from *rt* are written to the two right-most destination byte elements.

Restrictions:

No data-dependent exceptions are possible.

The operands must be a value in the specified format. If they are not, the results are **UNPREDICTABLE** and the values of the operand vectors become **UNPREDICTABLE**.

Operation:

```

ValidateAccessToDSP2Resources()
tempD7..0 ← GPR[rs]23..16
tempC7..0 ← GPR[rs]7..0
tempB7..0 ← GPR[rt]23..16
tempA7..0 ← GPR[rt]7..0
GPR[rd]31..0 ← tempD7..0 || tempC7..0 || tempB7..0 || tempA7..0

```

Exceptions:

Reserved Instruction, DSP Disabled

PREC.R.QB.PH

Precision Reduce Four Integer Halfwords to Four Bytes

PRECR_SRA[_R].PH.W**Precision Reduce Two Integer Words to Halfwords after a Right Shift**

31	26	25	21	20	16	15	11	10	9	3	2	0
PRECR_SRA.PH.W												
P32A 001000	rt		rs		sa		0	1111001			101	
PRECR_SRA_R.PH.W												
P32A 001000	rt		rs		sa		1	1111001			101	
6	5		5		5		1	7			3	

Format: PRECR_SRA[_R].PH.W

PRECR_SRA.PH.W rt, rs, sa

DSP-R2

PRECR_SRA_R.PH.W rt, rs, sa

DSP-R2**Purpose:** Precision Reduce Two Integer Words to Halfwords after a Right Shift

Do an arithmetic right shift of two integer words with optional rounding, and then reduce the precision to halfwords.

Description: $rt \leftarrow (\text{round}(rt \gg \text{shift}))_{15..0} \parallel (\text{round}(rs \gg \text{shift}))_{15..0}$

The two words in registers *rs* and *rt* are right shifted arithmetically by the specified shift amount *sa* to create interim results. The 16 least-significant bits of each interim result are then written to the corresponding elements of destination register *rt*.

In the rounding version of the instruction, a value of 1 is added at the most-significant discarded bit position after the shift is performed. The 16 least-significant bits of each interim result are then written to the corresponding elements of destination register *rt*.

The shift amount *sa* is interpreted as a five-bit unsigned integer taking values between 0 and 31.

This instruction does not write any bits of the *ouflag* field in the *DSPControl* register.

Restrictions:

No data-dependent exceptions are possible.

The operands must be a value in the specified format. If they are not, the results are **UNPREDICTABLE** and the values of the operand vectors become **UNPREDICTABLE**.

Operation:

```

PRECR_SRA.PH.W
  ValidateAccessToDSP2Resources()
  if (sa4..0 = 0) then
    tempB15..0 ← GPR[rt]15..0
    tempA15..0 ← GPR[rs]15..0
  else
    tempB15..0 ← ( (GPR[rt]31)sa || GPR[rt]31..sa )
    tempA15..0 ← ( (GPR[rs]31)sa || GPR[rs]31..sa )
  endif
  GPR[rt]31..0 ← tempB15..0 || tempA15..0

PRECR_SRA_R.PH.W
  ValidateAccessToDSP2Resources()
  if (sa4..0 = 0) then
    tempB16..0 ← ( GPR[rt]15..0 || 0 )
    tempA16..0 ← ( GPR[rs]15..0 || 0 )
  else
    tempB32..0 ← ( (GPR[rt]31)sa || GPR[rt]31..sa-1 ) + 1
    tempA32..0 ← ( (GPR[rs]31)sa || GPR[rs]31..sa-1 ) + 1
  endif

```

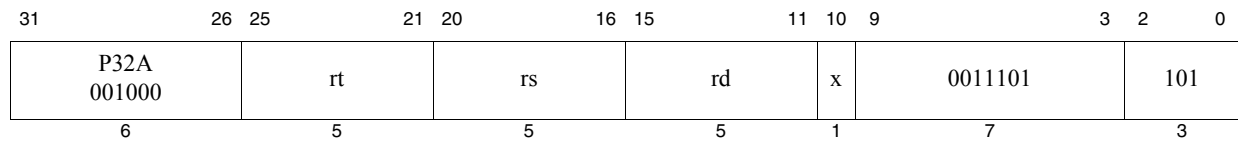

PRECR_SRA[_R].PH.W

Precision Reduce Two Integer Words to Halfwords after a Right Shift

$$\text{GPR}[\text{rt}]_{31..0} \leftarrow \text{tempB}_{16..1} \parallel \text{tempA}_{16..1}$$

Exceptions:

Reserved Instruction, DSP Disabled

PRECQR.PH.W**Precision Reduce Fractional Words to Fractional Halfwords****Format:** PRECQR.PH.W rd, rs, rt**DSP****Purpose:** Precision Reduce Fractional Words to Fractional Halfwords

Reduce the precision of two fractional words to produce two fractional halfword values.

Description: $rd \leftarrow rt_{31..16} \parallel rs_{31..16}$

The 16 most-significant bits from each of the Q31 fractional word values in registers *rs* and *rt* are written to destination register *rd*, creating a vector of two Q15 fractional values. The fractional word from the *rs* register is used to create the left-most Q15 fractional value in *rd*, and the fractional word from the *rt* register is used to create the right-most Q15 fractional value.

Restrictions:

No data-dependent exceptions are possible.

The operands must be values in the specified format. If they are not, the results are **UNPREDICTABLE** and the values of the operand vectors become **UNPREDICTABLE**.

Operation:

```

ValidateAccessToDSPResources()
tempB15..0 ← GPR[rs]31..16
tempA15..0 ← GPR[rt]31..16
GPR[rd]31..0 ← tempB15..0 || tempA15..0

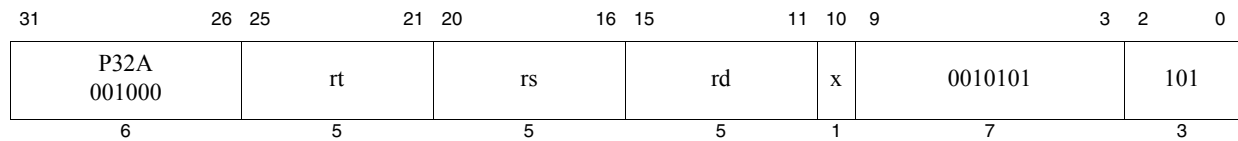
```

Exceptions:

Reserved Instruction, DSP Disabled

PRECQR.PH.W

Precision Reduce Fractional Words to Fractional Halfwords

PRECQR.QB.PH**Precision Reduce Four Fractional Halfwords to Four Bytes****Format:** PRECQR.QB.PH rd, rs, rt**DSP****Purpose:** Precision Reduce Four Fractional Halfwords to Four Bytes

Reduce the precision of four fractional halfwords to four byte values.

Description: $rd \leftarrow rs_{31..24} \parallel rs_{15..8} \parallel rt_{31..24} \parallel rt_{15..8}$

The four Q15 fractional values in registers *rs* and *rt* are truncated by dropping the eight least significant bits from each value to produce four fractional byte values. The four fractional byte values are written to the four byte elements of destination register *rd*. The two values obtained from register *rt* are placed in the two right-most byte positions in the destination register, and the two values obtained from register *rs* are placed in the two remaining byte positions.

Restrictions:

No data-dependent exceptions are possible.

The operands must be values in the specified format. If they are not, the results are **UNPREDICTABLE** and the values of the operand vectors become **UNPREDICTABLE**.

Operation:

```

ValidateAccessToDSPResources()
tempD7..0 ← GPR[rs]31..24
tempC7..0 ← GPR[rs]15..8
tempB7..0 ← GPR[rt]31..24
tempA7..0 ← GPR[rt]15..8
GPR[rd]31..0 ← tempD7..0 || tempC7..0 || tempB7..0 || tempA7..0

```

Exceptions:

Reserved Instruction, DSP Disabled

PRECQR.QB.PH

Precision Reduce Four Fractional Halfwords to Four Bytes

PRECQRQU_S.QB.PH**Precision Reduce Fractional Halfwords to Unsigned Bytes With Saturation**

31	26	25	21	20	16	15	11	10	9	3	2	0	
P32A 001000			rt		rs		rd		x	0101101			101
6			5		5		5		1	7			3

Format: PRECQRQU_S.QB.PH rd, rs, rt**DSP****Purpose:** Precision Reduce Fractional Halfwords to Unsigned Bytes With Saturation

Reduce the precision of four fractional halfwords with saturation to produce four unsigned byte values, with saturation.

Description: $rd \leftarrow \text{sat}(\text{reduce_prec}(rs_{31..16})) \parallel \text{sat}(\text{reduce_prec}(rs_{15..0})) \parallel \text{sat}(\text{reduce_prec}(rt_{31..16})) \parallel \text{sat}(\text{reduce_prec}(rt_{15..0}))$

The four Q15 fractional halfwords from registers *rs* and *rt* are used to create four unsigned byte values that are written to corresponding elements of destination register *rd*. The two halfwords from the *rs* register and the two halfwords from the *rt* register are used to create the four unsigned byte values.

Each unsigned byte value is created from the Q15 fractional halfword input value after first examining the sign and magnitude of the halfword. If the sign of the halfword value is positive and the value is greater than 0x7F80 hexadecimal, the result is clamped to the maximum positive 8-bit value (255 decimal, 0xFF hexadecimal). If the sign of the halfword value is negative, the result is clamped to the minimum positive 8-bit value (0 decimal, 0x00 hexadecimal). Otherwise, the sign bit is discarded from the input and the result is taken from the eight most-significant bits that remain.

If clamping was needed to produce any of the unsigned output values, bit 22 of the *DSPControl* register is set to 1.

Restrictions:

No data-dependent exceptions are possible.

The operands must be values in the specified format. If they are not, the results are **UNPREDICTABLE** and the values of the operand vectors become **UNPREDICTABLE**.

Operation:

```

ValidateAccessToDSPResources()
tempD7..0 ← sat8ReducePrecision( GPR[rs]31..16 )
tempC7..0 ← sat8ReducePrecision( GPR[rs]15..0 )
tempB7..0 ← sat8ReducePrecision( GPR[rt]31..16 )
tempA7..0 ← sat8ReducePrecision( GPR[rt]15..0 )
GPR[rd]31..0 ← tempD7..0 || tempC7..0 || tempB7..0 || tempA7..0

function sat8ReducePrecision( a15..0 )
    sign ← a15
    mag14..0 ← a14..0
    if ( sign = 0 ) then
        if ( mag14..0 > 0x7F80 ) then
            temp7..0 ← 0xFF
            DSPControl_ouflag:22 ← 1
        else
            temp7..0 ← mag14..7
        endif
    else
        temp7..0 ← 0x00
        DSPControl_ouflag:22 ← 1
    endif
    return temp7..0
endfunction sat8ReducePrecision

```

PRECRQU_S.QB.PH

Precision Reduce Fractional Halfwords to Unsigned Bytes With Saturation

Exceptions:

Reserved Instruction, DSP Disabled

PRECQRQ_RS.PH.W Precision Reduce Fractional Words to Halfwords With Rounding and Saturation

31	26	25	21	20	16	15	11	10	9	3	2	0
P32A 001000	rt		rs		rd		x	0100101			101	
6	5		5		5		1	7			3	

Format: PRECQRQ_RS.PH.W rd, rs, rt**DSP****Purpose:** Precision Reduce Fractional Words to Halfwords With Rounding and Saturation

Reduce the precision of two fractional words to produce two fractional halfword values, with rounding and saturation.

Description: $rd \leftarrow \text{truncQ15SatRound}(rs_{31..0}) \mid \mid \text{truncQ15SatRound}(rt_{31..0})$

The two Q31 fractional word values in each of registers *rs* and *rt* are used to create two Q15 fractional halfword values that are written to the two halfword elements in destination register *rd*. The fractional word from the *rs* register is used to create the left-most Q15 fractional halfword result in *rd*, and the fractional word from the *rt* register is used to create the right-most halfword value.

Each input Q31 fractional value is rounded and saturated before being truncated to create the Q15 fractional halfword result. First, the value 0x00008000 is added to the input Q31 value to round even, creating an interim rounded result. If this addition causes overflow, the interim rounded result is saturated to the maximum Q31 value (0x7FFFFFFF hexadecimal). Then, the 16 least-significant bits of the interim rounded and saturated result are discarded and the 16 most-significant bits are written to the destination register in the appropriate position.

If either of the rounding operations results in overflow and saturation, a 1 is written to bit 22 in the *DSPControl* register within the *ouflag* field.

Restrictions:

No data-dependent exceptions are possible.

The operands must be values in the specified format. If they are not, the results are **UNPREDICTABLE** and the values of the operand vectors become **UNPREDICTABLE**.

Operation:

```

ValidateAccessToDSPResources()
tempB15..0 ← trunc16Sat16Round( GPR[rs]31..0 )
tempA15..0 ← trunc16Sat16Round( GPR[rt]31..0 )
GPR[rd]31..0 ← tempB15..0 || tempA15..0

function trunc16Sat16Round( a31..0 )
    temp32..0 ← ( a31 || a31..0 ) + 0x00008000
    if ( temp32 ≠ temp31 ) then
        temp32..0 ← 0 || 0x7FFFFFFF
        DSPControlouflag:22 ← 1
    endif
    return temp31..16
endfunction trunc16Sat16Round

```

Exceptions:

Reserved Instruction, DSP Disabled

PRECQR_RS.PH.W Precision Reduce Fractional Words to Halfwords With Rounding and Saturation

PREPEND**Right Shift and Prepend Bits to the MSB**

Format: PREPEND *rt*, *rs*, *sa*
 EXTW *rt*, *rs*, *rt*, *sa*

DSP-R2
 Replaced with EXTW in nanoMIPS

Purpose: Right Shift and Prepend Bits to the MSB

Logically right-shift the first source register, replacing the bits emptied by the shift with bits from the source register.

Description: $rt \leftarrow rs_{sa-1..0} \parallel (rt \gg sa)$

The word value in register *rt* is logically right-shifted by the specified shift amount *sa*, and *sa* bits from the least-significant positions of register *rs* are written into the *sa* most-significant bits emptied by the shift. The result is then written to destination register *rt*.

Restrictions:

No data-dependent exceptions are possible.

The operands must be a value in the specified format. If they are not, the results are **UNPREDICTABLE** and the values of the operand vectors become **UNPREDICTABLE**.

Operation:

```

ValidateAccessToDSP2Resources()
if ( sa4..0 = 0 ) then
    temp31..0 ← GPR[rt]31..0
else
    temp31..0 ← ( GPR[rs]sa-1..0 || GPR[rt]31..sa )
endif
GPR[rt]31..0 = temp31..0

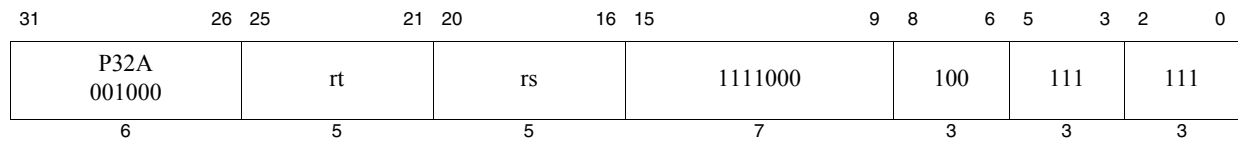
```

Exceptions:

Reserved Instruction, DSP Disabled

PREPEND

Right Shift and Prepend Bits to the MSB

RADDU.W.QB**Unsigned Reduction Add Vector Quad Bytes****Format:** RADDU.W.QB rt, rs**DSP****Purpose:** Unsigned Reduction Add Vector Quad Bytes

Reduction add of four unsigned byte values in a vector register to produce an unsigned word result.

Description: $rt \leftarrow \text{zero_extend}(rs_{31..24} + rs_{23..16} + rs_{15..8} + rs_{7..0})$ The unsigned byte elements in register *rs* are added together as unsigned 8-bit values, and the result is zero extended to a word and written to register *rt*.**Restrictions:**

No data-dependent exceptions are possible.

The operands must be in the specified format. If they are not, the results are **UNPREDICTABLE** and the values of the operand vectors become **UNPREDICTABLE**.**Operation:**

```

ValidateAccessToDSPResources()
temp9..0 ← ( 02 || GPR[rs]31..24 ) + ( 02 || GPR[rs]23..16 ) + ( 02 || GPR[rs]15..8 ) +
( 02 || GPR[rs]7..0 )
GPR[rt]31..0 ← 0(GPRELEN-10) || temp9..0

```

Exceptions:

Reserved Instruction, DSP Disabled

RADDU.W.QB

Unsigned Reduction Add Vector Quad Bytes

RDDSP**Read DSPControl Register Fields to a GPR**

31	26	25	21	20	14	13	12	11	9	8	6	5	3	2	0
P32A 001000			rt		mask			00	011	001	111		111		
6			5		7			2	3	3	3		3		

Format: RDDSP
 RDDSP rt, mask
 RDDSP rt

**DSP
 Assembly Idiom**

Purpose: Read DSPControl Register Fields to a GPR

To copy selected fields from the special-purpose *DSPControl* register to the specified GPR.

Description: $rt \leftarrow \text{select}(\text{mask}, \text{DSPControl})$

Selected fields in the special register *DSPControl* are copied into the corresponding bits of destination register *rt*. Each of bits 0 through 5 of the *mask* operand corresponds to a specific field in the *DSPControl* register. A mask bit value of 1 indicates that the bits from the corresponding field in *DSPControl* will be copied into the same bit positions in register *rt*, and a mask bit value of 0 indicates that the corresponding bit positions in *rt* will be set to zero. Bits 6 through 9 of the *mask* operand are ignored.

The table below shows the correspondence between the bits in the *mask* operand and the fields in the *DSPControl* register; mask bit 0 is the least-significant bit in *mask*.

Bit	31	24	23	16	15	14	13	12	7	6	5	0
DSPControl field	ccond		ouflag		0	EFI	C	scount			pos	
Mask bit	4		3			5	2	1			0	

For example, to copy only the bits from the *scount* field in *DSPControl*, the value of the *mask* operand used will be 2 decimal (0x02 hexadecimal). After execution of the instruction, bits 7 through 12 of register *rt* will have the value of bits 7 through 12 from the *scount* field in *DSPControl*. The remaining bits in register *rt* will be set to zero.

The one-operand version of the instruction provides a convenient assembly idiom that allows the programmer to read all fields in the *DSPControl* register into the destination register, i.e., it is equivalent to specifying a *mask* value of 31 decimal (0x1F hexadecimal).

Restrictions:

No data-dependent exceptions are possible.

The operands must be values in the specified format. If they are not, the results are **UNPREDICTABLE** and the values of the operand vectors become **UNPREDICTABLE**.

Operation:

```

ValidateAccessToDSPResources()
temp31..0 ← 032
if ( mask0 = 1 ) then
    temp5..0 ← DSPControlpos:5..0
endif
if ( mask1 = 1 ) then
    temp12..7 ← DSPControlscount:12..7
endif
if ( mask2 = 1 ) then
    temp13 ← DSPControlc:13
endif
if ( mask3 = 1 ) then
    temp23..16 ← DSPControlouflag:23..16

```

RDDSP**Read DSPControl Register Fields to a GPR**

```

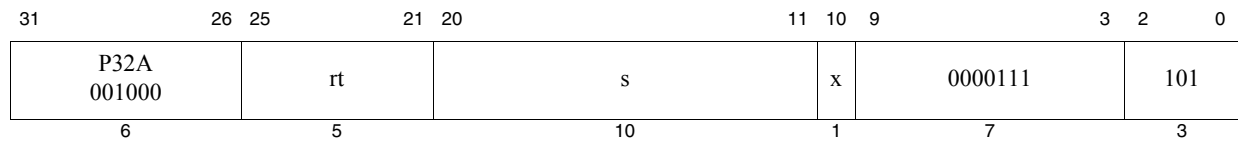
endif
if ( mask4 = 1 ) then
    temp27..24 ← DSPControlccond:27..24
endif
if ( mask5 = 1 ) then
    temp14 ← DSPControlefi:14
endif

GPR[rt]31..0 ← temp31..0

```

Exceptions:

Reserved Instruction, DSP Disabled

REPL.PH**Replicate Immediate Integer into all Vector Element Positions****Format:** REPL.PH rd, immediate**DSP****Purpose:** Replicate Immediate Integer into all Vector Element Positions

Replicate a sign-extended, 10-bit signed immediate integer value into the two halfwords in a halfword vector.

Description: $rd \leftarrow \text{sign_extend}(\text{immediate}) \parallel \text{sign_extend}(\text{immediate})$ The specified 10-bit signed immediate integer value is sign-extended to 16 bits and replicated into the two halfword positions in destination register *rd*.**Restrictions:**

No data-dependent exceptions are possible.

The operands must be values in the specified format. If they are not, the results are **UNPREDICTABLE** and the values of the operand vectors become **UNPREDICTABLE**.**Operation:**

```

ValidateAccessToDSPResources()
temp15..0 ← (immediate9)6 || immediate9..0
GPR[rd]31..0 ← temp15..0 || temp15..0

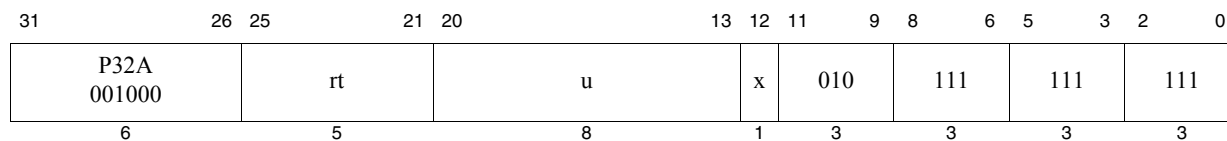
```

Exceptions:

Reserved Instruction, DSP Disabled

REPL.PH

Replicate Immediate Integer into all Vector Element Positions

REPL.QB**Replicate Immediate Integer into all Vector Element Positions****Format:** REPL.QB rt, immediate**DSP****Purpose:** Replicate Immediate Integer into all Vector Element Positions

Replicate a immediate byte into all elements of a quad byte vector.

Description: $rt \leftarrow \text{immediate} \parallel \text{immediate} \parallel \text{immediate} \parallel \text{immediate}$ The specified 8-bit signed immediate value is replicated into the four byte elements of destination register *rt*.**Restrictions:**

No data-dependent exceptions are possible.

The operands must be values in the specified format. If they are not, the results are **UNPREDICTABLE** and the values of the operand vectors become **UNPREDICTABLE**.**Operation:**

```

ValidateAccessToDSPResources()
temp7..0 ← immediate7..0
GPR[rt]31..0 ← temp7..0 || temp7..0 || temp7..0 || temp7..0

```

Exceptions:

Reserved Instruction, DSP Disabled

REPL.QB

Replicate Immediate Integer into all Vector Element Positions

REPLV.PH**Replicate a Halfword into all Vector Element Positions**

31	26	25	21	20	16	15	9	8	6	5	3	2	0
P32A 001000			rt		rs		0000001		100		111		111
6			5		5		7		3		3		3

Format: REPLV.PH *rt*, *rs***DSP****Purpose:** Replicate a Halfword into all Vector Element Positions

Replicate a variable halfword into the elements of a halfword vector.

Description: $rt \leftarrow (rs_{15..0} \parallel rs_{15..0})$ The halfword value in register *rs* is replicated into the two halfword elements of destination register *rt*.**Restrictions:**

No data-dependent exceptions are possible.

The operands must be values in the specified format. If they are not, the results are **UNPREDICTABLE** and the values of the operand vectors become **UNPREDICTABLE**.**Operation:**

```

ValidateAccessToDSPResources()
temp15..0 ← GPR[rs]15..0
GPR[rt]31..0 ← temp15..0 || temp15..0

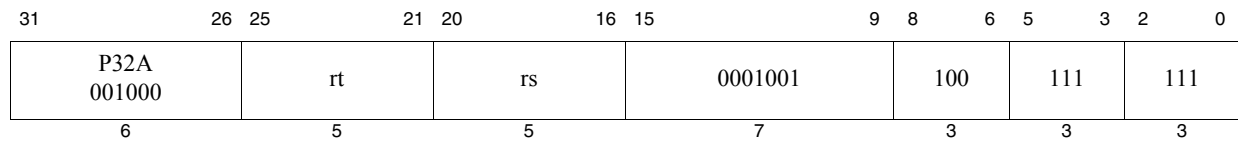
```

Exceptions:

Reserved Instruction, DSP Disabled

REPLV.PH

Replicate a Halfword into all Vector Element Positions

REPLV.QB**Replicate Byte into all Vector Element Positions****Format:** REPLV.QB rt, rs**DSP****Purpose:** Replicate Byte into all Vector Element Positions

Replicate a variable byte into all elements of a quad byte vector.

Description: $rt \leftarrow rs_{7..0} \parallel rs_{7..0} \parallel rs_{7..0} \parallel rs_{7..0}$ The byte value in register *rs* is replicated into the four byte elements of destination register *rt*.**Restrictions:**

No data-dependent exceptions are possible.

The operands must be values in the specified format. If they are not, the results are **UNPREDICTABLE** and the values of the operand vectors become **UNPREDICTABLE**.**Operation:**

```

ValidateAccessToDSPResources()
temp7..0 ← GPR[rs]7..0
GPR[rt]31..0 ← temp7..0 || temp7..0 || temp7..0 || temp7..0

```

Exceptions:

Reserved Instruction, DSP Disabled

REPLV.QB

Replicate Byte into all Vector Element Positions

SHILO**Shift an Accumulator Value Leaving the Result in the Same Accumulator**

31	26	25	22	21	16	15	14	13	10	9	3	2	0
P32A 001000		x		s		ac	x		0000011			101	
6		4		6		2	4		7			3	

Format: SHILO *ac*, *shift***DSP****Purpose:** Shift an Accumulator Value Leaving the Result in the Same AccumulatorShift the *HI/LO* paired value in a 64-bit accumulator either left or right, leaving the result in the same accumulator.**Description:** $ac \leftarrow (\text{shift} \geq 0) ? (ac \gg \text{shift}) : (ac \ll -\text{shift})$

The *HI/LO* register pair is treated as a single 64-bit accumulator that is shifted logically by *shift* bits, with the result of the shift written back to the source accumulator. The *shift* argument is a six-bit signed integer value: a positive argument results in a right shift of up to 31 bits, and a negative argument results in a left shift of up to 32 bits.

The value of *ac* can range from 0 to 3. When *ac*=0, this refers to the original *HI/LO* register pair of the MIPS32 architecture.

Restrictions:

No data-dependent exceptions are possible.

The operands must be values in the specified format. If they are not, the results are **UNPREDICTABLE** and the values of the operand vectors become **UNPREDICTABLE**.

Operation:

```

ValidateAccessToDSPResources()
sign ← shift5
shift5..0 ← ( sign = 0 ? shift5..0 : -shift5..0 )
if ( shift5..0 = 0 ) then
    temp63..0 ← (HI[ac]31..0 || LO[ac]31..0)
else
    if (sign = 0) then
        temp63..0 ← 0shift || (( HI[ac]31..0 || LO[ac]31..0 ) >> shift )
    else
        temp63..0 ← (( HI[ac]31..0 || LO[ac]31..0 ) << shift ) || 0shift
    endif
endif
( HI[ac]31..0 || LO[ac]31..0 ) ← temp63..32 || temp31..0

```

Exceptions:

Reserved Instruction, DSP Disabled

SHILO

Shift an Accumulator Value Leaving the Result in the Same Accumulator

SHILOV**Variable Shift of Accumulator Value Leaving the Result in the Same Accumulator**

31	26	25	21	20	16	15	14	13	12	11	9	8	6	5	3	2	0
P32A 001000		x		rs		ac	01	001		001		111		111			
6		5		5		2	2	3		3		3		3			

Format: SHILOV ac, rs**DSP****Purpose:** Variable Shift of Accumulator Value Leaving the Result in the Same Accumulator

Shift the *HI/LO* paired value in an accumulator either left or right by the amount specified in a GPR, leaving the result in the same accumulator.

Description: $ac \leftarrow (GPR[rs]_{6..0} \geq 0) ? (ac \gg GPR[rs]_{6..0}) : (ac \ll -GPR[rs]_{6..0})$

The *HI/LO* register pair is treated as a single 64-bit accumulator that is shifted logically by *shift* bits, with the result of the shift written back to the source accumulator. The *shift* argument is provided by the six least-significant bits of register *rs*; the remaining bits of *rs* are ignored. The *shift* argument is interpreted as a six-bit signed integer: a positive argument results in a right shift of up to 31 bits, and a negative argument results in a left shift of up to 32 bits.

The value of *ac* can range from 0 to 3. When *ac*=0, this refers to the original *HI/LO* register pair of the MIPS32 architecture.

Restrictions:

No data-dependent exceptions are possible.

The operands must be values in the specified format. If they are not, the results are **UNPREDICTABLE** and the values of the operand vectors become **UNPREDICTABLE**.

Operation:

```

ValidateAccessToDSPResources()
sign ← GPR[rs]5
shift5..0 ← ( sign = 0 ? GPR[rs]5..0 : -GPR[rs]5..0 )
if ( shift5..0 = 0 ) then
    temp63..0 ← ( HI[ac]31..0 || LO[ac]31..0 )
else
    if ( sign = 0 ) then
        temp63..0 ← 0shift || (( HI[ac]31..0 || LO[ac]31..0 ) >> shift )
    else
        temp63..0 ← (( HI[ac]31..0 || LO[ac]31..0 ) << shift ) || 0shift
    endif
endif
( HI[ac]31..0 || LO[ac]31..0 ) ← temp63..32 || temp31..0

```

Exceptions:

Reserved Instruction, DSP Disabled

SHILOV

Variable Shift of Accumulator Value Leaving the Result in the Same Accumulator

SHLL[_S].PH**Shift Left Logical Vector Pair Halfwords**

31	26	25	21	20	16	15	11	10	9	3	2	0
SHLL.PH												
P32A 001000	rt		rs		sa		00		1110110		101	
SHLL_S.PH												
P32A 001000	rt		rs		sa		10		1110110		101	
6	5		5		4		2		7		3	

Format: SHLL[_S].PH
 SHLL.PH rt, rs, sa
 SHLL_S.PH rt, rs, sa

DSP
DSP

Purpose: Shift Left Logical Vector Pair Halfwords

Element-wise shift of two independent halfwords in a vector data type by a fixed number of bits, with optional saturation.

Description: $rt \leftarrow \text{sat16}(rs_{31..16} \ll sa) \parallel (rs_{15..0} \ll sa)$

The two halfword values in register *rs* are each independently shifted left, inserting zeros into the least-significant bit positions emptied by the shift. In the saturating version of the instruction, if the shift results in an overflow the intermediate result is saturated to either the maximum positive or the minimum negative 16-bit value, depending on the sign of the original unshifted value. The two independent results are then written to the corresponding halfword elements of destination register *rt*.

This instruction writes a 1 to bit 22 in the *DSPControl* register in the *ouflag* field if any of the left shift operations results in an overflow or saturation.

Restrictions:

No data-dependent exceptions are possible.

The operands must be values in the specified format. If they are not, the results are **UNPREDICTABLE** and the values of the operand vectors become **UNPREDICTABLE**.

Operation:

```
SHLL.PH
  ValidateAccessToDSPResources()
  tempB15..0 ← shift16Left( GPR[rs]31..16, sa )
  tempA15..0 ← shift16Left( GPR[rs]15..0, sa )
  GPR[rt]31..0 ← tempB15..0 || tempA15..0

SHLL_S.PH
  ValidateAccessToDSPResources()
  tempB15..0 ← sat16ShiftLeft( GPR[rs]31..16, sa )
  tempA15..0 ← sat16ShiftLeft( GPR[rs]15..0, sa )
  GPR[rt]31..0 ← tempB15..0 || tempA15..0

function shift16Left( a15..0, s3..0 )
  if ( s3..0 = 0 ) then
    temp15..0 ← a15..0
  else
    sign ← a15
    temp15..0 ← ( a15-s..0 || 0s )
    discard15..0 ← ( sign(16-s) || a14..14-(s-1) )
    if (( discard15..0 ≠ 0x0000 ) and ( discard15..0 ≠ 0xFFFF )) then
```

SHLL[_S].PH**Shift Left Logical Vector Pair Halfwords**

```

        DSPControl_ouflag:22 ← 1
    endif
endif
return temp15..0
endfunction shift16Left

function sat16ShiftLeft( a15..0, s3..0 )
    if ( s3..0 = 0 ) then
        temp15..0 ← a15..0
    else
        sign ← a15
        temp15..0 ← ( a15-s..0 || 0s )
        discard15..0 ← ( sign(16-s) || a14..14-(s-1) )
        if (( discard15..0 ≠ 0x0000 ) and ( discard15..0 ≠ 0xFFFF )) then
            temp15..0 ← ( sign = 0 ? 0x7FFF : 0x8000 )
            DSPControl_ouflag:22 ← 1
        endif
    endif
endif
return temp15..0
endfunction sat16ShiftLeft

```

Exceptions:

Reserved Instruction, DSP Disabled

SHLL.QB**Shift Left Logical Vector Quad Bytes**

31	26	25	21	20	16	15	9	8	6	5	3	2	0
P32A 001000		rt		rs		sa	0	100	001	111	111		
6		5		5		3	1	3	3	3	3		

Format: SHLL.QB *rt*, *rs*, *sa***DSP****Purpose:** Shift Left Logical Vector Quad Bytes

Element-wise left shift of four independent bytes in a vector data type by a fixed number of bits.

Description: $rt \leftarrow (rs_{31..24} \ll sa) \parallel (rs_{23..16} \ll sa) \parallel (rs_{15..8} \ll sa) \parallel (rs_{7..0} \ll sa)$

The four byte values in register *rs* are each independently shifted left by *sa* bits and the *sa* least significant bits of each value are set to zero. The four independent results are then written to the corresponding byte elements of destination register *rt*.

This instruction writes a 1 to bit 22 in the *DSPControl* register in the *ouflag* field if any of the left shift operations results in an overflow.

Restrictions:

No data-dependent exceptions are possible.

The operands must be values in the specified format. If they are not, the results are **UNPREDICTABLE** and the values of the operand vectors become **UNPREDICTABLE**.

Operation:

```

ValidateAccessToDSPResources()
tempD7..0 ← shift8Left( GPR[rs]31..24, sa2..0 )
tempC7..0 ← shift8Left( GPR[rs]23..16, sa2..0 )
tempB7..0 ← shift8Left( GPR[rs]15..8, sa2..0 )
tempA7..0 ← shift8Left( GPR[rs]7..0, sa2..0 )
GPR[rt]31..0 ← tempD7..0 || tempC7..0 || tempB7..0 || tempA7..0

function shift8Left( a7..0, s2..0 )
  if ( s2..0 = 0 ) then
    temp7..0 ← a7..0
  else
    sign ← a7
    temp7..0 ← ( a7-s..0 || 0s )
    discard7..0 ← ( sign(8-s) || a6..6-(s-1) )
    if ( discard7..0 ≠ 0x00 ) then
      DSPControlouflag:22 ← 1
    endif
  endif
  return temp7..0
endfunction shift8Left

```

Exceptions:

Reserved Instruction, DSP Disabled

SHLL.QB

Shift Left Logical Vector Quad Bytes

SHLLV[_S].PH**Shift Left Logical Variable Vector Pair Halfwords**

31	26	25	21	20	16	15	11	10	9	3	2	0
SHLLV.PH												
P32A 001000		rt		rs		rd		0	1110001		101	
SHLLV_S.PH												
P32A 001000		rt		rs		rd		1	1110001		101	
6		5		5		5		1	7		3	

Format: SHLLV[_S].PH
 SHLLV.PH rd, rt, rs
 SHLLV_S.PH rd, rt, rs

DSP
DSP

Purpose: Shift Left Logical Variable Vector Pair Halfwords

Element-wise left shift of the two right-most independent halfwords in a vector data type by a variable number of bits, with optional saturation.

Description: $rd \leftarrow \text{sat16}(rt_{31..16} \ll rs_{3..0}) \parallel \text{sat16}(rt_{15..0} \ll rs_{3..0})$

The two halfword values in register *rt* are each independently shifted left by *shift* bits, inserting zeros into the least-significant bit positions emptied by the shift. In the saturating version of the instruction, if the shift results in an overflow the intermediate result is saturated to either the maximum positive or the minimum negative 16-bit value, depending on the sign of the original unshifted value. The two independent results are then written to the corresponding halfword elements of destination register *rd*.

The four least-significant bits of *rs* provide the shift value, interpreted as a four-bit unsigned integer; the remaining bits of *rs* are ignored.

This instruction writes a 1 to bit 22 in the *DSPControl* register in the ouflag field if any of the left shift operations results in an overflow or saturation.

Restrictions:

No data-dependent exceptions are possible.

The operands must be values in the specified format. If they are not, the results are **UNPREDICTABLE** and the values of the operand vectors become **UNPREDICTABLE**.

Operation:

```
SHLLV.PH
  ValidateAccessToDSPResources()
  tempB15..0 ← shift16Left( GPR[rt]31..16, GPR[rs]3..0 )
  tempA15..0 ← shift16Left( GPR[rt]15..0, GPR[rs]3..0 )
  GPR[rd]31..0 ← tempB15..0 || tempA15..0

SHLLV_S.PH
  ValidateAccessToDSPResources()
  tempB15..0 ← sat16ShiftLeft( GPR[rt]31..16, GPR[rs]3..0 )
  tempA15..0 ← sat16ShiftLeft( GPR[rt]15..0, GPR[rs]3..0 )
  GPR[rd]31..0 ← tempB15..0 || tempA15..0
```

Exceptions:

Reserved Instruction, DSP Disabled

SHLLV[_S].PH

Shift Left Logical Variable Vector Pair Halfwords

SHLLV.QB**Shift Left Logical Variable Vector Quad Bytes**

31	26	25	21	20	16	15	11	10	9	3	2	0
P32A 001000						rt		rs		rd		x
6						5		5		5		1
										1110010		101
										7		3

Format: SHLLV.QB rd, rt, rs**DSP****Purpose:** Shift Left Logical Variable Vector Quad Bytes

Element-wise left shift of four independent bytes in a vector data type by a variable number of bits.

Description: $rd \leftarrow (rt_{31..24} \ll rs_{2..0}) \parallel (rt_{23..16} \ll rs_{2..0}) \parallel (rt_{15..8} \ll rs_{2..0}) \parallel (rt_{7..0} \ll rs_{2..0})$

The four byte values in register *rt* are each independently shifted left by *sa* bits, inserting zeros into the least-significant bit positions emptied by the shift. The four independent results are then written to the corresponding byte elements of destination register *rd*.

The three least-significant bits of *rs* provide the shift value, interpreted as a three-bit unsigned integer; the remaining bits of *rs* are ignored.

This instruction writes a 1 to bit 22 in the *DSPControl* register in the *ouflag* field if any of the left shift operations results in an overflow.

Restrictions:

No data-dependent exceptions are possible.

The operands must be values in the specified format. If they are not, the results are **UNPREDICTABLE** and the values of the operand vectors become **UNPREDICTABLE**.

Operation:

```

ValidateAccessToDSPResources()
tempD7..0 ← shift8Left( GPR[rt]31..24, GPR[rs]2..0 )
tempC7..0 ← shift8Left( GPR[rt]23..16, GPR[rs]2..0 )
tempB7..0 ← shift8Left( GPR[rt]15..8, GPR[rs]2..0 )
tempA7..0 ← shift8Left( GPR[rt]7..0, GPR[rs]2..0 )
GPR[rd]31..0 ← tempD7..0 || tempC7..0 || tempB7..0 || tempA7..0

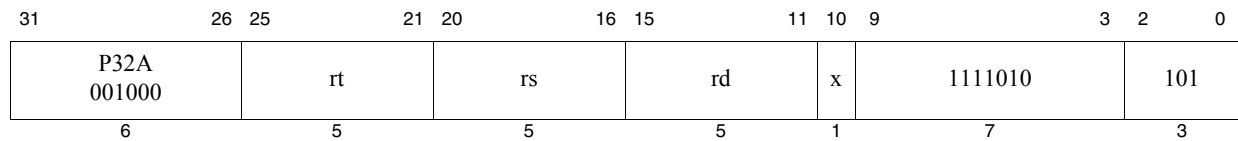
```

Exceptions:

Reserved Instruction, DSP Disabled

SHLLV.QB

Shift Left Logical Variable Vector Quad Bytes

SHLLV_S.W**Shift Left Logical Variable Vector Word****Format:** SHLLV_S.W rd, rt, rs**DSP****Purpose:** Shift Left Logical Variable Vector Word

A left shift of the word in a vector data type by a variable number of bits, with optional saturation.

Description: $rd \leftarrow \text{sat32}(rt_{31..0} \ll rs_{4..0})$

The word element in register *rt* is shifted left by *shift* bits, inserting zeros into the least-significant bit positions emptied by the shift. If the shift results in an overflow the intermediate result is saturated to either the maximum positive or the minimum negative 32-bit value, depending on the sign of the original unshifted value.

The shifted result is then written to destination register *rd*.

The five least-significant bits of *rs* are used as the shift value, interpreted as a five-bit unsigned integer; the remaining bits of *rs* are ignored.

This instruction writes a 1 to bit 22 in the *DSPControl* register in the *ouflag* field if either of the left shift operations results in an overflow or saturation.

Restrictions:

No data-dependent exceptions are possible.

The operands must be values in the specified format. If they are not, the results are **UNPREDICTABLE** and the values of the operand vectors become **UNPREDICTABLE**.

Operation:

```

ValidateAccessToDSPResources()
temp31..0 ← sat32ShiftLeft( GPR[rt]31..0, GPR[rs]4..0 )
GPR[rd]31..0 ← temp31..0

```

Exceptions:

Reserved Instruction, DSP Disabled

SHLLV_S.W

Shift Left Logical Variable Vector Word

SHLL_S.W**Shift Left Logical Word with Saturation**

31	26	25	21	20	16	15	11	10	9	3	2	0
P32A 001000	rt		rs		sa		x	1111110			101	
6	5		5		5		1	7			3	

Format: SHLL_S.W rt, rs, sa**DSP****Purpose:** Shift Left Logical Word with Saturation

To execute a left shift of a word with saturation by a fixed number of bits.

Description: $rt \leftarrow \text{sat32}(rs \ll sa)$

The 32-bit word in register *rs* is shifted left by *sa* bits, with zeros inserted into the bit positions emptied by the shift. If the shift results in a signed overflow, the shifted result is saturated to either the maximum positive (hexadecimal 0x7FFFFFFF) or minimum negative (hexadecimal 0x80000000) 32-bit value, depending on the sign of the original unshifted value. The shifted result is then written to destination register *rt*.

The instruction's *sa* field specifies the shift value, interpreted as a five-bit unsigned integer.

If the shift operation results in an overflow and saturation, this instruction writes a 1 to bit 22 of the *DSPControl* register within the *ouflag* field.

Restrictions:

No data-dependent exceptions are possible.

The operands must be values in the specified format. If they are not, the results are **UNPREDICTABLE** and the values of the operand vectors become **UNPREDICTABLE**.

Operation:

```

ValidateAccessToDSPResources()
temp31..0 ← sat32ShiftLeft( GPR[rs]31..0, sa4..0 )
GPR[rt]31..0 ← temp31..0

function sat32ShiftLeft( a31..0, s4..0 )
  if ( s = 0 ) then
    temp31..0 ← a
  else
    sign ← a31
    temp31..0 ← ( a31-s..0 || 0s )
    discard31..0 ← ( sign(32-s) || a30..30-(s-1) )
    if ( ( discard31..0 ≠ 0x00000000 ) and ( discard31..0 ≠ 0xFFFFFFFF ) ) then
      temp31..0 ← ( sign = 0 ? 0x7FFFFFFF : 0x80000000 )
      DSPControlouflag:22 ← 1
    endif
  endif
  return temp31..0
endfunction sat32ShiftLeft

```

Exceptions:

Reserved Instruction, DSP Disabled

Programming Notes:

To do a logical left shift of a word in a register without saturation, use the MIPS32 SLL instruction.

SHLL_S.W

Shift Left Logical Word with Saturation

SHRA[_R].QB**Shift Right Arithmetic Vector of Four Bytes**

31	26	25	21	20	16	15	13	12	11	9	8	6	5	3	2	0
SHRA.QB																
P32A 001000	rt			rs			sa		0	000		111		111		111
SHRA_R.QB																
P32A 001000	rt			rs			sa		1	000		111		111		111
6	5			5			3		1	3		3		3		3

Format: SHRA[_R].QB

SHRA.QB rt, rs, sa

DSP-R2

SHRA_R.QB rt, rs, sa

DSP-R2**Purpose:** Shift Right Arithmetic Vector of Four Bytes

To execute an arithmetic right shift on four independent bytes by a fixed number of bits.

Description: $rt \leftarrow \text{round}(rs_{31..24} \gg sa) \parallel \text{round}(rs_{23..16} \gg sa) \parallel \text{round}(rs_{15..8} \gg sa) \parallel \text{round}(rs_{7..0} \gg sa)$

The four byte elements in register *rs* are each shifted right arithmetically by *sa* bits, then written to the corresponding vector elements in destination register *rt*. The *sa* argument is interpreted as an unsigned three-bit integer taking values from zero to seven.

In the rounding variant of the instruction, a value of 1 is added at the most significant discarded bit position of each result prior to writing the rounded result to the destination register.

Restrictions:

No data-dependent exceptions are possible.

The operands must be a value in the specified format. If they are not, the results are **UNPREDICTABLE** and the values of the operand vectors become **UNPREDICTABLE**.

Operation:

```

SHRA.QB
  ValidateAccessToDSP2Resources()
  tempD7..0 ← ( GPR[rs]31 )sa || GPR[rs]31..24+sa )
  tempC7..0 ← ( GPR[rs]23 )sa || GPR[rs]23..16+sa )
  tempB7..0 ← ( GPR[rs]15 )sa || GPR[rs]15..8+sa )
  tempA7..0 ← ( GPR[rs]7 )sa || GPR[rs]7..sa )
  GPR[rt]31..0 ← tempD7..0 || tempC7..0 || tempB7..0 || tempA7..0

SHRA_R.QB
  ValidateAccessToDSP2Resources()
  if ( sa2..0 = 0 ) then
    tempD7..0 ← GPR[rs]31..24
    tempC7..0 ← GPR[rs]23..16
    tempB7..0 ← GPR[rs]15..8
    tempA7..0 ← GPR[rs]7..0
  else
    tempD8..0 ← ( GPR[rs]31 )sa || GPR[rs]31..24+sa-1 ) + 1
    tempC8..0 ← ( GPR[rs]23 )sa || GPR[rs]23..16+sa-1 ) + 1
    tempB8..0 ← ( GPR[rs]15 )sa || GPR[rs]15..8+sa-1 ) + 1
    tempA8..0 ← ( GPR[rs]7 )sa || GPR[rs]7..sa-1 ) + 1
  endif
  GPR[rt]31..0 ← tempD8..1 || tempC8..1 || tempB8..1 || tempA8..1

```


SHRA[_R].QB

Shift Right Arithmetic Vector of Four Bytes

endif

Exceptions:

Reserved Instruction, DSP Disabled

SHRA[_R].PH**Shift Right Arithmetic Vector Pair Halfwords**

31	26	25	21	20	16	15	12	11	10	9	3	2	0
SHRA.PH													
P32A 001000		rt		rs		sa		x	0	1100110			101
SHRA_R.PH													
P32A 001000		rt		rs		sa		x	1	1100110			101
6		5		5		4		1	1	7			3

Format: SHRA[_R].PH
 SHRA.PH rt, rs, sa
 SHRA_R.PH rt, rs, sa

DSP
DSP

Purpose: Shift Right Arithmetic Vector Pair Halfwords

Element-wise arithmetic right-shift of two independent halfwords in a vector data type by a fixed number of bits, with optional rounding.

Description: $rt \leftarrow \text{rnd16}(rs_{31..16} \gg sa) \parallel \text{rnd16}(rs_{15..0} \gg sa)$

The two halfword values in register *rt* are each independently shifted right by *sa* bits, with each value's original sign bit duplicated into the *sa* most-significant bits emptied by the shift.

In the non-rounding variant of this instruction, the two independent results are then written to the corresponding halfword elements of destination register *rd*.

In the rounding variant of the instruction, a 1 is added at the most-significant discarded bit position before the results are written to destination register *rd*.

Restrictions:

No data-dependent exceptions are possible.

The operands must be values in the specified format. If they are not, the results are **UNPREDICTABLE** and the values of the operand vectors become **UNPREDICTABLE**.

Operation:

```
SHRA.PH
  ValidateAccessToDSPResources()
  tempB15..0 ← shift16RightArithmetic( GPR[rs]31..16, sa )
  tempA15..0 ← shift16RightArithmetic( GPR[rs]15..0, sa )
  GPR[rt]31..0 ← tempB15..0 || tempA15..0

SHRA_R.PH
  ValidateAccessToDSPResources()
  tempB15..0 ← rnd16ShiftRightArithmetic( GPR[rs]31..16, sa )
  tempA15..0 ← rnd16ShiftRightArithmetic( GPR[rs]15..0, sa )
  GPR[rt]31..0 ← tempB15..0 || tempA15..0

function shift16RightArithmetic( a15..0, s3..0 )
  if ( s3..0 = 0 ) then
    temp15..0 ← a15..0
  else
    sign ← a15
    temp15..0 ← ( signs || a15..s )
  endif
  return temp15..0
endfunction shift16RightArithmetic
```

SHRA[_R].PH**Shift Right Arithmetic Vector Pair Halfwords**

```

function rnd16ShiftRightArithmetic( a15..0, s3..0 )
  if ( s3..0 = 0 ) then
    temp16..0 ← ( a15..0 || 0 )
  else
    sign ← a15
    temp16..0 ← ( signs || a15..s-1 )
  endif
  temp16..0 ← temp + 1
  return temp16..1
endfunction rnd16ShiftRightArithmetic

```

Exceptions:

Reserved Instruction, DSP Disabled

SHRAV[_R].PH**Shift Right Arithmetic Variable Vector Pair Halfwords**

31	26	25	21	20	16	15	11	10	9	3	2	0	
SHRAV.PH													
P32A 001000	rt				rs				rd	0	0110001		101
SHRAV_R.PH													
P32A 001000	rt				rs				rd	1	0110001		101
6	5				5				5	1	7		3

Format: SHRAV[_R].PH

SHRAV.PH rd, rt, rs

SHRAV_R.PH rd, rt, rs

DSP**DSP****Purpose:** Shift Right Arithmetic Variable Vector Pair Halfwords

Element-wise arithmetic right shift of two independent halfwords in a vector data type by a variable number of bits, with optional rounding.

Description: $rd \leftarrow \text{rnd16}(rt_{31..16} \gg rs_{3..0}) \ || \ \text{rnd16}(rt_{15..0} \gg rs_{3..0})$

The two halfword values in register *rt* are each independently shifted right, with each value's original sign bit duplicated into the most-significant bits emptied by the shift. In the non-rounding variant of this instruction, the two independent results are then written to the corresponding halfword elements of destination register *rd*.

In the rounding variant of this instruction, a 1 is added at the most-significant discarded bit position before the results are written to destination register *rd*.

The shift amount *sa* is given by the four least-significant bits of register *rs*; the remaining bits of *rs* are ignored.

Restrictions:

No data-dependent exceptions are possible.

The operands must be values in the specified format. If they are not, the results are **UNPREDICTABLE** and the values of the operand vectors become **UNPREDICTABLE**.

Operation:

```

SHRAV.PH
    ValidateAccessToDSPResources()
    tempB15..0 ← shift16RightArithmetic( GPR[rt]31..16, GPR[rs]3..0 )
    tempA15..0 ← shift16RightArithmetic( GPR[rt]15..0, GPR[rs]3..0 )
    GPR[rd]31..0 ← tempB15..0 || tempA15..0

SHRAV_R.PH
    ValidateAccessToDSPResources()
    tempB15..0 ← rnd16ShiftRightArithmetic( GPR[rt]31..16, GPR[rs]3..0 )
    tempA15..0 ← rnd16ShiftRightArithmetic( GPR[rt]15..0, GPR[rs]3..0 )
    GPR[rd]31..0 ← tempB15..0 || tempA15..0

```

Exceptions:

Reserved Instruction, DSP Disabled

SHRAV[_R].PH

Shift Right Arithmetic Variable Vector Pair Halfwords

SHRAV[_R].QB**Shift Right Arithmetic Variable Vector of Four Bytes**

31		26 25		21 20		16 15		11 10 9			3 2		0
SHRAV.QB													
P32A 001000		rt		rs		rd		0	0111001			101	
SHRAV_R.QB													
P32A 001000		rt		rs		rd		1	0111001			101	
6		5		5		5		1	7			3	

Format: SHRAV[_R].QB

SHRAV.QB rd, rt, rs

DSP-R2

SHRAV_R.QB rd, rt, rs

DSP-R2**Purpose:** Shift Right Arithmetic Variable Vector of Four Bytes

To execute an arithmetic right shift on four independent bytes by a variable number of bits.

Description: $rd \leftarrow \text{round}(rt_{31..24} \gg rs_{2..0}) \parallel \text{round}(rt_{23..16} \gg rs_{2..0}) \parallel \text{round}(rt_{15..8} \gg rs_{2..0}) \parallel \text{round}(rt_{7..0} \gg rs_{2..0})$

The four byte elements in register *rt* are each shifted right arithmetically by *sa* bits, then written to the corresponding byte elements in destination register *rd*. The *sa* argument is provided by the three least-significant bits of register *rs*, interpreted as an unsigned three-bit integer taking values from zero to seven. The remaining bits of *rs* are ignored.

In the rounding variant of the instruction, a value of 1 is added at the most significant discarded bit position of each result prior to writing the rounded result to the destination register.

Restrictions:

No data-dependent exceptions are possible.

The operands must be a value in the specified format. If they are not, the results are **UNPREDICTABLE** and the values of the operand vectors become **UNPREDICTABLE**.

Operation:

```

SHRAV.QB
    ValidateAccessToDSP2Resources()
    sa2..0 ← GPR[rs]2..0
    if ( sa2..0 = 0 ) then
        tempD7..0 ← GPR[rt]31..24
        tempC7..0 ← GPR[rt]23..16
        tempB7..0 ← GPR[rt]15..8
        tempA7..0 ← GPR[rt]7..0
    else
        tempD7..0 ← ( GPR[rt]31 )sa || GPR[rt]31..24+sa
        tempC7..0 ← ( GPR[rt]23 )sa || GPR[rt]23..16+sa
        tempB7..0 ← ( GPR[rt]15 )sa || GPR[rt]15..8+sa
        tempA7..0 ← ( GPR[rt]7 )sa || GPR[rt]7..sa
    endif
    GPR[rd]31..0 ← tempD7..0 || tempC7..0 || tempB7..0 || tempA7..0

SHRAV_R.QB
    ValidateAccessToDSP2Resources()
    sa2..0 ← GPR[rs]2..0
    if ( sa2..0 = 0 ) then
        tempD8..0 ← ( GPR[rt]31..24 || 0 )
        tempC8..0 ← ( GPR[rt]23..16 || 0 )

```

SHRAV[_R].QB**Shift Right Arithmetic Variable Vector of Four Bytes**

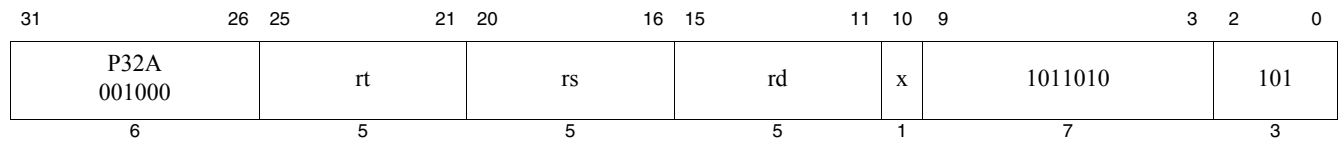
```

    tempB8..0 ← ( GPR[rt]15..8 || 0 )
    tempA8..0 ← ( GPR[rt]7..0 || 0 )
else
    tempD8..0 ← ( GPR[rt]31sa || GPR[rt]31..24+sa-1 ) + 1
    tempC8..0 ← ( GPR[rt]23sa || GPR[rt]23..16+sa-1 ) + 1
    tempB8..0 ← ( GPR[rt]15sa || GPR[rt]15..8+sa-1 ) + 1
    tempA8..0 ← ( GPR[rt]7sa || GPR[rt]7..sa-1 ) + 1
endif
GPR[rd]31..0 ← tempD8..1 || tempC8..1 || tempB8..1 || tempA8..1

```

Exceptions:

Reserved Instruction, DSP Disabled

SHRAV_R.W**Shift Right Arithmetic Variable Word with Rounding****Format:** SHRAV_R.W rd, rt, rs**DSP****Purpose:** Shift Right Arithmetic Variable Word with Rounding

Arithmetic right shift with rounding of a signed 32-bit word by a variable number of bits.

Description: $rd \leftarrow \text{rnd32}(rt_{31..0} \gg rs_{4..0})$

The word value in register *rt* is shifted right, with the value's original sign bit duplicated into the most-significant bits emptied by the shift. A 1 is then added at the most-significant discarded bit position before the result is written to destination register *rd*.

The shift amount *sa* is given by the five least-significant bits of register *rs*; the remaining bits of *rs* are ignored.

Restrictions:

No data-dependent exceptions are possible.

The operands must be values in the specified format. If they are not, the results are **UNPREDICTABLE** and the values of the operand vectors become **UNPREDICTABLE**.

Operation:

```

ValidateAccessToDSPResources()
temp31..0 ← rnd32ShiftRightArithmetic( GPR[rt]31..0, GPR[rs]4..0 )
GPR[rd]31..0 ← temp31..0

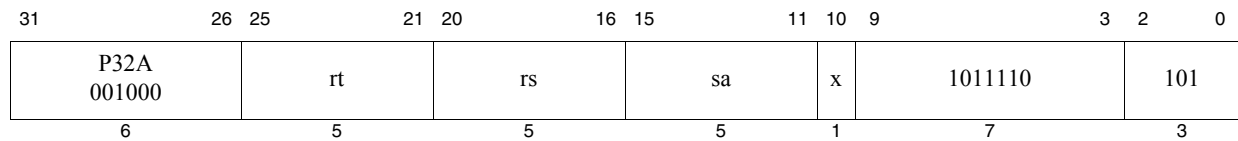
```

Exceptions:

Reserved Instruction, DSP Disabled

SHRAV_R.W

Shift Right Arithmetic Variable Word with Rounding

SHRA_R.W**Shift Right Arithmetic Word with Rounding****Format:** SHRA_R.W rt, rs, sa**DSP****Purpose:** Shift Right Arithmetic Word with Rounding

To execute an arithmetic right shift with rounding on a word by a fixed number of bits.

Description: $rt \leftarrow \text{rnd32}(rs_{31:0} \gg sa)$

The word in register *rs* is shifted right by *sa* bits, and the sign bit is duplicated into the *sa* bits emptied by the shift. The shifted result is then rounded by adding a 1 bit to the most-significant discarded bit. The rounded result is then written to the destination register *rt*.

Restrictions:

No data-dependent exceptions are possible.

The operands must be values in the specified format. If they are not, the results are **UNPREDICTABLE** and the values of the operand vectors become **UNPREDICTABLE**.

Operation:

```

ValidateAccessToDSPResources()
temp31..0 ← rnd32ShiftRightArithmetic( GPR[rt]31..0, sa4..0 )
GPR[rt]31..0 ← temp32..1

function rnd32ShiftRightArithmetic( a31..0, s4..0 )
  if ( s4..0 = 0 ) then
    temp32..0 ← ( a31..0 || 0 )
  else
    sign ← a31
    temp32..0 ← ( signs || a31..s-1 )
  endif
  temp32..0 ← temp + 1
  return temp32..1
endfunction rnd32ShiftRightArithmetic

```

Exceptions:

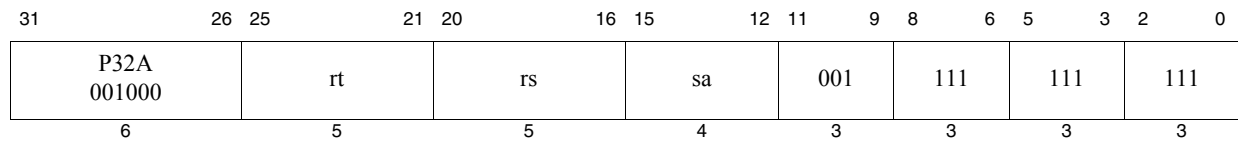
Reserved Instruction, DSP Disabled

Programming Notes:

To do an arithmetic right shift of a word in a register without rounding, use the MIPS32 SRA instruction.

SHRA_R.W

Shift Right Arithmetic Word with Rounding

SHRL.PH**Shift Right Logical Two Halfwords****Format:** SHRL.PH *rt*, *rs*, *sa***DSP-R2****Purpose:** Shift Right Logical Two Halfwords

To execute a right shift of two independent halfwords in a vector data type by a fixed number of bits.

Description: $rt \leftarrow (rs_{31..16} \gg sa) \parallel (rs_{15..0} \gg sa)$

The two halfwords in register *rs* are independently logically shifted right, inserting zeros into the bit positions emptied by the shift. The two halfword results are then written to the corresponding halfword elements in destination register *rt*.

The shift amount is provided by the *sa* field, which is interpreted as a four bit unsigned integer taking values between 0 and 15.

Restrictions:

No data-dependent exceptions are possible.

The operands must be a value in the specified format. If they are not, the results are **UNPREDICTABLE** and the values of the operand vectors become **UNPREDICTABLE**.

Operation:

```

ValidateAccessToDSP2Resources()
tempB15..0 ← 0sa || GPR[rs]31..sa+16
tempA15..0 ← 0sa || GPR[rs]15..sa
GPR[rt]31..0 ← tempB15..0 || tempA15..0

```

Exceptions:

Reserved Instruction, DSP Disabled

SHRL.PH

Shift Right Logical Two Halfwords

SHRL.QB**Shift Right Logical Vector Quad Bytes**

31	26	25	21	20	16	15	12	11	9	8	6	5	3	2	0
P32A 001000						rt	rs	sa	1	100	001	111	111		
6						5	5	4		3	3	3	3		

Format: SHRL.QB *rt*, *rs*, *sa***DSP****Purpose:** Shift Right Logical Vector Quad Bytes

Element-wise logical right shift of four independent bytes in a vector data type by a fixed number of bits.

Description: $rt \leftarrow rs_{31..24} \gg sa \parallel (rs_{23..16} \gg sa) \parallel (rs_{15..8} \gg sa) \parallel (rs_{7..0} \gg sa)$

The four byte values in register *rs* are each independently shifted right by *sa* bits and the *sa* most-significant bits of each value are set to zero. The four independent results are then written to the corresponding byte elements of destination register *rt*.

Restrictions:

No data-dependent exceptions are possible.

The operands must be values in the specified format. If they are not, the results are **UNPREDICTABLE** and the values of the operand vectors become **UNPREDICTABLE**.

Operation:

```

ValidateAccessToDSPResources()
tempD7..0 ← shift8Right( GPR[rs]31..24, sa )
tempC7..0 ← shift8Right( GPR[rs]23..16, sa )
tempB7..0 ← shift8Right( GPR[rs]15..8, sa )
tempA7..0 ← shift8Right( GPR[rs]7..0, sa )
GPR[rt]31..0 ← tempD7..0 || tempC7..0 || tempB7..0 || tempA7..0

function shift8Right( a7..0, s2..0 )
  if ( s2..0 = 0 ) then
    temp7..0 ← a7..0
  else
    temp7..0 ← ( 0s || a7..s )
  endif
  return temp7..0
endfunction shift8Right

```

Exceptions:

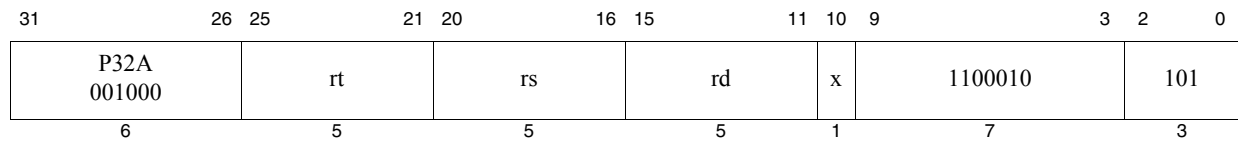
Reserved Instruction, DSP Disabled

Programming Notes:

To do a logical left shift of a word in a register without saturation, use the MIPS32 SLL instruction.

SHRL.QB

Shift Right Logical Vector Quad Bytes

SHRLV.PH**Shift Variable Right Logical Pair of Halfwords****Format:** SHRLV.PH rd, rt, rs**DSP-R2****Purpose:** Shift Variable Right Logical Pair of Halfwords

To execute a right shift of two independent halfwords in a vector data type by a variable number of bits.

Description: $rd \leftarrow (rt_{31..16} \gg rs_{3..0}) \parallel (rt_{15..0} \gg rs_{3..0})$

The two halfwords in register *rt* are independently logically shifted right, inserting zeros into the bit positions emptied by the shift. The two halfword results are then written to the corresponding halfword elements in destination register *rd*.

The shift amount is provided by the four least-significant bits of register *rs*, which is interpreted as a four bit unsigned integer taking values between 0 and 15. The remaining bits of *rs* are ignored.

Restrictions:

No data-dependent exceptions are possible.

The operands must be a value in the specified format. If they are not, the results are **UNPREDICTABLE** and the values of the operand vectors become **UNPREDICTABLE**.

Operation:

```

ValidateAccessToDSP2Resources()
sa3..0 ← GPR[rs]3..0
tempB15..0 ← 0sa || GPR[rt]31..sa+16
tempA15..0 ← 0sa || GPR[rt]15..sa
GPR[rd]31..0 ← tempB15..0 || tempA15..0

```

Exceptions:

Reserved Instruction, DSP Disabled

SHRLV.PH

Shift Variable Right Logical Pair of Halfwords

SHRLV.QB**Shift Right Logical Variable Vector Quad Bytes**

31	26	25	21	20	16	15	11	10	9	3	2	0	
P32A 001000			rt		rs		rd		x	1101010			101
6			5		5		5		1	7			3

Format: SHRLV.QB rd, rt, rs**DSP****Purpose:** Shift Right Logical Variable Vector Quad Bytes

Element-wise logical right shift of four independent bytes in a vector data type by a variable number of bits.

Description: $rd \leftarrow (rt_{31..24} \gg rs_{2..0}) \parallel (rt_{23..16} \gg rs_{2..0}) \parallel (rt_{15..8} \gg rs_{2..0}) \parallel (rt_{7..0} \gg rs_{2..0})$

The four byte values in register *rt* are each independently shifted right, inserting zeros into the most-significant bit positions emptied by the shift. The four independent results are then written to the corresponding byte elements of destination register *rd*.

The three least-significant bits of *rs* provide the shift value, interpreted as an unsigned integer; the remaining bits of *rs* are ignored.

Restrictions:

No data-dependent exceptions are possible.

The operands must be values in the specified format. If they are not, the results are **UNPREDICTABLE** and the values of the operand vectors become **UNPREDICTABLE**.

Operation:

```

ValidateAccessToDSPResources()
tempD7..0 ← shift8Right( GPR[rt]31..24, GPR[rs]2..0 )
tempC7..0 ← shift8Right( GPR[rt]23..16, GPR[rs]2..0 )
tempB7..0 ← shift8Right( GPR[rt]15..8, GPR[rs]2..0 )
tempA7..0 ← shift8Right( GPR[rt]7..0, GPR[rs]2..0 )
GPR[rd]31..0 ← tempD7..0 || tempC7..0 || tempB7..0 || tempA7..0

```

Exceptions:

Reserved Instruction, DSP Disabled

SHRLV.QB

Shift Right Logical Variable Vector Quad Bytes

SUBQ[_S].PH**Subtract Fractional Halfword Vector**

31	26	25	21	20	16	15	11	10	9	3	2	0
SUBQ.PH												
P32A 001000	rt		rs		rd		0	1000001			101	
SUBQ_S.PH												
P32A 001000	rt		rs		rd		1	1000001			101	
6	5		5		5		1	7			3	

Format: SUBQ[_S].PH
 SUBQ.PH rd, rs, rt
 SUBQ_S.PH rd, rs, rt

DSP
DSP

Purpose: Subtract Fractional Halfword Vector

Element-wise subtraction of one vector of Q15 fractional halfword values from another to produce a vector of Q15 fractional halfword results, with optional saturation.

Description: $rd \leftarrow \text{sat16}(rs_{31..16} - rt_{31..16}) \parallel \text{sat16}(rs_{15..0} - rt_{15..0})$

The two fractional halfwords in register *rt* are subtracted from the corresponding fractional halfword elements in register *rs*.

For the non-saturating version of this instruction, each result is written to the corresponding element in register *rd*. In the case of overflow or underflow, the result modulo 2 is written to register *rd*.

For the saturating version of the instruction, the subtraction is performed using signed saturating arithmetic. If the operation results in an overflow or an underflow, the result is clamped to either the largest representable value (0x7FFF hexadecimal) or the smallest representable value (0x8000 hexadecimal), respectively, before being written to the destination register *rd*.

For both instructions, if any of the individual subtractions result in underflow, overflow, or saturation, a 1 is written to bit 20 in the *DSPControl* register within the *ouflag* field.

Restrictions:

No data-dependent exceptions are possible.

The operands must be values in the specified format. If they are not, the results are **UNPREDICTABLE** and the values of the operand vectors become **UNPREDICTABLE**.

Operation:

SUBQ.PH:

```
ValidateAccessToDSPResources()
tempB15..0 ← subtract16( GPR[rs]31..16 , GPR[rt]31..16 )
tempA15..0 ← subtract16( GPR[rs]15..0 , GPR[rt]15..0 )
GPR[rd]31..0 ← tempB15..0 || tempA15..0
```

SUBQ_S.PH:

```
ValidateAccessToDSPResources()
tempB15..0 ← sat16Subtract( GPR[rs]31..16 , GPR[rt]31..16 )
tempA15..0 ← sat16Subtract( GPR[rs]15..0 , GPR[rt]15..0 )
GPR[rd]31..0 ← tempB15..0 || tempA15..0
```

```
function subtract16( a15..0, b15..0 )
  temp16..0 ← ( a15 || a15..0 ) - ( b15 || b15..0 )
  if ( temp16 ≠ temp15 ) then
    DSPControlouflag:20 ← 1
  endif
```

SUBQ[_S].PH**Subtract Fractional Halfword Vector**

```

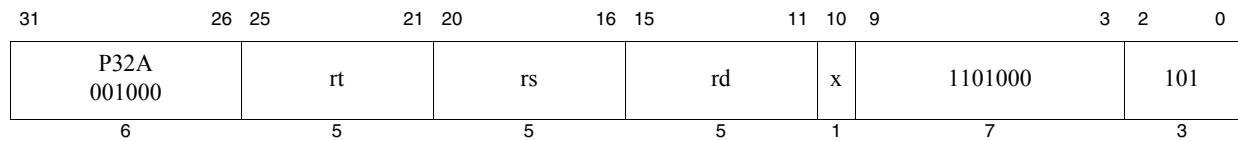
    return temp15..0
endfunction subtract16

function sat16Subtract( a15..0, b15..0 )
    temp16..0 ← ( a15 || a15..0 ) - ( b15 || b15..0 )
    if ( temp16 ≠ temp15 ) then
        if ( temp16 = 0 ) then
            temp ← 0x7FFF
        else
            temp ← 0x8000
        endif
        DSPControlouflag:20 ← 1
    endif
    return temp15..0
endfunction sat16Subtract

```

Exceptions:

Reserved Instruction, DSP Disabled

SUBQ_S.W**Subtract Fractional Word****Format:** SUBQ_S.W rd, rs, rt**DSP****Purpose:** Subtract Fractional Word

One Q31 fractional word is subtracted from another to produce a Q31 fractional result, with saturation.

Description: $rd \leftarrow \text{sat32}(rs_{31..0} - rt_{31..0})$

The Q31 fractional word in register *rt* is subtracted from the corresponding fractional word in register *rs*, and the 32-bit result is written to destination register *rd*. The subtraction is performed using signed saturating arithmetic. If the operation results in an overflow or an underflow, the result is clamped to either the largest representable value (0x7FFFFFFF hexadecimal) or the smallest representable value (0x80000000 hexadecimal), respectively, before being sign-extended and written to the destination register *rd*.

If the subtraction results in underflow, overflow, or saturation, a 1 is written to bit 20 in the *DSPControl* register within the *ouflag* field.

Restrictions:

No data-dependent exceptions are possible.

The operands must be values in the specified format. If they are not, the results are **UNPREDICTABLE** and the values of the operand vectors become **UNPREDICTABLE**.

Operation:

```

ValidateAccessToDSPResources()
temp31..0 ← sat32Subtract( GPR[rs]31..0 , GPR[rt]31..0 )
GPR[rd]31..0 ← temp31..0

function sat32Subtract( a31..0, b31..0 )
    temp32..0 ← ( a31 || a31..0 ) - ( b31 || b31..0 )
    if ( temp32 ≠ temp31 ) then
        if ( temp32 = 0 ) then
            temp31..0 ← 0x7FFFFFFF
        else
            temp31..0 ← 0x80000000
        endif
        DSPControlouflag:20 ← 1
    endif
    return temp31..0
endfunction sat32Subtract

```

Exceptions:

Reserved Instruction, DSP Disabled

SUBQ_S.W

Subtract Fractional Word

SUBQH[_R].PH**Subtract Fractional Halfword Vectors And Shift Right to Halve Results**

31	26	25	21	20	16	15	11	10	9	3	2	0
SUBQH.PH												
P32A 001000	rt		rs		rd		0	1001001			101	
SUBQH_R.PH												
P32A 001000	rt		rs		rd		1	1001001			101	
6		5		5		5		1	7			3

Format: SUBQH[_R].PH

SUBQH.PH rd, rs, rt

DSP-R2

SUBQH_R.PH rd, rs, rt

DSP-R2**Purpose:** Subtract Fractional Halfword Vectors And Shift Right to Halve Results

Element-wise fractional subtraction of halfword vectors, with a right shift by one bit to halve each result, with optional rounding.

Description: $rd \leftarrow \text{round}((rs_{31..16} - rt_{31..16}) \gg 1) \parallel \text{round}((rs_{15..0} - rt_{15..0}) \gg 1)$

Each element from the two halfword values in register *rt* is subtracted from the corresponding halfword element in register *rs* to create an interim 17-bit result.

In the non-rounding instruction variant, each interim result is then shifted right by one bit before being written to the corresponding halfword element of destination register *rd*.

In the rounding version of the instruction, a value of 1 is added at the least-significant bit position of each interim result; the interim result is then right-shifted by one bit and written to the destination register.

This instruction does not modify the *DSPControl* register.

Restrictions:

No data-dependent exceptions are possible.

The operands must be a value in the specified format. If they are not, the results are **UNPREDICTABLE** and the values of the operand vectors become **UNPREDICTABLE**.

Operation:

```

ADDQH.PH
    ValidateAccessToDSP2Resources()
    tempB15..0 ← rightShift1SubQ16( GPR[rs]31..16 , GPR[rt]31..16 )
    tempA15..0 ← rightShift1SubQ16( GPR[rs]15..0 , GPR[rt]15..0 )
    GPR[rd]31..0 ← tempB15..0 || tempA15..0

ADDQH_R.PH
    ValidateAccessToDSP2Resources()
    tempB15..0 ← roundRightShift1SubQ16( GPR[rs]31..16 , GPR[rt]31..16 )
    tempA15..0 ← roundRightShift1SubQ16( GPR[rs]15..0 , GPR[rt]15..0 )
    GPR[rd]31..0 ← tempB15..0 || tempA15..0

function rightShift1SubQ16( a15..0 , b15..0 )
    temp16..0 ← (( a15 || a15..0 ) - ( b15 || b15..0 ))
    return temp16..1
endfunction rightShift1SubQ16

function roundRightShift1SubQ16( a15..0 , b15..0 )
    temp16..0 ← (( a15 || a15..0 ) - ( b15 || b15..0 ))
    temp16..0 ← temp16..0 + 1

```


SUBQH[_R].PH

Subtract Fractional Halfword Vectors And Shift Right to Halve Results

```
    return temp16..1  
endfunction roundRightShift1SubQ16
```

Exceptions:

Reserved Instruction, DSP Disabled

SUBQH[_R].W**Subtract Fractional Words And Shift Right to Halve Results**

31	26	25	21	20	16	15	11	10	9	3	2	0
SUBQH.W												
P32A 001000		rt		rs		rd		0	1010001			101
SUBQH_R.W												
P32A 001000		rt		rs		rd		1	1010001			101
6		5		5		5		1	7			3

Format: SUBQH[_R].W
 SUBQH.W rd, rs, rt
 SUBQH_R.W rd, rs, rt

DSP-R2**DSP-R2**

Purpose: Subtract Fractional Words And Shift Right to Halve Results

Fractional subtraction of word vectors, with a right shift by one bit to halve the result, with optional rounding.

Description: $rd \leftarrow \text{round}((rs_{31..0} - rt_{31..0}) \gg 1)$

The word in register *rt* is subtracted from the word in register *rs* to create an interim 33-bit result.

In the non-rounding instruction variant, the interim result is then shifted right by one bit before being written to the destination register *rd*.

In the rounding version of the instruction, a value of 1 is added at the least-significant bit position of the interim result; the interim result is then right-shifted by one bit and written to the destination register.

This instruction does not modify the *DSPControl* register.

Restrictions:

No data-dependent exceptions are possible.

The operands must be a value in the specified format. If they are not, the results are **UNPREDICTABLE** and the values of the operand vectors become **UNPREDICTABLE**.

Operation:

```

ADDQH.W
  ValidateAccessToDSP2Resources()
  tempA31..0 ← rightShift1SubQ32( GPR[rs]31..0 , GPR[rt]31..0 )
  GPR[rd]31..0 ← tempA31..0

ADDQH_R.W
  ValidateAccessToDSP2Resources()
  tempA31..0 ← roundRightShift1SubQ32( GPR[rs]31..0 , GPR[rt]31..0 )
  GPR[rd]31..0 ← tempA31..0

function rightShift1SubQ32( a31..0 , b31..0 )
  temp32..0 ← (( a31 || a31..0 ) - ( b31 || b31..0 ))
  return temp32..1
endfunction rightShift1SubQ32

function roundRightShift1SubQ32( a31..0 , b31..0 )
  temp32..0 ← (( a31 || a31..0 ) - ( b31 || b31..0 ))
  temp32..0 ← temp32..0 + 1
  return temp32..1
endfunction roundRightShift1SubQ32

```

SUBQH[_R].W

Subtract Fractional Words And Shift Right to Halve Results

Exceptions:

Reserved Instruction, DSP Disabled

SUBU[_S].PH**Subtract Unsigned Integer Halfwords**

31	26	25	21	20	16	15	11	10	9	3	2	0	
SUBU.PH													
P32A 001000	rt		rs		rd		0	1100001			101		
SUBU_S.PH													
P32A 001000	rt		rs		rd		1	1100001			101		
6		5		5		5		1	7			3	

Format: SUBU[_S].PH
 SUBU.PH rd, rs, rt
 SUBU_S.PH rd, rs, rt

DSP-R2**DSP-R2****Purpose:** Subtract Unsigned Integer Halfwords

Element-wise subtraction of pairs of unsigned integer halfwords, with optional saturation.

Description: $rd \leftarrow \text{sat16}(rs_{31..16} - rt_{31..16}) \parallel \text{sat16}(rs_{15..0} - rt_{15..0})$

The two unsigned integer halfwords in register *rs* are subtracted from the corresponding unsigned integer halfwords in register *rt*. The unsigned results are then written to the corresponding element in destination register *rd*.

In the saturating version of the instruction, if either subtraction results in an underflow the result is clamped to the minimum unsigned integer halfword value (0x0000 hexadecimal), before being written to the destination register *rd*.

For both instruction variants, if either subtraction causes an underflow the instruction writes a 1 to bit 20 in the *DSPControl* register in the *ouflag* field.

Restrictions:

No data-dependent exceptions are possible.

The operands must be a value in the specified format. If they are not, the results are **UNPREDICTABLE** and the values of the operand vectors become **UNPREDICTABLE**.

Operation:

```

SUBU.PH
  ValidateAccessToDSP2Resources()
  tempB15..0 ← subtractU16U16( GPR[rt]31..16 , GPR[rs]31..16 )
  tempA15..0 ← subtractU16U16( GPR[rt]15..0 , GPR[rs]15..0 )
  GPR[rd]31..0 ← tempB15..0 || tempA15..0

SUBU_S.PH
  ValidateAccessToDSPResources()
  tempB15..0 ← satU16SubtractU16U16( GPR[rt]31..16 , GPR[rs]31..16 )
  tempA15..0 ← satU16SubtractU16U16( GPR[rt]15..0 , GPR[rs]15..0 )
  GPR[rd]31..0 ← tempB15..0 || tempA15..0

function subtractU16U16( a15..0, b15..0 )
  temp16..0 ← ( 0 || a15..0 ) - ( 0 || b15..0 )
  if ( temp16 = 1 ) then
    DSPControlouflag:20 ← 1
  endif
  return temp15..0
endfunction subtractU16U16

```

SUBU[_S].PH**Subtract Unsigned Integer Halfwords**

```

function satU16SubtractU16U16( a15..0, b15..0 )
    temp16..0 ← ( 0 || a15..0 ) - ( 0 || b15..0 )
    if ( temp16 = 1 ) then
        temp15..0 ← 0x0000
        DSPControlouflag:20 ← 1
    endif
    return temp15..0
endfunction satU16SubtractU16U16

```

Exceptions:

Reserved Instruction, DSP Disabled

SUBU[_S].QB**Subtract Unsigned Quad Byte Vector**

31	26	25	21	20	16	15	11	10	9	3	2	0	
SUBU.QB													
P32A 001000	rt		rs		rd		0	1011001			101		
SUBU.QB													
P32A 001000	rt		rs		rd		1	1011001			101		
6		5		5		5		1	7			3	

Format: SUBU[_S].QB
 SUBU.QB rd, rs, rt
 SUBU_S.QB rd, rs, rt

DSP
DSP

Purpose: Subtract Unsigned Quad Byte Vector

Element-wise subtraction of one vector of unsigned byte values from another to produce a vector of unsigned byte results, with optional saturation.

Description: $rd \leftarrow \text{sat8}(rs_{31..24} - rt_{31..24}) \parallel \text{sat8}(rs_{23..16} - rt_{23..16}) \parallel \text{sat8}(rs_{15..8} - rt_{15..8}) \parallel \text{sat8}(rs_{7..0} - rt_{7..0})$

The four byte elements in *rt* are subtracted from the corresponding byte elements in register *rs*.

For the non-saturating version of the instruction, the result modulo 256 is written into the corresponding position in register *rd*.

For the saturating version of the instruction the subtraction is performed using unsigned saturating arithmetic. If the subtraction results in underflow, the value is clamped to the smallest representable value (0 decimal, 0x00 hexadecimal) before being written to the destination register *rd*.

For each instruction, if any of the individual subtractions result in underflow or saturation, a 1 is written to bit 20 in the *DSPControl* register within the *ouflag* field.

Restrictions:

No data-dependent exceptions are possible.

The operands must be values in the specified format. If they are not, the results are **UNPREDICTABLE** and the values of the operand vectors become **UNPREDICTABLE**.

Operation:

```

SUBU.QB:
    ValidateAccessToDSPResources()
    tempD7..0 ← subtractU8( GPR[rs]31..24 , GPR[rt]31..24 )
    tempC7..0 ← subtractU8( GPR[rs]23..16 , GPR[rt]23..16 )
    tempB7..0 ← subtractU8( GPR[rs]15..8 , GPR[rt]15..8 )
    tempA7..0 ← subtractU8( GPR[rs]7..0 , GPR[rt]7..0 )
    GPR[rd]31..0 ← tempD7..0 || tempC7..0 || tempB7..0 || tempA7..0

SUBU_S.QB:
    ValidateAccessToDSPResources()
    tempD7..0 ← satU8Subtract( GPR[rs]31..24 , GPR[rt]31..24 )
    tempC7..0 ← satU8Subtract( GPR[rs]23..16 , GPR[rt]23..16 )
    tempB7..0 ← satU8Subtract( GPR[rs]15..8 , GPR[rt]15..8 )
    tempA7..0 ← satU8Subtract( GPR[rs]7..0 , GPR[rt]7..0 )
    GPR[rd]31..0 ← tempD7..0 || tempC7..0 || tempB7..0 || tempA7..0

function subtractU8( a7..0, b7..0 )
    temp8..0 ← ( 0 || a7..0 ) - ( 0 || b7..0 )

```

SUBU[_S].QB**Subtract Unsigned Quad Byte Vector**

```

    if ( temp8 = 1 ) then
        DSPControlouflag:20 ← 1
    endif
    return temp7..0
endfunction subtractU8

function satU8Subtract( a7..0, b7..0 )
    temp8..0 ← ( 0 || a7..0 ) - ( 0 || b7..0 )
    if ( temp8 = 1 ) then
        temp7..0 ← 0x00
        DSPControlouflag:20 ← 1
    endif
    return temp7..0
endfunction satU8Subtract

```

Exceptions:

Reserved Instruction, DSP Disabled

SUBUH[_R].QB**Subtract Unsigned Bytes And Right Shift to Halve Results**

31	26	25	21	20	16	15	11	10	9	3	2	0
SUBUH.QB												
P32A 001000		rt		rs		rd		0	1101001		101	
SUBUH_R.QB												
P32A 001000		rt		rs		rd		1	1101001		101	
6		5		5		5		1	7		3	

Format: SUBUH[_R].QB
 SUBUH.QB rd, rs, rt
 SUBUH_R.QB rd, rs, rt

DSP-R2**DSP-R2**

Purpose: Subtract Unsigned Bytes And Right Shift to Halve Results

Element-wise subtraction of two vectors of unsigned bytes, with a one-bit right shift to halve results and optional rounding.

Description: $rd \leftarrow \text{round}((rs_{31..24} - rt_{31..24}) \gg 1) \parallel \text{round}((rs_{23..16} - rt_{23..16}) \gg 1) \parallel \text{round}((rs_{15..8} - rt_{15..8}) \gg 1) \parallel \text{round}((rs_{7..0} - rt_{7..0}) \gg 1)$

The four unsigned byte values in register *rt* are subtracted from the corresponding unsigned byte values in register *rs*. Each unsigned result is then halved by shifting right by one bit position. The byte results are then written to the corresponding elements of destination register *rd*.

In the rounding variant of the instruction, a value of 1 is added to the result of each subtraction at the discarded bit position before the right shift.

The results of this instruction never overflow; no bits of the *ouflag* field in the *DSPControl* register are written.

Restrictions:

No data-dependent exceptions are possible.

The operands must be a value in the specified format. If they are not, the results are **UNPREDICTABLE** and the values of the operand vectors become **UNPREDICTABLE**.

Operation:

```

SUBUH.QB
  ValidateAccessToDSPResources()
  tempD7..0 ← ( ( 0 || GPR[rs]31..24 ) - ( 0 || GPR[rt]31..24 ) ) >> 1
  tempC7..0 ← ( ( 0 || GPR[rs]23..16 ) - ( 0 || GPR[rt]23..16 ) ) >> 1
  tempB7..0 ← ( ( 0 || GPR[rs]15..8 ) - ( 0 || GPR[rt]15..8 ) ) >> 1
  tempA7..0 ← ( ( 0 || GPR[rs]7..0 ) - ( 0 || GPR[rt]7..0 ) ) >> 1
  GPR[rd]31..0 ← tempD7..0 || tempC7..0 || tempB7..0 || tempA7..0

SUBUH_R.QB
  ValidateAccessToDSPResources()
  tempD7..0 ← ( ( 0 || GPR[rs]31..24 ) - ( 0 || GPR[rt]31..24 ) + 1 ) >> 1
  tempC7..0 ← ( ( 0 || GPR[rs]23..16 ) - ( 0 || GPR[rt]23..16 ) + 1 ) >> 1
  tempB7..0 ← ( ( 0 || GPR[rs]15..8 ) - ( 0 || GPR[rt]15..8 ) + 1 ) >> 1
  tempA7..0 ← ( ( 0 || GPR[rs]7..0 ) - ( 0 || GPR[rt]7..0 ) + 1 ) >> 1
  GPR[rd]31..0 ← tempD7..0 || tempC7..0 || tempB7..0 || tempA7..0

```

Exceptions:

Reserved Instruction, DSP Disabled

SUBUH[_R].QB

Subtract Unsigned Bytes And Right Shift to Halve Results

WRDSP**Write Fields to DSPControl Register from a GPR**

31	26	25	21	20	14	13	12	11	9	8	6	5	3	2	0
P32A 001000	rt		mask			01	011		001	111		111			
6	5		7			2	3		3	3		3			

Format: WRDSP

WRDSP rt, mask

WRDSP rt

**DSP
Assembly Idiom****Purpose:** Write Fields to DSPControl Register from a GPR

To copy selected fields from the specified GPR to the special-purpose DSPControl register.

Description: $\text{DSPControl} \leftarrow \text{select}(\text{mask}, \text{GPR}[\text{rt}])$

Selected fields in the special register *DSPControl* are overwritten with the corresponding bits from the source GPR *rt*. Each of bits 0 through 5 of the *mask* operand corresponds to a specific field in the *DSPControl* register. A mask bit value of 1 indicates that the field will be overwritten using the bits from the same bit positions in register *rt*, and a mask bit value of 0 indicates that the corresponding field will be unchanged. Bits 6 through 9 of the *mask* operand are ignored.

The table below shows the correspondence between the bits in the *mask* operand and the fields in the *DSPControl* register; mask bit 0 is the least-significant bit in *mask*.

Bit	31	24	23	16	15	14	13	12	7	6	5	0
DSPControl field	ccond			ouflag			0	EFI	C	scount		pos
Mask bit	4			3				5	2	1		0

For example, to overwrite only the *scount* field in *DSPControl*, the value of the *mask* operand used will be 2 decimal (0x02 hexadecimal). After execution of the instruction, the *scount* field in *DSPControl* will have the value of bits 7 through 12 of the specified source register *rt* and the remaining bits in *DSPControl* are unmodified.

The one-operand version of the instruction provides a convenient assembly idiom that allows the programmer to write all the allowable fields in the *DSPControl* register from the source GPR, i.e., it is equivalent to specifying a *mask* value of 31 decimal (0x1F hexadecimal).

Restrictions:

No data-dependent exceptions are possible.

The operands must be values in the specified format. If they are not, the results are **UNPREDICTABLE** and the values of the operand vectors become **UNPREDICTABLE**.

Operation:

```

ValidateAccessToDSPResources()
newbits31..0 ← 032
overwrite31..0 ← 0xFFFFFFFF
if ( mask0 = 1 ) then
    overwrite5..0 ← 06
    newbits5..0 ← GPR[rt]5..0
endif
if ( mask1 = 1 ) then
    overwrite12..7 ← 06
    newbits12..7 ← GPR[rt]12..7
endif
if ( mask2 = 1 ) then
    overwrite13 ← 0

```

WRDSP**Write Fields to DSPControl Register from a GPR**

```

        newbits13 ← GPR[rt]13
    endif
    if ( mask3 = 1 ) then
        overwrite23..16 ← 08
        newbits23..16 ← GPR[rt]23..16
    endif
    if ( mask4 = 1 ) then
        overwrite31..24 ← 08
        newbits31..24 ← GPR[rt]31..24
    endif
    if ( mask5 = 1 ) then
        overwrite14 ← 0
        newbits14 ← GPR[rt]14
    endif

    DSPControl ← DSPControl and overwrite31..0
    DSPControl ← DSPControl or new31..0

```

Exceptions:

Reserved Instruction, DSP Disabled

Endian-Agnostic Reference to Register Elements

A.1 Using Endian-Agnostic Instruction Names

Certain instructions being proposed in the Module only operate on a subset of the operands in the register. In most cases, this is simply the left (**L**) or right (**R**) half of the register. Some instructions refer to the left alternating (**LA**) or right alternating (**RA**) elements of the register. But this type of reference does not take the endian-ness of the processor and memory into account. Since the DSP Module instructions do not take the endian-ness into account and simply use the left or right part of the register, this section describes a method by which users can take advantage of user-defined macros to translate the given instruction to the appropriate one for a given processor endian-ness.

An example is given below that uses actual element numbers in the mnemonics to be endian-agnostic.

In the microMIPS32 architecture, the following conventions could be used:

- PH0 refers to halfword element 0 (from a pair in the specified register).
- PH1 refers to halfword element 1.
- QB01 refers to byte elements 0 and 1 (from a quad in the specified register).
- QB23 refers to byte elements 2 and 3.
- QB02 refers to (even) byte elements 0 and 2.
- QB13 refers to (odd) byte elements 1 and 3.

The even and odd subsets are mainly used in storing, computing on, and loading complex numbers that have a real and imaginary part. If the real and imaginary parts of a complex number are stored in consecutive memory locations, then computations that involve only the real or only the imaginary parts must first extract these to a different register. This can most effectively be done using the even and odd formats of the relevant operations.

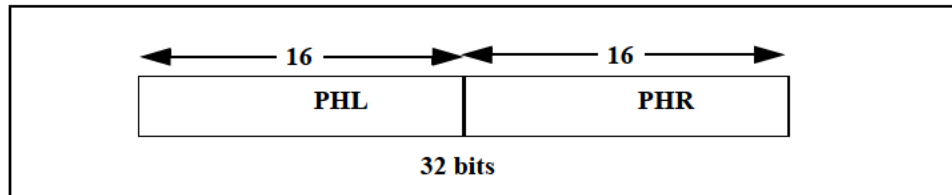
Note that these mnemonics are translated by the assembler to underlying real instructions that operate on absolute element positions in the register based on the endian-ness of the processor.

A.2 Mapping Endian-Agnostic Instruction Names to DSP Module Instructions

To illustrate this process, we will use one instruction as an example. This can be repeated for all the relevant instructions in the Module.

The **MULEQ_S** instruction multiplies fractional data operands to expanded full-size results in a destination register with optional saturation. Since the result occupies twice the width of the input operands, only half the operands from the source registers are operated on at a time. So the complete instruction mnemonic would be given as **MULEQ_S.W.PH0 rd, rs, rt** where the second part (after the first dot) indicates the size of the result, and the third part (after the second dot) indicates the element of the source register being used, which in this example is the 0th element. The real instructions that the hardware implements are **MULEQ_S.W.PHL** and **MULEQ_S.W.PHR** which operate on the left halfword element and the right halfword element respectively, of the given source registers, as shown in [Figure A.1](#). The user can map the user instruction (with **.PH0**) to the **MULEQ_S.W.PHL** real instruction if the processor is big-endian or to the real instruction **MULEQ_S.W.PHR** if the processor is little-endian.

Figure A.1 The Endian-Independent PHL and PHR Elements in a GPR for the microMIPS32 Architecture



Then **MULEQ_S.W.PH1 rd, rs, rt** instruction can be mapped to **MULEQ_S.W.PHR** if the processor is big-endian (see [Figure A.2](#)), and to **MULEQ_S.W.PHL** real instruction if the processor is little-endian (see [Figure A.3](#)).

Figure A.2 The Big-Endian PH0 and PH1 Elements in a GPR for the microMIPS32 Architecture

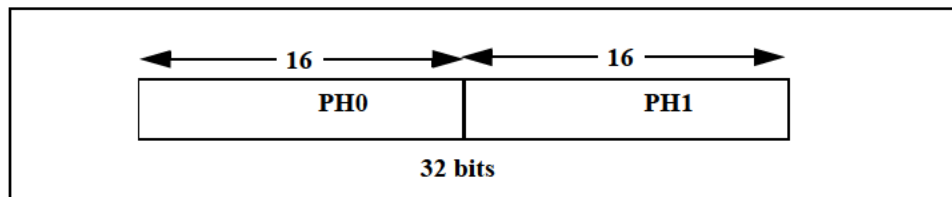
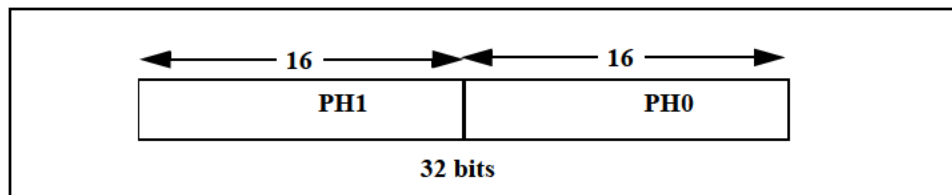


Figure A.3 The Little-Endian PH0 and PH1 Elements in a GPR for the microMIPS32 Architecture



To specify the even and odd type operations, a user instruction (to use odd elements) such as **PRECEQ_S.PH.QB02** (which precision expands the values) would be mapped to **PRECEQ_S.PH.QBLA** or **PRECEQ_S.PH.QBRA** depending on whether the endian-ness of the processor was big or little, respectively. (**LA** stands for left-alternating and **RA** for right-alternating).

Figure A.4 The Endian-Independent QBL and QBR Elements in a GPR for the microMIPS32 Architecture

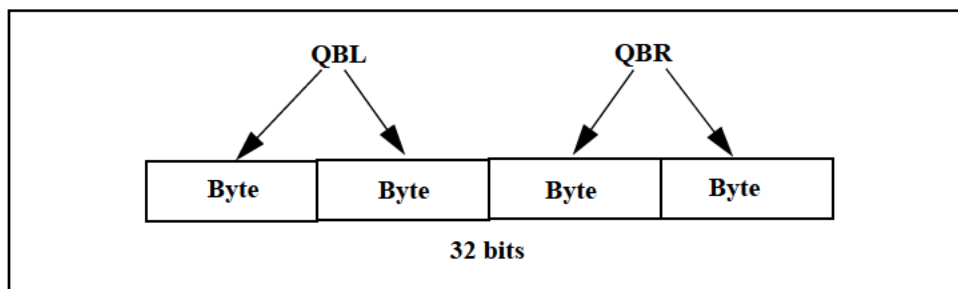
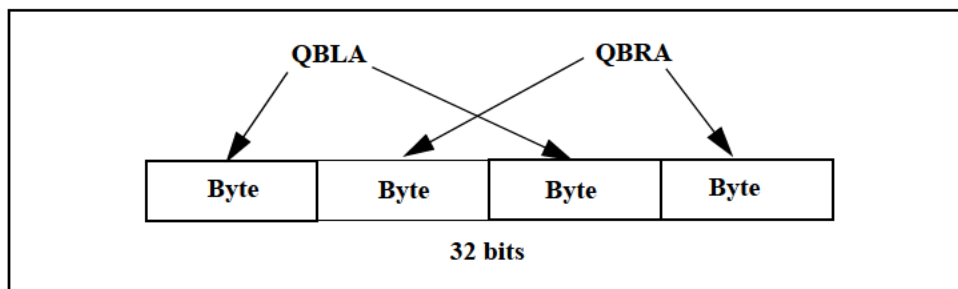


Figure A.5 The Endian-Independent QBLA and QBRA Elements in a GPR for the microMIPS32 Architecture



Revision History

Vertical change bars in the left page margin note the location of changes to this document since its last release.

NOTE: Change bars on figure titles are used to denote a potential change in the figure itself.

Version	Date	Comments
0.01	August 30, 2017	Initial revision.
0.02	November 17, 2017	Updated cover and formatting.
0.03	March 26, 2018	<ul style="list-style-type: none">• Fixed bits and typo in ADDQH_R.H• Fixed broken cross references in the Overview chapter.
0.04	April 27, 2018	Changed confidentiality level to Public.